



Arm® Cortex®-A78C Core

Revision: r0p1

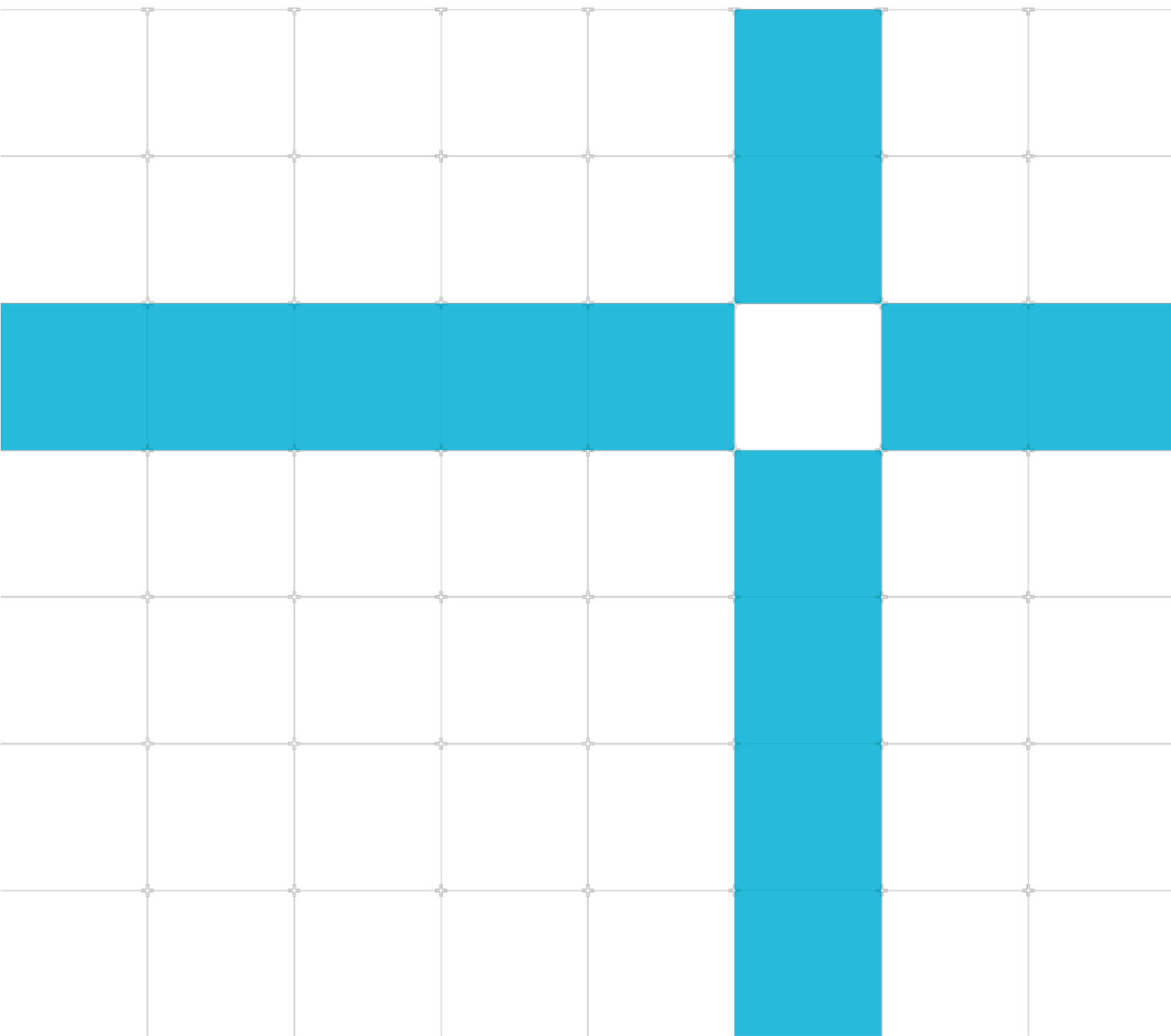
Software Optimization Guide

Non-Confidential

Copyright © 2020 Arm Limited (or its affiliates).
All rights reserved.

Issue 1.0

PJDOC-466751330-14664



Arm® Cortex®-A78C Core Software Optimization Guide

Copyright © 2020 Arm Limited (or its affiliates). All rights reserved.

Release information

Document history

Issue	Date	Confidentiality	Change
1.0	02 Nov 2020	Non-Confidential	First release for r0p1

Non-Confidential Proprietary Notice

This document is protected by copyright and other related rights and the practice or implementation of the information contained in this document may be protected by one or more patents or pending patent applications. No part of this document may be reproduced in any form by any means without the express prior written permission of Arm. No license, express or implied, by estoppel or otherwise to any intellectual property rights is granted by this document unless specifically stated.

Your access to the information in this document is conditional upon your acceptance that you will not use or permit others to use the information for the purposes of determining whether implementations infringe any third party patents.

THIS DOCUMENT IS PROVIDED "AS IS". ARM PROVIDES NO REPRESENTATIONS AND NO WARRANTIES, EXPRESS, IMPLIED OR STATUTORY, INCLUDING, WITHOUT LIMITATION, THE IMPLIED WARRANTIES OF MERCHANTABILITY, SATISFACTORY QUALITY, NON-INFRINGEMENT OR FITNESS FOR A PARTICULAR PURPOSE WITH RESPECT TO THE DOCUMENT. For the avoidance of doubt, Arm makes no representation with respect to, and has undertaken no analysis to identify or understand the scope and content of, patents, copyrights, trade secrets, or other rights.

This document may include technical inaccuracies or typographical errors.

TO THE EXTENT NOT PROHIBITED BY LAW, IN NO EVENT WILL ARM BE LIABLE FOR ANY DAMAGES, INCLUDING WITHOUT LIMITATION ANY DIRECT, INDIRECT, SPECIAL, INCIDENTAL, PUNITIVE, OR CONSEQUENTIAL DAMAGES, HOWEVER CAUSED AND REGARDLESS OF THE THEORY OF LIABILITY, ARISING OUT OF ANY USE OF THIS DOCUMENT, EVEN IF ARM HAS BEEN ADVISED OF THE POSSIBILITY OF SUCH DAMAGES.

This document consists solely of commercial items. You shall be responsible for ensuring that any use, duplication or disclosure of this document complies fully with any relevant export laws and regulations to assure that this document or any portion thereof is not exported, directly or indirectly, in violation of such export laws. Use of the word "partner" in reference to Arm's customers is not intended to create or refer to any partnership relationship with any other company. Arm may make changes to this document at any time and without notice.

If any of the provisions contained in these terms conflict with any of the provisions of any click through or signed written agreement covering this document with Arm, then the click through or signed written agreement prevails over and supersedes the conflicting provisions of these terms. This document may be translated into other languages for convenience, and you agree that if there is any conflict between the English version of this document and any translation, the terms of the English version of the Agreement shall prevail.

The Arm corporate logo and words marked with ® or ™ are registered trademarks or trademarks of Arm Limited (or its affiliates) in the US and/or elsewhere. All rights reserved. Other brands and names mentioned in this document may be the trademarks of their respective owners. Please follow Arm's trademark usage guidelines at <http://www.arm.com/company/policies/trademarks>.

Copyright © 2020 Arm Limited (or its affiliates). All rights reserved.

Arm Limited. Company 02557590 registered in England.

110 Fulbourn Road, Cambridge, England CB1 9NJ.

(LES-PRE-20349)

Confidentiality Status

This document is Non-Confidential. The right to use, copy and disclose this document may be subject to license restrictions in accordance with the terms of the agreement entered into by Arm and the party that Arm delivered this document to.

Unrestricted Access is an Arm internal classification.

Product Status

The information in this document is Final, that is for a developed product.

Web Address

www.arm.com

Contents

1 Introduction.....	6
1.1 Product revision status.....	6
1.2 Intended audience.....	6
1.3 Scope.....	6
1.4 Conventions	6
1.4.1 Glossary.....	6
1.4.2 Typographical conventions	8
1.5 Additional reading.....	9
1.6 Feedback.....	9
1.6.1 Feedback on this product.....	9
1.6.2 Feedback on content	9
2 Overview	10
2.1 Pipeline overview	11
3 Instruction characteristics	13
3.1 Instruction tables.....	13
3.2 Legend for reading the utilized pipelines	13
3.3 Branch instructions.....	14
3.4 Arithmetic and logical instructions	14
3.5 Move and shift instructions	16
3.6 Divide and multiply instructions	17
3.7 Saturating and parallel arithmetic instructions	18
3.8 Pointer Authentication instructions.....	20
3.9 Miscellaneous data-processing instructions.....	21
3.10 Load instructions.....	23
3.11 Store instructions	25
3.12 FP data processing instructions.....	27
3.13 FP miscellaneous instructions	29
3.14 FP load instructions.....	30
3.15 FP store instructions.....	32
3.16 ASIMD integer instructions.....	34

3.17 ASIMD floating-point instructions	37
3.18 ASIMD miscellaneous instructions.....	41
3.19 ASIMD load instructions	43
3.20 ASIMD store instructions	46
3.21 Cryptography extensions.....	49
3.22 CRC	50
4 Special considerations.....	51
4.1 Dispatch constraints.....	51
4.2 Dispatch stall.....	51
4.3 Optimizing general-purpose register spills and fills	51
4.4 Optimizing memory routines.....	51
4.5 Load/Store alignment.....	53
4.6 AES encryption/decryption.....	53
4.7 Region based fast forwarding.....	54
4.8 Branch instruction alignment.....	55
4.9 FPCR self-synchronization.....	55
4.10 Special register access.....	55
4.11 Register forwarding hazards.....	57
4.12 IT blocks	57
4.13 Instruction fusion.....	57
4.14 Mixing Arm and Thumb code	58
4.15 Cache maintenance operations	58

1 Introduction

1.1 Product revision status

The *rm**pn* identifier indicates the revision status of the product described in this book, for example, *r1p2*, where:

rm

Identifies the major revision of the product, for example, *r1*.

pn

Identifies the minor revision or modification status of the product, for example, *p2*.

1.2 Intended audience

This document is for system designers, system integrators, and programmers who are designing or programming a System-on-Chip (SoC) that uses an Arm core.

1.3 Scope

This document describes aspects of the Cortex®-A78C core micro-architecture that influence software performance. Micro-architectural detail is limited to that which is useful for software optimization.

Documentation extends only to software visible behavior of the Cortex-A78C core and not to the hardware rationale behind the behavior.

1.4 Conventions

The following subsections describe conventions used in Arm documents.

1.4.1 Glossary

The Arm Glossary is a list of terms used in Arm documentation, together with definitions for those terms. The Arm Glossary does not contain terms that are industry standard unless the Arm meaning differs from the generally accepted meaning.







See the Arm Glossary for more information: <https://developer.arm.com/glossary>.

1.4.1.1 Terms and Abbreviations

This document uses the following terms and abbreviations.

Term	Meaning
ALU	Arithmetic and Logical Unit
ASIMD	Advanced SIMD
MOP	Macro-Operation
μOP	Micro-Operation
SQRT	Square Root
T32	AArch32 Thumb® instruction set
FP	Floating-point

1.4.2 Typographical conventions

Convention	Use
<i>italic</i>	Introduces citations.
bold	Highlights interface elements, such as menu names. Denotes signal names. Also used for terms in descriptive lists, where appropriate.
<code>monospace</code>	Denotes text that you can enter at the keyboard, such as commands, file and program names, and source code.
<code>monospace bold</code>	Denotes language keywords when used outside example code.
<code>monospace <u>underline</u></code>	Denotes a permitted abbreviation for a command or option. You can enter the underlined text instead of the full command or option name.
<and>	Encloses replaceable terms for assembler syntax where they appear in code or code fragments. For example: <code>MRC p15, 0, <Rd>, <CRn>, <CRm>, <Opcode_2></code>
SMALL CAPITALS	Used in body text for a few terms that have specific technical meanings, that are defined in the Arm® Glossary. For example, IMPLEMENTATION DEFINED, IMPLEMENTATION SPECIFIC, UNKNOWN, and UNPREDICTABLE.
 Caution	This represents a recommendation which, if not followed, might lead to system failure or damage.
 Warning	This represents a requirement for the system that, if not followed, might result in system failure or damage.
 Danger	This represents a requirement for the system that, if not followed, will result in system failure or damage.
 Note	This represents an important piece of information that needs your attention.
 Tip	This represents a useful tip that might make it easier, better or faster to perform a task.
 Remember	This is a reminder of something important that relates to the information you are reading.

