# ARM<sup>®</sup> CoreLink<sup>™</sup> CCI-500 Cache Coherent Interconnect

Revision: r0p0

**Technical Reference Manual** 



### ARM<sup>®</sup> CoreLink<sup>™</sup> CCI-500 Cache Coherent Interconnect

#### **Technical Reference Manual**

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**Release Information** 

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# Preface

This preface introduces the *ARM*<sup>®</sup> *CoreLink*<sup>™</sup> *CCI-500 Cache Coherent Interconnect Technical Reference Manual.* 

It contains the following:

- *About this book* on page 7.
- *Feedback* on page 10.

### About this book

This book is for the CoreLink CCI-500 Cache Coherent Interconnect.

#### Product revision status

The *rmpn* identifier indicates the revision status of the product described in this book, for example, r1p2, where:

- rm Identifies the major revision of the product, for example, r1.
- pn Identifies the minor revision or modification status of the product, for example, p2.

#### Intended audience

This book is written for system designers, system integrators, and programmers who are designing or programming a System-on-Chip (SoC) that uses the CoreLink CCI-500 Cache Coherent Interconnect.

#### Using this book

This book is organized into the following chapters:

#### **Chapter 1 Introduction**

This chapter provides an overview of the CoreLink CCI-500 Cache Coherent Interconnect.

#### **Chapter 2 Functional Description**

This chapter describes the functionality of the CoreLink CCI-500 Cache Coherent Interconnect.

#### Chapter 3 Programmers Model

This chapter describes the programmers model of the CoreLink CCI-500 Cache Coherent Interconnect

#### Appendix A Signal Descriptions

This appendix describes the external signals of the CoreLink CCI-500 Cache Coherent Interconnect.

#### **Appendix B Revisions**

This appendix describes the technical changes between released issues of this book.

#### Glossary

The ARM Glossary is a list of terms used in ARM documentation, together with definitions for those terms. The ARM Glossary does not contain terms that are industry standard unless the ARM meaning differs from the generally accepted meaning.

See the ARM Glossary for more information.

#### Typographic conventions

#### italic

Introduces special terminology, denotes cross-references, and citations.

bold

Highlights interface elements, such as menu names. Denotes signal names. Also used for terms in descriptive lists, where appropriate.

#### monospace

Denotes text that you can enter at the keyboard, such as commands, file and program names, and source code.

#### <u>mono</u>space

Denotes a permitted abbreviation for a command or option. You can enter the underlined text instead of the full command or option name.

#### monospace italic

Denotes arguments to monospace text where the argument is to be replaced by a specific value.

#### monospace bold

Denotes language keywords when used outside example code.

<and>

Encloses replaceable terms for assembler syntax where they appear in code or code fragments. For example:

MRC p15, 0 <Rd>, <CRn>, <CRm>, <Opcode\_2>

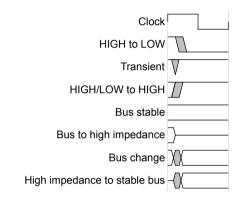
SMALL CAPITALS

Used in body text for a few terms that have specific technical meanings, that are defined in the *ARM glossary*. For example, IMPLEMENTATION DEFINED, IMPLEMENTATION SPECIFIC, UNKNOWN, and UNPREDICTABLE.

#### **Timing diagrams**

The following figure explains the components used in timing diagrams. Variations, when they occur, have clear labels. You must not assume any timing information that is not explicit in the diagrams.

Shaded bus and signal areas are undefined, so the bus or signal can assume any value within the shaded area at that time. The actual level is unimportant and does not affect normal operation.



#### Figure 1 Key to timing diagram conventions

#### Signals

The signal conventions are:

#### Signal level

The level of an asserted signal depends on whether the signal is active-HIGH or active-LOW. Asserted means:

- HIGH for active-HIGH signals.
- LOW for active-LOW signals.

#### Lower-case n

At the start or end of a signal name denotes an active-LOW signal.

#### Additional reading

This section lists publications by ARM and by third parties.

See Infocenter http://infocenter.arm.com, for access to ARM documentation.