1.5 Additional reading

This document contains information that is specific to this product. See the following documents for other relevant information:

Table 1-1 Arm publications

Document name	Document ID	Licensee only
Arm® Architecture Reference Manual, Armv8, for Armv8-A architecture profile	DDI 0487	No
Arm® Cortex®-A78C Core Technical Reference Manual	101971	No

1.6 Feedback

Arm welcomes feedback on this product and its documentation.

1.6.1 Feedback on this product

If you have any comments or suggestions about this product, contact your supplier and give:

- The product name.
- The product revision or version.
- An explanation with as much information as you can provide. Include symptoms and diagnostic procedures if appropriate.

1.6.2 Feedback on content

If you have comments on content, send an email to errata@arm.com and give:

- The title Arm® Cortex®-A78C Core Software Optimization Guide.
- The number PJDOC-466751330-14664.
- If applicable, the page number(s) to which your comments refer.
- A concise explanation of your comments.

Arm also welcomes general suggestions for additions and improvements.



Arm tests the PDF only in Adobe Acrobat and Acrobat Reader and cannot guarantee the quality of the represented document when used with any other PDF reader.

2 Overview

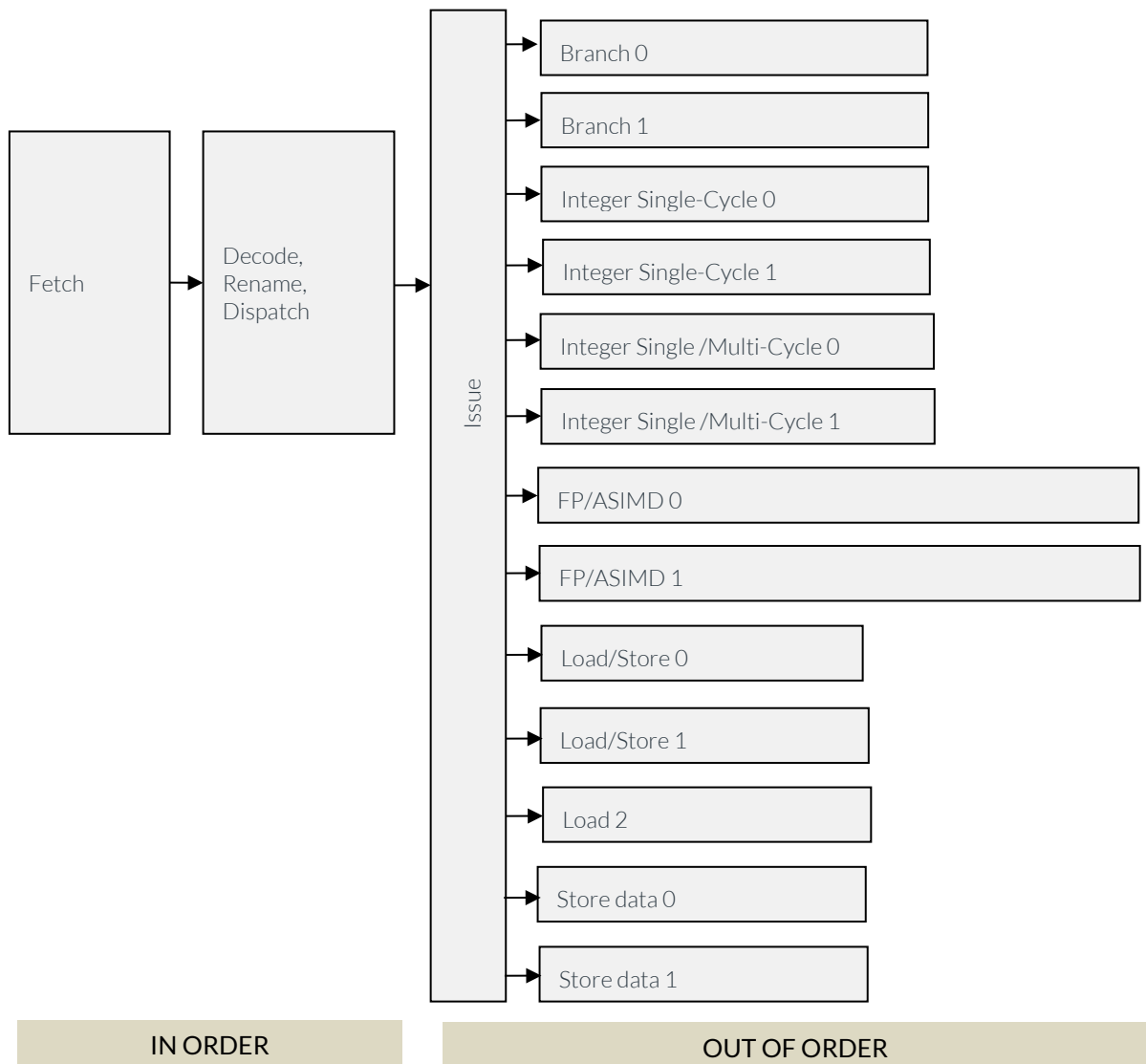
The Cortex-A78C core is a high-performance, low-power core that implements the Armv8-A architecture with support for the Armv8.1-A extension, Armv8.2-A extension, including the RAS extension, the Load acquire (LDAPR) instructions and Pointer Authentication introduced in the Armv8.3-A extension, the Dot Product and the multi-OS support introduced in the Armv8.4-A extension, and the Enhanced Pointer Authentication (excluding the optional FPAC extension) introduced in the Armv8.6-A extension.

This document describes elements of the Cortex-A78C core micro-architecture that influence software performance so that software and compilers can be optimized accordingly.

2.1 Pipeline overview

The following figure describes the high-level Cortex-A78C instruction processing pipeline. Instructions are first fetched and then decoded into internal Macro-Operations (MOPs). From there, the MOPs proceed through register renaming and dispatch stages. A MOP can be split into two Micro-Operations (μOPs) further down the pipeline after the decode stage. Once dispatched, μOPs wait for their operands and issue out-of-order to one of thirteen issue pipelines. Each issue pipeline can accept one μOP per cycle.

Figure 2-1 Cortex-A78C core pipeline



The execution pipelines support different types of operations, as follows:

Table 2-1 Cortex-A78C core operations

Instruction groups	Instructions
Branch 0/1	Branch μ Ops
Integer Single-Cycle 0/1	Integer ALU μ Ops
Integer Single/Multi-cycle 0/1	Integer shift-ALU, multiply, divide, CRC and sum-of-absolute-differences μ Ops
Load/Store 0/1	Load, Store address generation and special memory μ Ops
Load 2	Load μ Ops
Store data 0/1	Integer store data μ Ops
FP/ASIMD-0	ASIMD ALU, ASIMD misc, ASIMD integer multiply, FP convert, FP misc, FP add, FP multiply, FP divide, FP sqrt, crypto μ Ops, Vector store data μ Ops
FP/ASIMD-1	ASIMD ALU, ASIMD misc, FP misc, FP add, FP multiply, ASIMD shift μ Ops, Vector store data μ Ops, crypto μ Ops.

3 Instruction characteristics

3.1 Instruction tables

This chapter describes high-level performance characteristics for most Armv8.2-A A32, T32, and A64 instructions. A series of tables summarize the effective execution latency and throughput (instruction bandwidth per cycle), pipelines utilized, and special behaviours associated with each group of instructions. Utilized pipelines correspond to the execution pipelines described in chapter 2.

In the tables below, Execution Latency is defined as the minimum latency seen by an operation dependent on an instruction in the described group.

In the tables below, Execution Throughput is defined as the maximum throughput (in instructions per cycle) of the specified instruction group that can be achieved in the entirety of the Cortex-A78C microarchitecture.

3.2 Legend for reading the utilized pipelines

Table 3-1 Cortex-A78C core pipeline names and symbols

Pipeline name	Symbol used in tables
Branch 0/1	B
Integer single Cycle 0/1	S
Integer single Cycle 0/1 and single/multicycle 0/1	I
Integer single/multicycle 0/1	M
Integer multicycle 0	M0
Load/Store 0/1 and Load 2	L
Store data 0/1	D
FP/ASIMD 0/1	V
FP/ASIMD 0	V0
FP/ASIMD 1	V1

3.3 Branch instructions

Table 3-2 AArch64 Branch instructions

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Branch, immed	B	1	2	B	-
Branch, register	BR, RET	1	2	B	-
Branch and link, immed	BL	1	2	B	-
Branch and link, register	BLR	1	2	B	-
Compare and branch	CBZ, CBNZ, TBZ, TBNZ	1	2	B	-

Table 3-3 AArch32 Branch instructions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Branch, immed	B	1	2	B	-
Branch, register	BX	1	2	B	-
Branch and link, immed	BL, BLX	1	2	B	-
Branch and link, register	BLX	1	2	B	-
Compare and branch	CBZ, CBNZ	1	2	B	-

3.4 Arithmetic and logical instructions

Table 3-4 AArch64 Arithmetic and logical instructions

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ALU, basic	ADD, ADC, AND, BIC, EON, EOR, ORN, ORR, SUB, SBC	1	4	I	-
ALU, basic, flagset	ADDS, ADCS, ANDS, BICS, SUBS, SBCS	1	3	I	-
ALU, extend and shift	ADD{S}, AND{S}, BIC{S}, EON, EOR, ORN, ORR, SUB{S}	2	2	M	-
Arithmetic, LSL shift, shift <= 4	ADD, SUB	1	4	I	-
Arithmetic, flagset, LSL shift, shift <= 4	ADDS, SUBS	1	3	I	-
Arithmetic, LSR/ASR/ROR shift or LSL shift > 4	ADD{S}, SUB{S}	2	2	M	-

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Conditional compare	CCMN, CCMMP	1	3	I	-
Conditional select	CSEL, CSINC, CSINV, CSNEG	1	4	I	-
Logical, shift, no flagset	AND, BIC, EON, EOR, ORN, ORR	1	4	I	-
Logical, shift, flagset	AND{S}, BIC{S}	2	2	M	-
Flag manipulation instructions	SETF8, SETF16, RMIF, CFINV	1	3	I	-

Table 3-5 AArch32 Arithmetic and logical instructions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ALU, basic, unconditional, no flagset	ADD, ADC, ADR, AND, BIC, EOR, ORN, ORR, RSB, RSC, SUB, SBC	1	4	I	-
ALU, basic, unconditional, flagset	ADDS, ADCS, ANDS, BICS, CMN, CMP, EORS, ORNS, ORRS, RSBS, RSCS, SUBS, SBBS, TEQ, TST	1	3	I	-
ALU, basic, conditional	ADD{S}, ADC{S}, AND{S}, BIC{S}, CMN, CMP, EOR{S}, ORN{S}, ORR{S}, RSB{S}, RSC{S}, SUB{S}, SBC{S}, TEQ, TST	1	1	M0	-
ALU, basic, shift by register, conditional	(same as ALU basic, conditional)	2	1	I, M0	-
ALU, basic, shift by register, unconditional, flagset	(same as ALU, basic, unconditional, flagset)	2	1	M0	-
Arithmetic, shift by register, unconditional, no flagset	ADD, ADC, RSB, RSC, SUB, SBC	2	1	M0	-
Logical, shift by register, unconditional, no flagset	AND, BIC, EOR, ORN, ORR	1	1	M0	-
Arithmetic, LSL shift by immed, shift <= 4, unconditional, no flagset	ADD, ADC, RSB, RSC, SUB, SBC	1	4	I	-
Arithmetic, LSL shift by immed, shift <= 4, unconditional, flagset	ADDS, ADCS, RSBS, RSCS, SUBS, SBBS	1	3	I	-
Arithmetic, LSL shift by immed, shift <= 4, conditional	ADD{S}, ADC{S}, RSB{S}, RSC{S}, SUB{S}, SBC{S}	1	1	M0	-

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Arithmetic, LSR/ASR/ROR shift by immed or LSL shift by immed > 4, unconditional	ADD{S}, ADC{S}, RSB{S}, RSC{S}, SUB{S}, SBC{S}	2	2	M	-
Arithmetic, LSR/ASR/ROR shift by immed or LSL shift by immed > 4, conditional	ADD{S}, ADC{S}, RSB{S}, RSC{S}, SUB{S}, SBC{S}	2	1	M0	-
Logical, shift by immed, no flagset, unconditional	AND, BIC, EOR, ORN, ORR	1	4	I	-
Logical, shift by immed, no flagset, conditional	AND, BIC, EOR, ORN, ORR	1	1	M0	-
Logical, shift by immed, flagset, unconditional	ANDS, BICS, EORS, ORNS, ORRS	2	2	M	-
Logical, shift by immed, flagset, conditional	ANDS, BICS, EORS, ORNS, ORRS	2	1	M0	-
Test/Compare, shift by immed	CMN, CMP, TEQ, TST	2	2	M	-
Branch forms	-	+1	2	+B	1

Notes:

1. Branch forms are possible when the instruction destination register is the PC. For those cases, an additional branch μ OP is required. This adds 1 cycle to the latency.

3.5 Move and shift instructions

Table 3-6 AArch32 Move and shift instructions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Move, basic	MOV{S}, MOVW, MVN{S}	1	4	I	-
Move, shift by immed, no flagset	ASR, LSL, LSR, ROR, RRX, MVN	1	4	I	-
Move, shift by immed, flagset	ASRS, LSLS, LSRS, RORS, RRXS, MVNS	2	2	M	-
Move, shift by register, no flagset, unconditional	ASR, LSL, LSR, ROR, RRX, MVN	1	4	I	-
Move, shift by register, no flagset, conditional	ASR, LSL, LSR, ROR, RRX, MVN	2	2	I	-
Move, shift by register, flagset	ASRS, LSLS, LSRS, RORS, RRXS, MVNS	2	1	M0	-
Move, top	MOVT	1	4	I	-
Move, branch forms	-	+1	2	+B	-

3.6 Divide and multiply instructions

Table 3-7 AArch64 Divide and multiply instructions

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Divide, W-form	SDIV, UDIV	5 to 12	1/12 to 1/5	M0	1
Divide, X-form	SDIV, UDIV	5 to 20	1/20 to 1/5	M0	1
Multiply accumulate, W-form	MADD, MSUB	2(1)	1	M0	2
Multiply accumulate, X-form	MADD, MSUB	2(1)	1	M0	2
Multiply accumulate long	SMADDL, SMSUBL, UMADDL, UMSUBL	2(1)	2	M	2
Multiply high	SMULH, UMULH	3	2	M	2

Table 3-8 AArch32 Divide and multiply instructions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Divide	SDIV, UDIV	5 to 12	1/12 to 1/5	M0	1
Multiply, unconditional	MUL, SMULBB, SMULBT, SMULTB, SMULTT, SMULWB, SMULWT, SMMUL{R}, SMUAD{X}, SMUSD{X}	2	2	M	-
Multiply, conditional	MUL, SMULBB, SMULBT, SMULTB, SMULTT, SMULWB, SMULWT, SMMUL{R}, SMUAD{X}, SMUSD{X}	2	1	M0	-
Multiply accumulate, conditional	MLA, MLS, SMLABB, SMLABT, SMLATB, SMLATT, SMLAWB, SMLAWT, SMLAD{X}, SMLSD{X}, SMMLA{R}, SMMLS{R}	3	1	M0, I	-
Multiply accumulate, unconditional	MLA, MLS, SMLABB, SMLABT, SMLATB, SMLATT, SMLAWB, SMLAWT, SMLAD{X}, SMLSD{X}, SMMLA{R}, SMMLS{R}	2(1)	1	M0	2

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Multiply accumulate accumulate long, conditional	UMAAL	4	1	I, M0	-
Multiply accumulate accumulate long, unconditional	UMAAL	3	1	I, M0	-
Multiply accumulate long, no flagset	SMLAL, SMLALBB, SMLALBT, SMLALTB, SMLALTT, SMLALD{X}, SMLSLD{X}, UMLAL	3	1	M0, I	-
Multiply accumulate long, flagset	SMLAL, SMLALBB, SMLALBT, SMLALTB, SMLALTT, SMLALD{X}, SMLSLD{X}, UMLAL	4	1	M0, I	-
Multiply long, unconditional, no flagset	SMULL, UMULL	2	2	M	-
Multiply long, unconditional, flagset	SMULLS, UMULLS	3	1	M, I	-
Multiply long, conditional,	SMULL{S}, UMULL{S}	3	1	M, I	-

Notes:

1. Integer divides are performed using an iterative algorithm and block any subsequent divide operations until complete. Early termination is possible, depending upon the data values.
2. Multiply-accumulate pipelines support late-forwarding of accumulate operands from similar μ OPs, allowing a typical sequence of multiply-accumulate μ OPs to issue one every N cycles (accumulate latency N shown in parentheses). Accumulator forwarding is not supported for consumers of 64 bit multiply high operations.

3.7 Saturating and parallel arithmetic instructions

Table 3-9 AArch32 Saturating and parallel arithmetic instructions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Parallel arith, unconditional	SADD16, SADD8, SSUB16, SSUB8, UADD16, UADD8, USUB16, USUB8	2	1	M	-
Parallel arith, conditional	SADD16, SADD8, SSUB16, SSUB8, UADD16, UADD8, USUB16, USUB8	2(4)	1	M0, I	1
Parallel arith with exchange, unconditional	SASX, SSAX, UASX, USAX	3	2	I, M	-

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Parallel arith with exchange, conditional	SASX, SSAX, UASX, USAX	3(5)	1	I, M0	1
Parallel halving arith, unconditional	SHADD16, SHADD8, SHSUB16, SHSUB8, UHADD16, UHADD8, UHSUB16, UHSUB8	2	2	M	-
Parallel halving arith, conditional	SHADD16, SHADD8, SHSUB16, SHSUB8, UHADD16, UHADD8, UHSUB16, UHSUB8	2	1	M0	-
Parallel halving arith with exchange	SHASX, SHSAX, UHASX, UHSAX	3	1	I, M0	-
Parallel saturating arith, unconditional	QADD16, QADD8, QSUB16, QSUB8, UQADD16, UQADD8, UQSUB16, UQSUB8	2	2	M	-
Parallel saturating arith, conditional	QADD16, QADD8, QSUB16, QSUB8, UQADD16, UQADD8, UQSUB16, UQSUB8	2	1	M0	-
Parallel saturating arith with exchange, unconditional	QASX, QSAX, UQASX, UQSAX	3	2	I, M	-
Parallel saturating arith with exchange, conditional	QASX, QSAX, UQASX, UQSAX	3	1	I, M0	-
Saturate, unconditional	SSAT, SSAT16, USAT, USAT16	2	2	M	-
Saturate, conditional	SSAT, SSAT16, USAT, USAT16	2	1	M0	-
Saturating arith, unconditional	QADD, QSUB	2	2	M	-
Saturating arith, conditional	QADD, QSUB	2	1	M0	-
Saturating doubling arith, unconditional	QDADD, QDSUB	3	1	M, M	-
Saturating doubling arith conditional	QDADD, QDSUB	3	1	M, M0	-

Notes:

1. Conditional GE-setting instructions require three extra μ OPs and two additional cycles to conditionally update the GE field (GE latency shown in parentheses).

3.8 Pointer Authentication instructions

Table 3-10 AArch64 pointer authentication instructions

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Authenticate data address	AUTDA, AUTDB, AUTDZA, AUTDZB	5	1	M0	-
Authenticate instruction address	AUTIA, AUTIB, AUTIA1716, AUTIB1716, AUTIASP, AUTIBSP, AUTIAZ, AUTIBZ, AUTIZA, AUTIZB	5	1	M0	-
Branch and link, register, with pointer authentication	BLRAA, BLRAAZ, BLRAB, BLRABZ	6	1	M0, B	
Branch, register, with pointer authentication	BRAA, BRAAZ, BRAB, BRABZ	6	1	M0, B	
Branch, return, with pointer authentication	RETA, RETB	6	1	M0, B	
Compute pointer authentication code for data address	PACDA, PACDB, PACDZA, PACDZB	5	1	M0	
Compute pointer authentication code, using generic key	PACGA	5	1	M0	
Compute pointer authentication code for instruction address	PACIA, PACIB, PACIA1716, PACIB1716, PACIASP, PACIBSP, PACIAZ, PACIBZ, PACIZA, PACIZB	5	1	M0	
Load register, with pointer authentication	LDRAA, LDRAB	9	1	M0, L	
Strip pointer authentication code	XPACD, XPACI, XPACLRI	2	1	M0	

3.9 Miscellaneous data-processing instructions

Table 3-11 AArch64 Miscellaneous data-processing instructions

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Address generation	ADR, ADRP	1	4	I	-
Bitfield extract, one reg	EXTR	1	4	I	1
Bitfield extract, two regs	EXTR	3	2	I, M	-
Bitfield move, basic	SBFM, UBFM	1	4	I	-
Bitfield move, insert	BFM	2	2	M	-
Count leading	CLS, CLZ	1	4	I	-
Move immed	MOVN, MOVK, MOVZ	1	4	I	-
Reverse bits/bytes	RBIT, REV, REV16, REV32	1	4	I	-
Variable shift	ASRV, LSLV, LSRV, RORV	1	4	I	-

Notes:

1. One reg form is when Rn==Rm or imm==0, all other forms are considered two regs.

Table 3-12 AArch32 Miscellaneous data-processing instructions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Bit field extract	SBFX, UBFX	1	4	I	-
Bit field insert/clear, unconditional	BFI, BFC	2	2	M	-
Bit field insert/clear, conditional	BFI, BFC	2	1	M0	-
Count leading zeros	CLZ	1	4	I	-
Pack halfword, unconditional	PKH	2	2	M	-
Pack halfword, conditional	PKH	2	1	M0	-
Reverse bits/bytes	RBIT, REV, REV16, REVSH	1	4	I	-
Select bytes, unconditional	SEL	1	4	I	-
Select bytes, conditional	SEL	2	2	I	-
Sign/zero extend, normal	SXTB, SXTH, UXTB, UXTH	1	4	I	-
Sign/zero extend, parallel, unconditional	SXTB16, UXTB16	2	2	M	-
Sign/zero extend, parallel, conditional	SXTB16, UXTB16	2	1	M0	-
Sign/zero extend and add, normal, unconditional	SXTAB, SXTAH, UXTAB, UXTAH	2	2	M	-
Sign/zero extend and add, normal, conditional	SXTAB, SXTAH, UXTAB, UXTAH	2	1	M0	-
Sign/zero extend and add, parallel, unconditional	SXTAB16, UXTAB16	4	1	M	-
Sign/zero extend and add, parallel, conditional	SXTAB16, UXTAB16	4	1	M, M0	-
Sum of absolute differences	USAD8	2	1	M0	-
Sum of absolute differences accumulate, unconditional	USADA8	2	1	M0	-
Sum of absolute differences accumulate, conditional	USADA8	3	1	M0, I	-

3.10 Load instructions

The latencies shown assume the memory access hits in the Level 1 Data Cache.