#### **ARM** publications

This book contains information that is specific to this product. See the following documents for other relevant information:

- ARM<sup>®</sup> AMBA<sup>®</sup> AXI and ACE Protocol Specification (ARM IHI 0022).
- ARM<sup>®</sup> AMBA<sup>®</sup> APB Protocol Specification (ARM IHI 0024).
- Low Power Interface Specification, ARM<sup>®</sup> Q-Channel and P-Channel Interfaces (ARM IHI 0068).
- *ARM*<sup>®</sup> *CoreSight*<sup>™</sup> *Architecture Specification* (ARM IHI 0029).
- Principles of ARM<sup>®</sup> Memory Maps (ARM DEN 0001).

The following confidential books are only available to licensees:

- ARM<sup>®</sup> CoreLink<sup>™</sup> CCI-500 Cache Coherent Interconnect Configuration and Sign-off Guide (ARM 100024).
- *ARM*<sup>®</sup> *CoreLink*<sup>™</sup> *CCI-500 Cache Coherent Interconnect Integration Manual* (ARM 100025).

#### **Other publications**

This section lists relevant documents published by third parties:

• JEDEC Standard Manufacturer's Identification Code, JEP106 http://www.jedec.org.

# Feedback

#### Feedback on this product

If you have any comments or suggestions about this product, contact your supplier and give:

- The product name.
- The product revision or version.
- An explanation with as much information as you can provide. Include symptoms and diagnostic procedures if appropriate.

#### Feedback on content

If you have comments on content then send an e-mail to errata@arm.com. Give:

- The title.
- The number ARM 100023\_0000\_02\_en.
- The page number(s) to which your comments refer.
- A concise explanation of your comments.

ARM also welcomes general suggestions for additions and improvements.

\_\_\_\_\_ Note \_\_\_\_\_

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# Chapter 1 Introduction

This chapter provides an overview of the CoreLink CCI-500 Cache Coherent Interconnect.

It contains the following sections:

- 1.1 About the CoreLink CCI-500 Cache Coherent Interconnect on page 1-12.
- *1.2 Compliance* on page 1-13.
- 1.3 Features on page 1-14.
- 1.4 Interfaces on page 1-15.
- 1.5 CCI-500 operation on page 1-16.
- 1.6 Configurable options on page 1-17.
- 1.7 Test features on page 1-18.
- 1.8 Product design flow and documentation on page 1-19.
- 1.9 Product revisions on page 1-21.

# 1.1 About the CoreLink CCI-500 Cache Coherent Interconnect

The CCI-500 is a programmable high bandwidth interconnect that enables hardware-coherent systems.

Hardware-managed coherency can improve system performance and reduce system power by sharing onchip data. Managing coherency in hardware has the following benefits:

- Reduces external memory accesses.
- Reduces the software overhead and complexity.
- Enables use of big.LITTLE<sup>™</sup> processing with multiple processor clusters.

The CCI-500 is a configurable interconnect that supports connectivity of:

- Up to four AMBA 4 ACE masters, such as the Cortex®-A57 or Cortex-A53 processors.
- Up to six AMBA 4 ACE-Lite masters, such as the Mali<sup>™</sup>-T760.
- Up to six AMBA 4 AXI4 slaves, such as memory and system peripherals.

\_\_\_\_\_ Note \_\_\_\_\_

The CCI-500 permits combinations of ACE and ACE-Lite masters, up to a maximum total of seven.

The CCI-500 AXI4 master interfaces provide connection to memory and peripheral address space.

# 1.2 Compliance

The CCI-500 complies with the following specifications:

- ARM<sup>®</sup> AMBA<sup>®</sup> AXI and ACE Protocol Specification.
- Low Power Interface Specification, ARM® Q-Channel and P-Channel Interfaces.

This TRM complements architecture reference manuals, architecture specifications, protocol specifications, and relevant external standards. It does not duplicate information from these sources.

# 1.3 Features

The CCI-500 features combine to provide a programmable high-bandwidth interconnect that enables hardware coherent systems.

The CCI-500 provides:

- Data coherency between ACE masters.
- Input and Output (IO) coherency with ACE-Lite masters.
- Crossbar interconnect functionality between the masters and up to six slaves.
- A snoop filter to reduce snoop power and improve performance for snoop misses.
- DVM message transport between masters for communication between MMUs.
- Quality of Service (QoS) features for shaping traffic profiles.
- A Performance Monitoring Unit (PMU) to count performance-related events.
- Support for ARM TrustZone® to provide Secure, Non-secure, and protected states.
- A Programmers View (PV) to control coherency and interconnect functionality.