Table 3-13 AArch64 Load instructions

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Load register, literal	LDR, LDRSW, PRFM	4	3	L	-
Load register, unscaled immed	LDUR, LDURB, LDURH, LDURSB, LDURSH, LDURSW, PRFUM	4	3	L	-
Load register, immed post-index	LDR, LDRB, LDRH, LDRSB, LDRSH, LDRSW	4	3	L, I	-
Load register, immed pre-index	LDR, LDRB, LDRH, LDRSB, LDRSH, LDRSW	4	3	L, I	-
Load register, immed unprivileged	LDTR, LDTRB, LDTRH, LDTRSB, LDTRSH, LDTRSW	4	3	L	-
Load register, unsigned immed	LDR, LDRB, LDRH, LDRSB, LDRSH, LDRSW, PRFM	4	3	L	-
Load register, register offset, basic	LDR, LDRB, LDRH, LDRSB, LDRSH, LDRSW, PRFM	4	3	L	-
Load register, register offset, scale by 4/8	LDR, LDRSW, PRFM	4	3	L	-
Load register, register offset, scale by 2	LDRH, LDRSH	5	3	I, L	-
Load register, register offset, extend	LDR, LDRB, LDRH, LDRSB, LDRSH, LDRSW, PRFM	4	3	L	-
Load register, register offset, extend, scale by 4/8	LDR, LDRSW, PRFM	4	3	L	-
Load register, register offset, extend, scale by 2	LDRH, LDRSH	5	3	I, L	-
Load pair, signed immed offset, normal, W-form	LDP, LDNP	4	3	L	-
Load pair, signed immed offset, normal, X-form	LDP, LDNP	4	3/2	L	-
Load pair, signed immed offset, signed words	LDPSW	5	3/2	I, L	-
Load pair, immed post-index or immed pre-index, normal, W-form	LDP	4	3	L, I	-

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Load pair, immed post-index or immed pre-index, normal, X-form	LDP	4	3/2	L, I	-
Load pair, immed post-index, signed words	LDPSW	5	3/2	I, L	-

Table 3-14 AArch32 Load instructions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Load, immed offset	LDR{T}, LDRB{T}, LDRD, LDRH{T}, LDRSB{T}, LDRSH{T}	4	3	L	1, 2
Load, register offset, plus	LDR, LDRB, LDRD, LDRH, LDRSB, LDRSH	4	3	L	1, 2
Load, register offset, minus	LDR, LDRB, LDRD, LDRH, LDRSB, LDRSH	5	3	I, L	1, 2
Load, scaled register offset, plus, LSL2	LDR, LDRB	4	3	L	1, 2
Load, scaled register offset, other	LDR, LDRB, LDRH, LDRSB, LDRSH	5	3	I, L	1, 2
Load, immed pre-indexed	LDR, LDRB, LDRD, LDRH, LDRSB, LDRSH	4	3	L, I	1, 2
Load, register pre-indexed	LDRH, LDRSB, LDRSH	5	3	I, L, M0	1, 2, 3
Load, register pre-indexed	LDRD	4	3	L, M0	1, 2, 3
Load, scaled register pre-indexed, plus, LSL2	LDR, LDRB	4	3	L, M0	1, 2, 3
Load, scaled register pre-indexed, unshifted	LDR, LDRB	4	3	L, M0	1, 2, 3
Load, scaled register pre-indexed, other	LDR, LDRB	5	3	I, L, M0	1, 2, 3
Load, immed post-indexed	LDR{T}, LDRB{T}, LDRD, LDRH{T}, LDRSB{T}, LDRSH{T}	4	3	L, I	1, 2
Load, register post-indexed	LDR{T}, LDRB{T}, LDRH{T}, LDRSB{T}, LDRSH{T}	5	3	I, L, M0	1, 2, 3
Load, register post-indexed	LDRD	4	3	L, M0	1, 2, 3
Preload, immed offset	PLD, PLDW	4	3	L	-

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Preload, register offset, plus, LSL2 and unshifted	PLD, PLDW	4	3	L	-
Preload, register offset, minus	PLD, PLDW	5	3	I, L	-
Load multiple, no writeback, base reg not in list	LDMIA, LDMIB, LDMDA, LDMDB	N	3/R	L	1, 4, 5
Load multiple, no writeback, base reg in list	LDMIA, LDMIB, LDMDA, LDMDB	1+ N	3/R	I, L	1, 4, 5
Load multiple, writeback	LDMIA, LDMIB, LDMDA, LDMDB, POP	1+ N	3/R	L, I	1, 4, 5
(Load, all branch forms)	-	+1	-	+ B	6

Notes:

1. Conditional loads have extra μOP (s) which goes down pipeline 'I' and have 1 cycle extra latency compared to their unconditional counterparts.
2. Conditional loads go down L01 pipe and have an execution throughput of 2, whereas unconditional versions have a throughput of 3.
3. The address update op goes down pipeline 'I' if the load is unconditional.
4. N is floor $[(\text{num_reg}+5)/6]$.
5. R is floor $[(\text{num_reg}+1)/2]$.
6. Branch forms are possible when the instruction destination register is the PC. For those cases, an additional branch μOP is required. This adds 1 cycle to the latency.

3.11 Store instructions

The following table describes performance characteristics for standard store instructions. Stores μOP s are split into address and data μOP s. Once executed, stores are buffered and committed in the background.

Table 3-15 AArch64 Store instructions

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Store register, unscaled immed	STUR, STURB, STURH	1	2	L01, D	-
Store register, immed post-index	STR, STRB, STRH	1	2	L01, D	-
Store register, immed pre-index	STR, STRB, STRH	1	2	L01, D	-
Store register, immed unprivileged	STTR, STTRB, STTRH	1	2	L01, D	-
Store register, unsigned immed	STR, STRB, STRH	1	2	L01, D	-
Store register, register offset, basic	STR, STRB, STRH	1	2	L01, D	-
Store register, register offset, scaled by 4/8	STR	1	2	L01, D	-

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Store register, register offset, scaled by 2	STRH	2	2	I, L01, D	-
Store register, register offset, extend	STR, STRB, STRH	1	2	L01, D	-
Store register, register offset, extend, scale by 4/8	STR	1	2	L01, D	-
Store register, register offset, extend, scale by 1	STRH	2	2	I, L01, D	-
Store pair, immed offset	STP, STNP	1	2	L01, D	-
Store pair, immed post-index	STP	1	2	L01, D, I	-
Store pair, immed pre-index	STP	1	2	L01, D, I	-

Table 3-16 AArch32 Store instructions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Store, immed offset	STR{T}, STRB{T}, STRD, STRH{T}	1	2	L01, D	-
Store, register offset, plus	STR, STRB, STRD, STRH	1	2	L01, D	-
Store, register offset, minus	STR, STRB, STRD, STRH	1	2	L01, D	-
Store, scaled register offset, plus, no shift	STR, STRB	1	2	L01, D	-
Store, scaled register offset, plus, LSL2	STR, STRB	1	2	L01, D	-
Store, scaled register offset, plus, other	STR, STRB	2	2	I, L01, D	-
Store, scaled register offset, minus	STR, STRB	2	2	I, L01, D	-
Store, immed pre-indexed	STR, STRB, STRD, STRH	1	2	L01, D, I	-
Store, register pre-indexed, plus, no shift	STR, STRB, STRD, STRH	1	2	L01, D, MO	1
Store, register pre-indexed, minus	STR, STRB, STRD, STRH	2	2	I, L01, D, MO	1
Store, scaled register pre-indexed, plus LSL2	STR, STRB	1	2	L01, D, MO	1
Store, scaled register pre-indexed, other	STR, STRB	2	2	I, L01, D, MO	1
Store, immed post-indexed	STR{T}, STRB{T}, STRD, STRH{T}	1	2	L01, D, I	-
Store, register post-indexed	STRH{T}, STRD	1	2	L01, D, MO	1

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Store, register post-indexed	STR{T}, STRB{T}	1	2	L01, D, M0	1
Store, scaled register post-indexed	STR{T}, STRB{T}	1	2	L01, D, M0	2
Store multiple, no writeback	STMIA, STMIB, STMDA, STMDB	N	1/N	L01, D	3
Store multiple, writeback	STMIA, STMIB, STMDA, STMDB, PUSH	N	1/N	L01, D	3

Notes:

1. The address update op goes down pipeline 'I' if the store is unconditional.
2. The address update op goes down pipeline 'M' if the store is unconditional.
3. For store multiple instructions, $N = \text{floor}((\text{num_regs} + 3) / 4)$.

3.12 FP data processing instructions

Table 3-17 AArch64 FP data processing instructions

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
FP absolute value	FABS	2	2	V	-
FP arithmetic	FADD, FSUB	2	2	V	-
FP compare	FCCMP{E}, FCMP{E}	2	1/2	V0	-
FP divide, H-form	FDIV	7	4/7	V0	1
FP divide, S-form	FDIV	7 to 10	4/9 to 4/7	V0	1
FP divide, D-form	FDIV	7 to 15	1/7 to 2/7	V0	1
FP min/max	FMIN, FMINNM, FMAX, FMAXNM	2	2	V	-
FP multiply	FMUL, FNMUL	3	2	V	2
FP multiply accumulate	FMADD, FMSUB, FNMADD, FNMSUB	4 (2)	2	V	3
FP negate	FNEG	2	2	V	-
FP round to integral	FRINTA, FRINTI, FRINTM, FRINTN, FRINTP, FRINTX, FRINTZ	2	1	V0	-
FP select	FCSEL	2	2	V	-
FP square root, H-form	FSQRT	7	4/7	V0	1
FP square root, S-form	FSQRT	7 to 9	1/2 to 4/7	V0	1
FP square root, D-form	FSQRT	7 to 16	2/15 to 2/7	V0	1

Table 3-18 AArch32 FP data processing instructions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
VFP absolute value	VABS	2	2	V	-
VFP arith	VADD, VSUB	2	2	V	-
VFP compare, unconditional	VCMP, VCMPE	2	1	V0	-
VFP compare, conditional	VCMP, VCMPE	4	1	V, V0	-
VFP convert	VCVT{R}, VCVTB, VCVTT, VCVTA, VCVTM, VCVTN, VCVTP	2	1	V0	-
VFP divide, H-form	VDIV	7	4/7	V0	1
VFP divide, S-form	VDIV	7 to 10	4/9 to 4/7	V0	1
VFP divide, D-form	VDIV	7 to 15	1/7 to 2/7	V0	1
VFP max/min	VMAXNM, VMINNM	2	2	V	-
VFP multiply	VMUL, VNMUL	3	2	V	2
VFP multiply accumulate (chained)	VMLA, VMLS, VNMLA, VNMLS	5 (2)	2	V	3
VFP multiply accumulate (fused)	VFMA, VFMS, VFNMA, VFNMS	4 (2)	2	V	3
VFP negate	VNEG	2	2	V	-
VFP round to integral	VRINTA, VRINTM, VRINTN, VRINTP, VRINTR, VRINTX, VRINTZ	2	1	V0	-
VFP select	VSELEQ, VSELGE, VSELGT, VSELVS	2	2	V	-
VFP square root, H-form	VSQRT	7	4/7	V0	1
VFP square root, S-form	VSQRT	7 to 9	1/2 to 4/7	V0	1
VFP square root, D-form	VSQRT	7 to 16	2/15 to 2/7	V0	1

Notes:

1. FP divide and square root operations are performed using an iterative algorithm and block subsequent similar operations to the same pipeline until complete.
2. FP multiply-accumulate pipelines support late-forwarding of the result from FP multiply μ OPs to the accumulate operands of an FP multiply-accumulate μ OP. The latter can potentially be issued 1 cycle after the FP multiply μ OP has been issued.
3. FP multiply-accumulate pipelines support late-forwarding of accumulate operands from similar μ OPs, allowing a typical sequence of multiply-accumulate μ OPs to issue one every N cycles (accumulate latency N shown in parentheses).

3.13 FP miscellaneous instructions

Table 3-19 AArch64 FP miscellaneous instructions

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
FP convert, from vec to vec reg	FCVT, FCVTXN	2	1	V0	-
FP convert, from gen to vec reg	SCVTF, UCVTF	3	1	M0	-
FP convert, from vec to gen reg	FCVTAS, FCVTAU, FCVTMS, FCVTMU, FCVTNS, FCVTNU, FCVTPS, FCVTPU, FCVTZS, FCVTZU	3	1	V0	-
FP move, immed	FMOV	2	2	V	-
FP move, register	FMOV	2	2	V	-
FP transfer, from gen to vec reg	FMOV	3	1	M0	-
FP transfer, from vec to gen reg	FMOV	2	1	V1	-

Table 3-20 AArch32 FP miscellaneous instructions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
VFP move, immed	VMOV	2	2	V	-
VFP move, register	VMOV	2	2	V	-
VFP transfer, core to vfp, single reg to S-reg, cond	VMOV	5	1	M0, V	-
VFP transfer, core to vfp, single reg to S-reg, uncond	VMOV	3	1	M0	-
VFP transfer, core to vfp, single reg to upper/lower half of D-reg	VMOV	5	1	M0, V	-
VFP transfer, core to vfp, 2 regs to 2 S-reg, cond	VMOV	6	1/2	M0, V	-
VFP transfer, core to vfp, 2 regs to 2 S-reg, uncond	VMOV	4	1/2	M0	-
VFP transfer, core to vfp, 2 regs to D-reg, cond	VMOV	5	1	M0, V	-
VFP transfer, core to vfp, 2 regs to D-reg, uncond	VMOV	3	1	M0	-
VFP transfer, vfp S-reg or upper/lower half of vfp D-reg to core reg, cond	VMOV	3	1	V1, I	-
VFP transfer, vfp S-reg or upper/lower half of vfp D-reg to core reg, uncond	VMOV	2	1	V1	-

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
VFP transfer, vfp 2 S-reg or D-reg to 2 core regs, cond	VMOV	3	1	V1, I	-
VFP transfer, vfp 2 S-reg or D-reg to 2 core regs, uncond	VMOV	2	1	V1	-

3.14 FP load instructions

The latencies shown assume the memory access hits in the Level 1 Data Cache. Compared to standard loads, an extra cycle is required to forward results to FP/ASIMD pipelines.

Table 3-21 AArch64 FP load instructions

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Load vector reg, literal, S/D/Q forms	LDR	5	3	L	-
Load vector reg, unscaled immed	LDUR	5	3	L	-
Load vector reg, immed post-index	LDR	5	3	L, I	-
Load vector reg, immed pre-index	LDR	5	3	L, I	-
Load vector reg, unsigned immed	LDR	5	3	L	-
Load vector reg, register offset, basic	LDR	5	3	L, I	-
Load vector reg, register offset, scale, S/D-form	LDR	5	3	L, I	-
Load vector reg, register offset, scale, H/Q-form	LDR	6	3	I, L	-
Load vector reg, register offset, extend	LDR	5	3	L	-
Load vector reg, register offset, extend, scale, S/D-form	LDR	5	3	L	-
Load vector reg, register offset, extend, scale, H/Q-form	LDR	5	3	I, L	-
Load vector pair, immed offset, S/D-form	LDP, LDNP	5	3	L	-
Load vector pair, immed offset, Q-form	LDP, LDNP	5	3/2	L	-
Load vector pair, immed post-index, S/D-form	LDP	5	3	I, L	-

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Load vector pair, immed post-index, Q-form	LDP	5	3/2	L, I	-
Load vector pair, immed pre-index, S/D-form	LDP	5	3	I, L	-
Load vector pair, immed pre-index, Q-form	LDP	5	3/2	L, I	-

Table 3-22 AArch32 FP load instructions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
FP load, register	VLDR	5	3 (2)	L	1, 7
FP load multiple, S form	VLDmia, VLDMDB, VPOP	N (N*)	3/R (2/R)	L	1, 2, 3, 4, 6, 7
FP load multiple, D form	VLDmia, VLDMDB, VPOP	N (N*)	3/R (2/R)	L, V	1, 2, 3, 4, 6, 7
(FP load, writeback forms)	-	(1)	-	+ I	5

Notes:

1. Conditional loads have an extra uop which goes down pipeline 'V' and have 2 cycle extra latency compared to their unconditional counterparts.
2. N is $(\text{num_reg})/6 + 5$.
3. N* is $(\text{num_reg})/4 + 5$.
4. R is $\text{num_reg}/2$.
5. Writeback forms of load instructions require an extra μOP to update the base address. This update is typically performed in parallel with or prior to the load μOP (update latency shown in parentheses).
6. The number in parenthesis represents the latency and throughput of conditional loads.
7. Conditional loads go down L01 pipe.

3.15 FP store instructions

Stores MOPs are split into store address and store data μ OPs at dispatch time. Once executed, stores are buffered and committed in the background.

Table 3-23 AArch64 FP store instructions

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Store vector reg, unscaled immed, B/H/S/D-form	STUR	2	2	L01, V	-
Store vector reg, unscaled immed, Q-form	STUR	2	2	L01, V	-
Store vector reg, immed post-index, B/H/S/D-form	STR	2	2	L01, V	-
Store vector reg, immed post-index, Q-form	STR	2	2	L01, V	-
Store vector reg, immed pre-index, B/H/S/D-form	STR	2	2	L01, V	-
Store vector reg, immed pre-index, Q-form	STR	2	2	L01, V	-
Store vector reg, unsigned immed, B/H/S/D-form	STR	2	2	L01, V	-
Store vector reg, unsigned immed, Q-form	STR	2	2	L01, V	-
Store vector reg, register offset, basic, B/H/S/D-form	STR	2	2	L01, V	-
Store vector reg, register offset, basic, Q-form	STR	2	2	L01, V	-
Store vector reg, register offset, scale, H-form	STR	2	2	I, L01, V	-
Store vector reg, register offset, scale, S/D-form	STR	2	2	L01, V	-
Store vector reg, register offset, scale, Q-form	STR	2	2	L01, V	-
Store vector reg, register offset, extend, B/H/S/D-form	STR	2	2	L01, V	-
Store vector reg, register offset, extend, Q-form	STR	2	2	L01, V	-
Store vector reg, register offset, extend, scale, H-form	STR	2	2	I, L01, V	-
Store vector reg, register offset, extend, scale, S/D-form	STR	2	2	L01, V	-
Store vector reg, register offset, extend, scale, Q-form	STR	2	2	I, L01, V	-
Store vector pair, immed offset, S-form	STP, STNP	2	2	L01, V	-

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Store vector pair, immed offset, D-form	STP, STNP	2	2	L01, V	-
Store vector pair, immed offset, Q-form	STP, STNP	2	2	L01, V	-
Store vector pair, immed post-index, S-form	STP	2	2	I, L01, V	-
Store vector pair, immed post-index, D-form	STP	2	2	I, L01, V	-
Store vector pair, immed post-index, Q-form	STP	2	1	I, L01, V	-
Store vector pair, immed pre-index, S-form	STP	2	2	I, L01, V	-
Store vector pair, immed pre-index, D-form	STP	2	2	I, L01, V	-
Store vector pair, immed pre-index, Q-form	STP	2	1	I, L01, V	-

Table 3-24 AArch32 FP store instructions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
FP store, immed offset	VSTR	2	2	L01, V	-
FP store multiple, S-form	VSTMIA, VSTMDB, VPUSH	N + 1	2/R	L01, V	1, 2
FP store multiple, D-form	VSTMIA, VSTMDB, VPUSH	N + 1	2/R	L01, V	1, 2
(FP store, writeback forms)	-	(1)	-	+ I	3

Notes:

1. For store multiple instructions, N = (num_regs/2).
2. R is num_regs.
3. Writeback forms of store instructions require an extra μ OP to update the base address. This update is typically performed in parallel with or prior to the store μ OP (update latency shown in parentheses).

3.16 ASIMD integer instructions

Table 3-25 AArch64 ASIMD integer instructions

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD absolute diff	SABD, UABD	2	2	V	-
ASIMD absolute diff accum	SABA, UABA	3(1)	1	V1	2
ASIMD absolute diff accum long	SABAL(2), UABAL(2)	3(1)	1	V1	2
ASIMD absolute diff long	SABDL(2), UABDL(2)	2	2	V	-
ASIMD arith, basic	ABS, ADD, NEG, SADDL(2), SADDW(2), SHADD, SHSUB, SSUBL(2), SSUBW(2), SUB, UADDL(2), UADDW(2), UHADD, UHSUB, USUBL(2), USUBW(2)	2	2	V	-
ASIMD arith, complex	ADDHN(2), RADDHN(2), RSUBHN(2), SQABS, SQADD, SQNEG, SQSUB, SRHADD, SUBHN(2), SUQADD, UQADD, UQSUB, URHADD, USQADD	2	2	V	-
ASIMD arith, pair-wise	ADDP, SADDLP, UADDLP	2	2	V	-
ASIMD arith, reduce, 4H/4S	ADDV, SADDLV, UADDLV	2	1	V1	-
ASIMD arith, reduce, 8B/8H	ADDV, SADDLV, UADDLV	4	1	V1, V	-
ASIMD arith, reduce, 16B	ADDV, SADDLV, UADDLV	4	1/2	V1	-
ASIMD compare	CMEQ, CMGE, CMGT, CMHI, CMHS, CMLE, CMLT, CMTST	2	2	V	-
ASIMD dot product	SDOT, UDOT	2(1)	2	V	2
ASIMD logical	AND, BIC, EOR, MOV, MVN, ORN, ORR	2	2	V	-

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD max/min, basic and pair-wise	SMAX, SMAXP, SMIN, SMINP, UMAX, UMAXP, UMIN, UMINP	2	2	V	-
ASIMD max/min, reduce, 4H/4S	SMAXV, SMINV, UMAXV, UMINV	2	1	V1	-
ASIMD max/min, reduce, 8B/8H	SMAXV, SMINV, UMAXV, UMINV	4	1	V1, V	-
ASIMD max/min, reduce, 16B	SMAXV, SMINV, UMAXV, UMINV	4	1/2	V1	-
ASIMD multiply	MUL, SQDMULH, SQRDMULH	4	1	V0	-
ASIMD multiply accumulate	MLA, MLS	4(1)	1	V0	1
ASIMD multiply accumulate long	SMLAL(2), SMLSL(2), UMLAL(2), UMLSL(2)	4(1)	1	V0	1
ASIMD multiply accumulate saturating long	SQDMLAL(2), SQDMLSL(2)	4	1	V0	-
ASIMD multiply/multiply long (8x8) polynomial, D-form	PMUL, PMULL(2)	3	1	V0	3
ASIMD multiply/multiply long (8x8) polynomial, Q-form	PMUL, PMULL(2)	3	1	V0	3
ASIMD multiply long	SMULL(2), UMULL(2), SQDMULL(2)	3	1	V0	-
ASIMD pairwise add and accumulate long	SADALP, UADALP	4(1)	1	V1	2
ASIMD shift accumulate	SSRA, SRSRA, USRA, URSRA	4(1)	1	V1	2
ASIMD shift by immed, basic	SHL, SHLL(2), SHRN(2), SSHLL(2), SSHR, SXTL(2), USHLL(2), USHR, UXTL(2)	2	1	V1	-
ASIMD shift by immed and insert, basic	SLI, SRI	2	1	V1	-
ASIMD shift by immed, complex	RSHRN(2), SQRSHRN(2), SQRSHRUN(2), SQSHL{U}, SQSHRN(2), SQSHRUN(2), SRSHR, UQRSHRN(2), UQSHL, UQSHRN(2), URSR	4	1	V1	-

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD shift by register, basic	SSHL, USHL	2	1	V1	-
ASIMD shift by register, complex	SRSHL, SQRSHL, SQSHL, URSHL, UQRSHL, UQSHL	4	1	V1	-

Table 3-26 AArch32 ASIMD integer instructions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD absolute diff	VABD	2	2	V	-
ASIMD absolute diff accum	VABA	4(1)	1	V1	2
ASIMD absolute diff accum long	VABAL	4(1)	1	V1	2
ASIMD absolute diff long	VABDL	2	2	V	-
ASIMD arith, basic	VADD, VADDL, VADDW, VNEG, VSUB, VSUBL, VSUBW	2	2	V	-
ASIMD arith, complex	VABS, VADDHN, VHADD, VHSUB, VQABS, VQADD, VQNEG, VQSUB, VRADDHN, VRHADD, VRSUBHN, VSUBHN	2	2	V	-
ASIMD arith, pair-wise	VPADD, VPADDL	2	2	V	-
ASIMD compare	VCEQ, VCGE, VCGT, VCLE, VTST	2	2	V	-
ASIMD logical	VAND, VBIC, VMVN, VORR, VORN, VEOR	2	2	V	-
ASIMD max/min	VMAX, VMIN, VPMAX, VPMIN	2	2	V	-
ASIMD multiply, D-form	VMUL, VQDMULH, VQRDMULH	4	1	V0	-
ASIMD multiply accumulate	VMLA, VMLS	4(1)	1	V0	1
ASIMD multiply accumulate long	VMLAL, VMLSL	4(1)	1	V0	1
ASIMD multiply accumulate saturating long	VQDMLAL, VQDMLSL	4	1	V0	-
ASIMD multiply/multiply long (8x8) polynomial, D-form	VMUL (.P8), VMULL (.P8)	3	1	V0	-

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD multiply (8x8) polynomial, Q-form	VMUL (.P8)	3	1	V0	-
ASIMD multiply long	VMULL (.S,.I), VQDMULL	4	1	V0	-
ASIMD pairwise add and accumulate	VPADAL	4(1)	1	V1	1
ASIMD shift accumulate	VSRA, VRSRA	4(1)	1	V1	1
ASIMD shift by immed, basic	VMOVL, VSHL, VSHLL, VSHR, VSHRN	2	1	V1	-
ASIMD shift by immed and insert, basic	VSLI, VSRI	2	1	V1	-
ASIMD shift by immed, complex	VQRSHRN, VQRSHRUN, VQSHL{U}, VQSHRN, VQSHRUN, VRSHR, VRSHRN	4	1	V1	-
ASIMD shift by register, basic	VSHL	2	1	V1	-
ASIMD shift by register, complex	VQRSHL, VQSHL, VRSHL	4	1	V1	-

Notes:

1. Multiply-accumulate pipelines support late-forwarding of accumulate operands from similar μ OPs, allowing a typical sequence of integer multiply-accumulate μ OPs to issue one every cycle or one every other cycle (accumulate latency shown in parentheses).
2. Other accumulate pipelines also support late-forwarding of accumulate operands from similar μ OPs, allowing a typical sequence of such μ OPs to issue one every cycle (accumulate latency shown in parentheses).
3. This category includes instructions of the form “PMULL Vd.8H, Vn.8B, Vm.8B” and “PMULL2 Vd.8H, Vn.16B, Vm.16B”.

3.17 ASIMD floating-point instructions

Table 3-27 AArch64 ASIMD floating-point instructions

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD FP absolute value/difference	FABS, FABD	2	2	V	-
ASIMD FP arith, normal	FABD, FADD, FSUB, FADDP	2	2	V	-
ASIMD FP compare	FACGE, FACGT, FCMEQ, FCMGE, FCMGT, FCMLE, FCMLT	2	2	V	-

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD FP convert, long (F16 to F32)	FCVTL(2)	4	1/2	V0	-
ASIMD FP convert, long (F32 to F64)	FCVTL(2)	3	1	V0	-
ASIMD FP convert, narrow (F32 to F16)	FCVTN(2)	4	1/2	V0	-
ASIMD FP convert, narrow (F64 to F32)	FCVTN(2), FCVTXN(2)	3	1	V0	-
ASIMD FP convert, other, D-form F32 and Q-form F64	FCVTAS, FCVTAU, FCVTMS, FCVTMU, FCVTNS, FCVTNU, FCVTPS, FCVTPU, FCVTZS, FCVTZU, SCVTF, UCVTF	3	1	V0	-
ASIMD FP convert, other, D-form F16 and Q-form F32	FCVTAS, VCVTAU, FCVTMS, FCVTMU, FCVTNS, FCVTNU, FCVTPS, FCVTPU, FCVTZS, FCVTZU, SCVTF, UCVTF	4	1/2	V0	-
ASIMD FP convert, other, Q-form F16	FCVTAS, VCVTAU, FCVTMS, FCVTMU, FCVTNS, FCVTNU, FCVTPS, FCVTPU, FCVTZS, FCVTZU, SCVTF, UCVTF	6	1/4	V0	-
ASIMD FP divide, D-form, F16	FDIV	7	1/7	V0	3
ASIMD FP divide, D-form, F32	FDIV	7 to 10	2/9 to 2/7	V0	3
ASIMD FP divide, Q-form, F16	FDIV	10 to 13	1/13 to 1/10	V0	3
ASIMD FP divide, Q-form, F32	FDIV	7 to 10	1/9 to 1/7	V0	3
ASIMD FP divide, Q-form, F64	FDIV	7 to 15	1/14 to 1/7	V0	3
ASIMD FP max/min, normal	FMAX, FMAXNM, FMIN, FMINNM	2	2	V	-
ASIMD FP max/min, pairwise	FMAXP, FMAXNMP, FMINP, FMINNMP	2	2	V	-
ASIMD FP max/min, reduce, D-form, F16	FMAXV, FMAXNMV, FMINV, FMINNMV	3	2	V	-
ASIMD FP max/min, reduce, Q-form, F32 AND F64	FMAXV, FMAXNMV, FMINV, FMINNMV	5	1/2	V	-
ASIMD FP max/min, reduce, Q-form, F16	FMAXV, FMAXNMV, FMINV, FMINNMV	8	1/3	V	-
ASIMD FP multiply	FMUL, FMULX	3	2	V	2

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD FP multiply accumulate	FMLA, FMLS	4 (2)	2	V	1
ASIMD FP negate	FNEG	2	2	V	-
ASIMD FP round, D-form F32 and Q-form, F64	FRINTA, FRINTI, FRINTM, FRINTN, FRINTP, FRINTX, FRINTZ	3	1	V0	-
ASIMD FP round, D-form F16 and Q-form F32	FRINTA, FRINTI, FRINTM, FRINTN, FRINTP, FRINTX, FRINTZ	4	½	V0	-
ASIMD FP round, Q-form F16	FRINTA, FRINTI, FRINTM, FRINTN, FRINTP, FRINTX, FRINTZ	6	1/4	V0	-
ASIMD FP square root, D-form, F16	FSQRT	7	1/7	V0	3
ASIMD FP square root, D-form, F32	FSQRT	7 to 10	2/9 to 2/7	V0	3
ASIMD FP square root, Q-form, F16	FSQRT	11 to 13	1/13 to 1/11	V0	3
ASIMD FP square root, Q-form, F32	FSQRT	7 to 10	1/9 to 1/7	V0	3
ASIMD FP square root, Q-form, F64	FSQRT	7 to 16	1/15 to 1/7	V0	3

Table 3-28 AArch32 ASIMD floating-point instructions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD FP absolute value	VABS	2	2	V	-
ASIMD FP arith	VABD, VADD, VPADD, VSUB	2	2	V	-
ASIMD FP compare	VACGE, VACGT, VACLE, VACLT, VCEQ, VCGE, VCGT, VCLE	2	2	V	-
ASIMD FP convert, integer, D-form	VCVT, VCVTA, VCVTM, VCVTN, VCVTP	3	1	V0	-
ASIMD FP convert, integer, Q-form	VCVT, VCVTA, VCVTM, VCVTN, VCVTP	4	1/2	V0	-
ASIMD FP convert, fixed, D-form	VCVT	3	1	V0	-
ASIMD FP convert, fixed, Q-form	VCVT	4	1/2	V0	-
ASIMD FP convert, half-precision	VCVT	4	1/2	V0	-
ASIMD FP max/min	VMAX, VMIN, VPMAX, VPMIN, VMAXNM, VMINNM	2	2	V	-
ASIMD FP multiply	VMUL, VNMUL	3	2	V	2
ASIMD FP chained multiply accumulate	VMLA, VMLS	5(2)	2	V	1
ASIMD FP fused multiply accumulate	VFMA, VFMS	4(2)	2	V	1
ASIMD FP negate	VNEG	2	2	V	
ASIMD FP round to integral, D-form	VRINTA, VRINTM, VRINTN, VRINTP, VRINTX, VRINTZ	3	1	V0	-
ASIMD FP round to integral, Q-form	VRINTA, VRINTM, VRINTN, VRINTP, VRINTX, VRINTZ	4	1/2	V0	-

Notes:

1. ASIMD multiply-accumulate pipelines support late-forwarding of accumulate operands from similar μ OPs, allowing a typical sequence of floating-point multiply-accumulate μ OPs to issue one every N cycles (accumulate latency N shown in parentheses).
2. ASIMD multiply-accumulate pipelines support late-forwarding of the result from ASIMD FP multiply μ OPs to the accumulate operands of an ASIMD FP multiply-accumulate μ OP. The latter can potentially be issued 1 cycle after the ASIMD FP multiply μ OP has been issued.
3. ASIMD divide and square root operations are performed using an iterative algorithm and block subsequent similar operations to the same pipeline until complete.

3.18 ASIMD miscellaneous instructions

Table 3-29 AArch64 ASIMD miscellaneous instructions

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD bit reverse	RBIT	2	2	V	-
ASIMD bitwise insert	BIF, BIT, BSL	2	2	V	-
ASIMD count	CLS, CLZ, CNT	2	2	V	-
ASIMD duplicate, gen reg	DUP	3	1	M0	-
ASIMD duplicate, element	DUP	2	2	V	-
ASIMD extract	EXT	2	2	V	-
ASIMD extract narrow	XTN	2	2	V	-
ASIMD extract narrow, saturating	SQXTN(2), SQXTUN(2), UQXTN(2)	4	1	V1	-
ASIMD insert, element to element	INS	2	2	V	-
ASIMD move, FP immed	FMOV	2	2	V	-
ASIMD move, integer immed	MOVI	2	2	V	-
ASIMD reciprocal estimate, D-form F32 and F64	FRECPE, FRECPX, FRSQRT, URECPE, URSQRT	3	1	V0	-
ASIMD reciprocal estimate, D-form F16 and Q-form F32	FRECPE, FRECPX, FRSQRT, URECPE, URSQRT	4	½	V0	-
ASIMD reciprocal estimate, Q-form F16	FRECPE, FRECPX, FRSQRT, URECPE, URSQRT	6	¼	V0	-
ASIMD reciprocal step	FRECPS, FRSQRTS	4	2	V	-
ASIMD reverse	REV16, REV32, REV64	2	2	V	-
ASIMD table lookup, 1 or 2 table regs	TBL	2	2	V	-
ASIMD table lookup, 3 table regs	TBL	4	1	V	-
ASIMD table lookup, 4 table regs	TBL	4	2/3	V	-
ASIMD table lookup extension, 1 table reg	TBX	2	2	V	-
ASIMD table lookup extension, 2 table reg	TBX	4	1	V	-
ASIMD table lookup extension, 3 table reg	TBX	6	2/3	V	-

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD table lookup extension, 4 table reg	TBX	6	2/5	V	-
ASIMD transfer, element to gen reg	UMOV, SMOV	2	1	V1	-
ASIMD transfer, gen reg to element	INS	5	1	M0, V	-
ASIMD transpose	TRN1, TRN2	2	2	V	-
ASIMD unzip/zip	UZP1, UZP2, ZIP1, ZIP2	2	2	V	-

Table 3-30 AArch32 ASIMD miscellaneous instructions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD bitwise insert	VBIF, VBIT, VBSL	2	2	V	-
ASIMD count	VCLS, VCLZ, VCNT	2	2	V	-
ASIMD duplicate, core reg	VDUP	3	1	M0	-
ASIMD duplicate, scalar	VDUP	2	2	V	-
ASIMD extract	VEXT	2	2	V	-
ASIMD move, immed	VMOV	2	2	V	-
ASIMD move, register	VMOV	2	2	V	-
ASIMD move, narrowing	VMOVN	2	2	V	-
ASIMD move, saturating	VQMOVN, VQMOVUN	4	1	V1	-
ASIMD reciprocal estimate, D-form F32 and F64	VRECPE, VRSQRTE	3	1	V0	-
ASIMD reciprocal estimate, D-form F16 and Q-form F32	VRECPE, VRSQRTE	4	1/2	V0	-
ASIMD reciprocal estimate, Q-form F16	VRECPE, VRSQRTE	6	¼	V0	-
ASIMD reciprocal step	VRECPS, VRSQRTS	5	2	V	-
ASIMD reverse	VREV16, VREV32, VREV64	2	2	V	-
ASIMD swap	VSWP	4	2/3	V	-
ASIMD table lookup, 1 or 2 table regs	VTBL	2	2	V	-
ASIMD table lookup, 3 table regs	VTBL	4	1	V	-
ASIMD table lookup, 4 table regs	VTBL	6	2/3	V	-

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD table lookup extension, 1 reg	VTBX	2	2	V	-
ASIMD table lookup extension, 2 table reg	VTBX	4	1	V	-
ASIMD table lookup extension, 3 table reg	VTBX	6	2/3	V	-
ASIMD table lookup extension, 4 table reg	VTBX	6	2/5	V	-
ASIMD transfer, scalar to core reg, word	VMOV	2	1	V1	-
ASIMD transfer, scalar to core reg, byte/hword	VMOV	3	1	V1, I	-
ASIMD transfer, core reg to scalar	VMOV	5	1	M0, V	-
ASIMD transpose	VTRN	4	2/3	V	-
ASIMD unzip/zip	VUZP, VZIP	4	2/3	V	-

3.19 ASIMD load instructions

The latencies shown assume the memory access hits in the Level 1 Data Cache. Compared to standard loads, an extra cycle is required to forward results to FP/ASIMD pipelines.

Table 3-31 AArch64 ASIMD load instructions

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD load, 1 element, multiple, 1 reg, D-form	LD1	5	3	L	-
ASIMD load, 1 element, multiple, 1 reg, Q-form	LD1	5	3	L	-
ASIMD load, 1 element, multiple, 2 reg, D-form	LD1	5	3/2	L	-
ASIMD load, 1 element, multiple, 2 reg, Q-form	LD1	5	3/2	L	-
ASIMD load, 1 element, multiple, 3 reg, D-form	LD1	5	1	L	-
ASIMD load, 1 element, multiple, 3 reg, Q-form	LD1	5	1	L	-
ASIMD load, 1 element, multiple, 4 reg, D-form	LD1	6	4/3	L	-
ASIMD load, 1 element, multiple, 4 reg, Q-form	LD1	6	4/3	L	-

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD load, 1 element, one lane, B/H/S	LD1	7	2	L, V	-
ASIMD load, 1 element, one lane, D	LD1	7	2	L, V	-
ASIMD load, 1 element, all lanes, D-form, B/H/S	LD1R	7	2	L, V	-
ASIMD load, 1 element, all lanes, D-form, D	LD1R	7	2	L, V	-
ASIMD load, 1 element, all lanes, Q-form	LD1R	7	2	L, V	-
ASIMD load, 2 element, multiple, D-form, B/H/S	LD2	7	1	L, V	-
ASIMD load, 2 element, multiple, Q-form, B/H/S	LD2	7	1	L, V	-
ASIMD load, 2 element, multiple, Q-form, D	LD2	7	1	L, V	-
ASIMD load, 2 element, one lane, B/H	LD2	7	1	L, V	-
ASIMD load, 2 element, one lane, S	LD2	7	1	L, V	-
ASIMD load, 2 element, one lane, D	LD2	7	1	L, V	-
ASIMD load, 2 element, all lanes, D-form, B/H/S	LD2R	7	1	L, V	-
ASIMD load, 2 element, all lanes, D-form, D	LD2R	7	1	L, V	-
ASIMD load, 2 element, all lanes, Q-form	LD2R	7	1	L, V	-
ASIMD load, 3 element, multiple, D-form, B/H/S	LD3	8	2/3	L, V	-
ASIMD load, 3 element, multiple, Q-form, B/H/S	LD3	8	2/3	L, V	-
ASIMD load, 3 element, multiple, Q-form, D	LD3	8	2/3	L, V	-
ASIMD load, 3 element, one lane, B/H	LD3	8	2/3	L, V	-
ASIMD load, 3 element, one lane, S	LD3	8	2/3	L, V	-
ASIMD load, 3 element, one lane, D	LD3	8	2/3	L, V	-
ASIMD load, 3 element, all lanes, D-form, B/H/S	LD3R	8	2/3	L, V	-
ASIMD load, 3 element, all lanes, D-form, D	LD3R	8	2/3	L, V	-

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD load, 3 element, all lanes, Q-form, B/H/S	LD3R	8	2/3	L, V	-
ASIMD load, 3 element, all lanes, Q-form, D	LD3R	8	2/3	L, V	-
ASIMD load, 4 element, multiple, D-form, B/H/S	LD4	8	1/2	L, V	-
ASIMD load, 4 element, multiple, Q-form, B/H/S	LD4	10	1/4	L, V	-
ASIMD load, 4 element, multiple, Q-form, D	LD4	10	1/4	L, V	-
ASIMD load, 4 element, one lane, B/H	LD4	8	1/2	L, V	-
ASIMD load, 4 element, one lane, S	LD4	8	1/2	L, V	-
ASIMD load, 4 element, one lane, D	LD4	8	1/2	L, V	-
ASIMD load, 4 element, all lanes, D-form, B/H/S	LD4R	8	1/2	L, V	-
ASIMD load, 4 element, all lanes, D-form, D	LD4R	8	1/2	L, V	-
ASIMD load, 4 element, all lanes, Q-form, B/H/S	LD4R	8	1/2	L, V	-
ASIMD load, 4 element, all lanes, Q-form, D	LD4R	8	1/2	L, V	-
(ASIMD load, writeback form)	-	(1)	-	+ I	1

Table 3-32 AArch32 ASIMD load instructions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD load, 1 element, multiple, 1 reg	VLD1	5	3(2)	L	2
ASIMD load, 1 element, multiple, 2 reg	VLD1	5	3(2)	L	2
ASIMD load, 1 element, multiple, 3 reg	VLD1	5	3/2(1)	L	2
ASIMD load, 1 element, multiple, 4 reg	VLD1	5	3/2(1)	L	2
ASIMD load, 1 element, one lane	VLD1	7	2	L, V	2
ASIMD load, 1 element, all lanes, 1 reg	VLD1	7	2	L, V	2
ASIMD load, 1 element, all lanes, 2 reg	VLD1	7	1	L, V	2

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD load, 2 element, multiple, 2 reg	VLD2	7	1	L, V	2
ASIMD load, 2 element, multiple, 4 reg	VLD2	8	1/2	L, V	2
ASIMD load, 2 element, one lane, size 32	VLD2	7	1	L, V	2
ASIMD load, 2 element, one lane, size 8/16	VLD2	7	1	L, V	2
ASIMD load, 2 element, all lanes	VLD2	7	1	L, V	2
ASIMD load, 3 element, multiple, 3 reg	VLD3	8	2/3	L, V	2
ASIMD load, 3 element, one lane, size 32	VLD3	8	2/3	L, V	2
ASIMD load, 3 element, one lane, size 8/16	VLD3	8	2/3	L, V	2
ASIMD load, 3 element, all lanes	VLD3	8	2/3	L, V	2
ASIMD load, 4 element, multiple, 4 reg	VLD4	8	1/2	L, V	2
ASIMD load, 4 element, one lane, size 32	VLD4	8	1/2	L, V	2
ASIMD load, 4 element, one lane, size 8/16	VLD4	8	1/2	L, V	2
ASIMD load, 4 element, all lanes	VLD4	8	1/2	L, V	2
(ASIMD load, writeback form)	-	(1)	-	+I	1

Notes:

1. Writeback forms of load instructions require an extra μ OP to update the base address. This update is typically performed in parallel with the load μ OP (update latency shown in parentheses).
2. Conditional loads go down L01 pipe and the number in parenthesis represents their throughput when different from the unconditional forms.

3.20 ASIMD store instructions

Store MOPs are split into store address and store data μ OPs at dispatch time. Once executed, stores are buffered and committed in the background.

Table 3-33 AArch64 ASIMD store instructions

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD store, 1 element, multiple, 1 reg, D-form	ST1	2	2	L01, V	-
ASIMD store, 1 element, multiple, 1 reg, Q-form	ST1	2	2	L01, V	-

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD store, 1 element, multiple, 2 reg, D-form	ST1	2	2	L01, V	-
ASIMD store, 1 element, multiple, 2 reg, Q-form	ST1	2	1	L01, V	-
ASIMD store, 1 element, multiple, 3 reg, D-form	ST1	2	1	L01, V	-
ASIMD store, 1 element, multiple, 3 reg, Q-form	ST1	2	2/3	L01, V	-
ASIMD store, 1 element, multiple, 4 reg, D-form	ST1	2	1	L01, V	-
ASIMD store, 1 element, multiple, 4 reg, Q-form	ST1	2	1/2	L01, V	-
ASIMD store, 1 element, one lane, B/H/S	ST1	4	1	L01, V	-
ASIMD store, 1 element, one lane, D	ST1	4	1	L01, V	-
ASIMD store, 2 element, multiple, D-form, B/H/S	ST2	4	1	V, L01	-
ASIMD store, 2 element, multiple, Q-form, B/H/S	ST2	5	1/2	V, L01	-
ASIMD store, 2 element, multiple, Q-form, D	ST2	5	1/2	V, L01	-
ASIMD store, 2 element, one lane, B/H/S	ST2	4	1	V, L01	-
ASIMD store, 2 element, one lane, D	ST2	4	1	V, L01	-
ASIMD store, 3 element, multiple, D-form, B/H/S	ST3	5	1/2	V, L01	-
ASIMD store, 3 element, multiple, Q-form, B/H/S	ST3	6	1/3	V, L01	-
ASIMD store, 3 element, multiple, Q-form, D	ST3	6	1/3	V, L01	-
ASIMD store, 3 element, one lane, B/H	ST3	5	1/2	V, L01	-
ASIMD store, 3 element, one lane, S	ST3	5	1/2	V, L01	-
ASIMD store, 3 element, one lane, D	ST3	5	1/2	V, L01	-
ASIMD store, 4 element, multiple, D-form, B/H/S	ST4	5	1/3	V, L01	-
ASIMD store, 4 element, multiple, Q-form, B/H/S	ST4	8	1/6	V, L01	-
ASIMD store, 4 element, multiple, Q-form, D	ST4	8	1/4	V, L01	-

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD store, 4 element, one lane, B/H	ST4	5	2/3	V, L01	-
ASIMD store, 4 element, one lane, S	ST4	5	2/3	V, L01	-
ASIMD store, 4 element, one lane, D	ST4	5	1/2	V, L01	-
(ASIMD store, writeback form)	-	(1)	-	Add I	1

Table 3-34 AArch32 ASIMD store instructions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
ASIMD store, 1 element, multiple, 1 reg	VST1	2	2	L01, V	-
ASIMD store, 1 element, multiple, 2 reg	VST1	2	2	L01, V	-
ASIMD store, 1 element, multiple, 3 reg	VST1	2	1	L01, V	-
ASIMD store, 1 element, multiple, 4 reg	VST1	2	1	L01, V	-
ASIMD store, 1 element, one lane	VST1	4	1	V, L01	-
ASIMD store, 2 element, multiple, 2 reg	VST2	5	2/3	V, L01	-
ASIMD store, 2 element, multiple, 4 reg	VST2	5	1/3	V, L01	-
ASIMD store, 2 element, one lane	VST2	4	1	V, L01	-
ASIMD store, 3 element, multiple, 3 reg	VST3	5	1/2	V, L01	-
ASIMD store, 3 element, one lane, size 32	VST3	4	1/2	V, L01	-
ASIMD store, 3 element, one lane, size 8/16	VST3	4	1/2	V, L01	-
ASIMD store, 4 element, multiple, 4 reg	VST4	5	1/3	V, L01	-
ASIMD store, 4 element, one lane, size 32	VST4	5	2/3	V, L01	-
ASIMD store, 4 element, one lane, size 8/16	VST4	5	2/3	V, L01	-
(ASIMD store, writeback form)	-	(1)	-	+I	1

Notes:

1. Writeback forms of store instructions require an extra μ OP to update the base address. This update is typically performed in parallel with the store μ OP (update latency shown in parentheses).

3.21 Cryptography extensions

Table 3-35 AArch64 Cryptography extensions

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Crypto AES ops	AESD, AESE, AESIMC, AESMC	2	2	V	-
Crypto polynomial (64x64) multiply long	PMULL (2)	2	1	V0	-
Crypto SHA1 hash acceleration ops	SHA1H	2	1	V0	-
Crypto SHA1 hash acceleration ops	SHA1C, SHA1M, SHA1P	4	1	V0	-
Crypto SHA1 schedule acceleration ops	SHA1SU0, SHA1SU1	2	1	V0	-
Crypto SHA256 hash acceleration ops	SHA256H, SHA256H2	4	1	V0	-
Crypto SHA256 schedule acceleration ops	SHA256SU0, SHA256SU1	2	1	V0	-

Table 3-36 AArch32 Cryptography extensions

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
Crypto AES ops	AESD, AESE, AESIMC, AESMC	2	2	V	1
Crypto polynomial (64x64) multiply long	VMULL.P64	2	1	V0	-
Crypto SHA1 hash acceleration ops	SHA1H	2	1	V0	-
Crypto SHA1 hash acceleration ops	SHA1C, SHA1M, SHA1P	4	1	V0	-
Crypto SHA1 schedule acceleration ops	SHA1SU0, SHA1SU1	2	1	V0	-
Crypto SHA256 hash acceleration ops	SHA256H, SHA256H2	4	1	V0	-
Crypto SHA256 schedule acceleration ops	SHA256SU0, SHA256SU1	2	1	V0	-

Notes:

1. Adjacent AESE/AESMC instruction pairs and adjacent AESD/AESIMC instruction pairs will exhibit the performance characteristics described in Section 4.6.

3.22 CRC

Table 3-37 AArch64 CRC

Instruction Group	AArch64 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
CRC checksum ops	CRC32, CRC32C	2	1	M0	1

Table 3-38 AArch32 CRC

Instruction Group	AArch32 Instructions	Execution Latency	Execution Throughput	Utilized Pipelines	Notes
CRC checksum ops	CRC32, CRC32C	2	1	M0	1

Notes:

1. CRC execution supports late-forwarding of the result from a producer μ OP to a consumer μ OP. This results in a 1 cycle reduction in latency as seen by the consumer.

4 Special considerations

4.1 Dispatch constraints

Dispatch of μ OPs from the in-order portion to the out-of-order portion of the microarchitecture includes several constraints. It is important to consider these constraints during code generation to maximize the effective dispatch bandwidth and subsequent execution bandwidth of Cortex-A78C.

The dispatch stage can process up to 6 MOPs per cycle and dispatch up to 12 μ OPs per cycle, with the following limitations on the number of μ OPs of each type that may be simultaneously dispatched.

- Up to 4 μ OPs utilizing the S or B pipelines
- Up to 4 μ OPs utilizing the M pipelines
- Up to 2 μ OPs utilizing the M0 pipelines
- Up to 2 μ OPs utilizing the V0 pipeline
- Up to 2 μ OPs utilizing the V1 pipeline
- Up to 6 μ OPs utilizing the L pipelines

In the event there are more μ OPs available to be dispatched in a given cycle than can be supported by the constraints above, μ OPs will be dispatched in oldest to youngest age-order to the extent allowed by the above.

4.2 Dispatch stall

In the event of a V-pipeline μ OP containing more than 1 quad-word register source, a portion or all of which was previously written as one or multiple single words, that μ OP will stall in dispatch for three cycles. This stall occurs only on the first such instance, and subsequent consumers of the same register will not experience this stall.

4.3 Optimizing general-purpose register spills and fills

Register transfers between general-purpose registers (GPR) and ASIMD registers (VPR) are lower latency than reads and writes to the cache hierarchy, thus it is recommended that GPR registers be filled/spilled to the VPR rather to memory, when possible.

4.4 Optimizing memory routines

To achieve maximum throughput for memory copy (or similar loops), one should do the following:

- Unroll the loop to include multiple load and store operations per iteration, minimizing the overheads of looping.

- Use non-writeback forms of LDP and STP instructions interleaving them like shown in the example below:

```

Loop_start:
    SUBS    x2,x2,#96
    LDP     q3,q4,[x1,#0]
    STP     q3,q4,[x0,#0]
    LDP     q3,q4,[x1,#32]
    STP     q3,q4,[x0,#32]
    LDP     q3,q4,[x1,#64]
    STP     q3,q4,[x0,#64]
    ADD     x1,x1,#96
    ADD     x0,x0,#96
    BGT     Loop_start

```

A recommended copy routine for AArch32 would look like the sequence above but would use LDRD/STRD instructions. Avoid load-/store-multiple instruction encodings (such as LDM and STM).

To achieve maximum throughput on memset, it is recommended that one do the following:

- Unroll the loop to include multiple load and store operations per iteration, minimizing the overheads of looping.

```

Loop_start:
    STP     q1,q3,[x0,#0]
    STP     q1,q3,[x0,#0x20]
    STP     q1,q3,[x0,#0x40]
    STP     q1,q3,[x0,#0x60]
    ADD     x0,x0,#0x80
    SUBS    x2,x2,#0x80
    B.GT    Loop_start

```

To achieve maximum performance on memset to zero, it is recommended that one use DC ZVA instead of STP. An optimal routine might look something like the following:

```

Loop_start:
    SUBS    x2,x2,#0x80
    DC      ZVA,x0
    ADD     x0,x0,#0x40
    DC      ZVA,x0
    ADD     x0,x0,#0x40
    B.GT    Loop_start

```

4.5 Load/Store alignment

The Armv8.2-A architecture allows many types of load and store accesses to be arbitrarily aligned. The Cortex-A78C core handles most unaligned accesses without performance penalties. However, there are cases which could reduce bandwidth or incur additional latency, as described below:

- Load operations that cross a cache-line (64-byte) boundary
- Quad-word load operations that are not 4B aligned
- Store operations that cross a 32B boundary

4.6 AES encryption/decryption

Cortex-A78C can issue two AESE/AESMC/AESD/AESIMC instruction every cycle (fully pipelined) with an execution latency of two cycles. This means encryption or decryption for at least four data chunks should be interleaved for maximum performance:

```
AESE  data0, key0
AESMC data0, data0
AESE  data1, key0
AESMC data1, data1
AESE  data2, key0
AESMC data2, data2
AESE  data3, key1
AESMC data3, data3
AESE  data0, key0
AESMC data0, data0
...
```

Pairs of dependent AESE/AESMC and AESD/AESIMC instructions exhibit higher performance when they are adjacent in the program code and both instructions use the same destination register.

4.7 Region based fast forwarding

The forwarding logic in the V pipelines is optimized to provide optimal latency for instructions which are expected to commonly forward to one another. The effective latency of FP and ASIMD instructions as described in section 3 is increased by one cycle if the producer and consumer instructions are not part of the same forwarding region. These optimized forwarding regions are defined in the following table.

Table 4-1 Optimized forwarding regions

Region	Instruction Types	Notes
1	ASIMD ALU, ASIMD shift, ASIMD/ scalar insert and move, ASIMD abs/cmp/max/min and the ASIMD miscellaneous instructions in table 3-29 and 3-30.	1
2	FP multiply, FP multiply-accumulate, FP compare, FP add/sub and the ASIMD miscellaneous instructions in table 3-29 and 3-30.	1,2,3
3	Crypto SHA1/SHA256	

Notes:

1. Reciprocal step and estimate instructions are excluded from this region.
2. ASIMD extract narrow, saturating instructions are excluded from this region.
3. ASIMD miscellaneous instructions can only be consumers of this region.

The following instructions are not a part of any region:

- FP div/sqrt
- FP convert and rounding
- ASIMD integer mul/mac
- ASIMD reduction

In addition to the regions mentioned in the table above, all floating point and ASIMD instructions can fast forward to FP and ASIMD stores.

More special notes about the forwarding region in table 4-1:

- Fast forwarding will not occur in AArch32 mode if the consuming register's width is greater than that of the producer.
- Element sources used by FP multiply and multiply-accumulate operations cannot be consumers.
- Complex ASIMD shift by immediate/register and shift accumulate instructions cannot be producers (see section 3.16) in region 1.
- ASIMD extract narrow, saturating instructions cannot be producers (see section 3.18) in region 1.
- ASIMD absolute difference accumulate and pairwise add and accumulate instructions cannot be producers (see section 3.16) in region 1.
- For FP producer-consumer pairs, the precision of the instructions should match (single, double or half) in region 2.
- Pair-wise FP instructions cannot be producers or consumers in region 2.

It is not advisable to interleave instructions belonging to different regions. Also, certain instructions can only be producers or consumers in a particular region but not both (see footnote 3 for table 4-1). For example, the code below interleaves producers and consumers from regions 1 and 2. This will result in an additional latency of 1 cycle as seen by FMUL.

FSUB v27.2s, v28.2s, v20.2s – Region 2

FADD v20.2s, v28.2s, v20.2s – Region 2

MOV v27.s[1], v20.s[1] - Region 2 producer but not a region 2 consumer.

FMUL v26.2s, v27.2s, v6.2s – Region 2

4.8 Branch instruction alignment

Branch instruction and branch target instruction alignment and density can affect performance.



For best case performance, avoid placing more than four branch instructions within an aligned 32-byte instruction memory region.

4.9 FPCR self-synchronization

Programmers and compiler writers should note that writes to the FPCR register are self-synchronizing, i.e. its effect on subsequent instructions can be relied upon without an intervening context synchronizing operation.

4.10 Special register access

The Cortex-A78C core performs register renaming for general purpose registers to enable speculative and out-of-order instruction execution. But most special-purpose registers are not renamed. Instructions that read or write non-renamed registers are subjected to one or more of the following additional execution constraints:

- Non-Speculative Execution – Instructions may only execute non-speculatively.
- In-Order Execution – Instructions must execute in-order with respect to other similar instructions or in some cases all instructions.
- Flush Side-Effects – Instructions trigger a flush side-effect after executing for synchronization.

The table below summarizes various special-purpose register read accesses and the associated execution constraints or side-effects.

Table 4-2 Special-purpose register read accesses

Register Read	Non-Speculative	In-Order	Flush Side-Effect	Notes
APSR	Yes	Yes	No	3
CurrentEL	No	Yes	No	-
DAIF	No	Yes	No	-
DLR_ELO	No	Yes	No	-
DSPSR_ELO	No	Yes	No	-
ELR_*	No	Yes	No	-
FPCR	No	Yes	No	-
FPSCR	Yes	Yes	No	2
FPSR	Yes	Yes	No	2
NZCV	No	No	No	1
SP_*	No	No	No	1
SPSel	No	Yes	No	-
SPSR_*	No	Yes	No	-

Notes:

1. The NZCV and SP registers are fully renamed.
2. FPSR/FPSCR reads must wait for all prior instructions that may update the status flags to execute and retire.
3. APSR reads must wait for all prior instructions that may set the Q bit to execute and retire.

The table below summarizes various special-purpose register write accesses and the associated execution constraints or side-effects.

Table 4-3 Special-purpose register write accesses

Register Write	Non-Speculative	In-Order	Flush Side-Effect	Notes
APSR	Yes	Yes	No	4
DAIF	Yes	Yes	No	-
DLR_ELO	Yes	Yes	No	-
DSPSR_ELO	Yes	Yes	No	-
ELR_*	Yes	Yes	No	-
FPCR	Yes	Yes	Maybe	2
FPSCR	Yes	Yes	Maybe	2, 3
FPSR	Yes	Yes	No	3
NZCV	No	No	No	1
SP_*	No	No	No	1
SPSel	Yes	Yes	Yes	-
SPSR_*	Yes	Yes	No	-

Notes:

1. The NZCV and SP registers are fully renamed.

2. If the FPCR/FPSCR write is predicted to change the control field values, it will introduce a barrier which prevents subsequent instructions from executing. If the FPCR/FPSCR write is predicted to not change the control field values, it will execute without a barrier but trigger a flush if the values change.
3. FPSR/FPSCR writes must stall at dispatch if another FPSR/FPSCR write is still pending.
4. APSR writes that set the Q bit will introduce a barrier which prevents subsequent instructions from executing until the write completes.

4.11 Register forwarding hazards

The Armv8-A architecture allows FP/ASIMD instructions to read and write 32-bit S-registers. In AArch32, each S-register corresponds to one half (upper or lower) of an overlaid 64-bit D-register. A Q-register in turn consists of two overlaid D-registers. Register forwarding hazards may occur when one μ OP reads a Q-register operand that has recently been written with one or more S-register results. Consider the following scenario.

```
VADD    S0, S1, S2
VADD    Q6, Q5, Q0
```

The first instruction writes S0, which corresponds to the lowest part of Q0. The second instruction then requires Q0 as an input operand. In this scenario, there is a RAW dependency between the first and the second instructions. In most cases, Cortex-A78C performs slightly worse in such situations.

Cortex-A78C is able to avoid this register-hazard condition for certain cases. The following rules describe the conditions under which a register-hazard can occur.

- The producer writes an S-register (not a D[x] scalar)
- The consumer reads an overlapping Q-register (not as a D[x] scalar)
- The consumer is a FP/ASIMD μ OP (not a store or MOV μ OP)

To avoid unnecessary hazards, it is recommended that the programmer use D[x] scalar writes when populating registers prior to ASIMD operations. For example, either of the following instruction forms would safely prevent a subsequent hazard.

```
VLD1.32 D0[x], [address]
VADD    Q1, Q0, Q2
```

4.12 IT blocks

The Armv8-A architecture performance deprecates some uses of the IT instruction in such a way that software may be written using multiple naïve single instruction IT blocks. It is preferred that software instead generate multi instruction IT blocks rather than single instruction blocks.

4.13 Instruction fusion

Cortex-A78C can accelerate certain instruction pairs in an operation called fusion. Specific AArch64 instruction pairs that can be fused are as follows:

1. CMP/CMN (immediate) + B.cond
2. CMP/CMN (register) + B.cond

3. TST (immediate) + B.cond
4. TST (register) + B.cond
5. BICS (register) + B.cond
6. NOP + Any instruction

The following instruction pairs are fused in both Aarch32 and Aarch64 modes:

1. AESE + AESMC (see Section 4.6 on AES Encryption/Decryption)
2. AESD + AESIMC (see Section 4.6 on AES Encryption/Decryption)

These instruction pairs must be adjacent to each other in program code.

4.14 Mixing Arm and Thumb code

Mixing Arm and Thumb instructions in the same cache-line should be avoided. In particular, old-style interworking veneers to switch from Thumb to Arm state using BX pc may be very slow. This overhead can be reduced by inserting a direct branch or return between indirect branches in one state and code in the other state. For example:

```
BX pc    // Thumb to Arm veneer
B.-2     // never executed
... Arm code
```

However, it is preferable to remove the indirect branch by using only Thumb-2 or Arm code for each veneer.

4.15 Cache maintenance operations

While using set way invalidation operations on L1 cache, it is recommended that software be written to traverse the sets in the inner loop and ways in the out loop.