# 1.4 Interfaces

The CCI-500 has several interfaces to connect it to a wider system.

The CCI-500 is highly-configurable. You can select how many master and slave components to include in your system. The following figure shows an example CCI-500-based system:

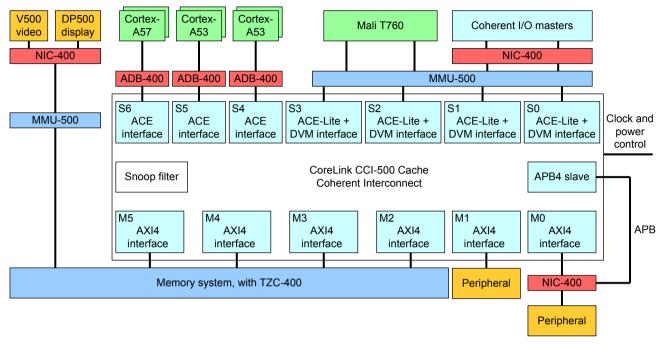


Figure 1-1 Example system with a CCI-500

In this example, slave interfaces S4 to S6 support the ACE protocol for connecting masters such as the Cortex-A53 or Cortex-A57 processors. The CCI-500 manages full coherency and data sharing between processors. Optionally, you can use the ADB-400 asynchronous bridge between components to integrate multiple power domains or clock domains.

Slave interfaces S0 to S3 support ACE-Lite and DVM signaling for connecting I/O coherent devices such as the Mali-T760 graphics unit. You can use DVM signaling for MMUs such as the MMU-500.

You can use the APB4 slave programming interface to program the CCI-500 registers.

Typically, up to four AXI4 master interfaces are connected to compatible memory controllers for LPDDR4 and LPDDR3 memory. Interfaces M5-M2 in the figure show these connections.

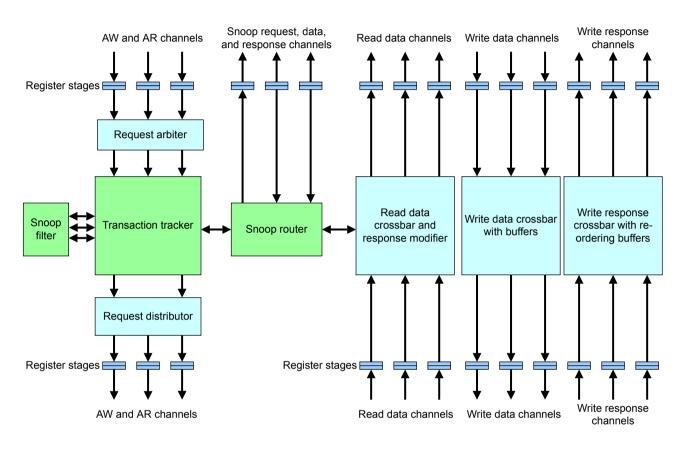
Typically, up to two AXI4 master interfaces are connected to system components, as shown by interfaces M1 and M0.

The clock and power control is achieved using the Q-Channel and the P-Channel.

# 1.5 CCI-500 operation

The CCI-500 has dual-layer request channels, with a central *Transaction Tracker* (TT) that handles coherency and ordering. This is non-blocking and can re-order requests based on QoS requirements.

The following figure shows the CCI-500 high-level operation.



CCI-500 slave interfaces

CCI-500 master interfaces

#### Figure 1-2 CCI-500 high-level operation

The TT uses a snoop filter to determine where to send snoop requests. To maximize throughput:

- The snoop filter has four partitions.
- The read data and write data interconnects are fully-connected crossbars.

Write responses also use a crossbar interconnect and the re-order buffer helps the CCI-500 to meet ordering requirements without stalling requests.

Each interface has a configurable number of register stages, with a minimum of one stage for each interface.

# 1.6 Configurable options

The CCI-500 is highly-configurable. Design-time configuration options enable you to meet your functional requirements with the smallest possible area and power.

Reset-time configuration options enable you to change the functionality of the interconnect for different applications.

When using the CCI-500 in your system, you can configure:

- The number of slave and master interfaces.
- The number of pipeline stages on interfaces to aid timing closure for large designs.
- Write buffering to achieve trade-off between area and write bandwidth.
- The number of outstanding transactions, to achieve trade-off between performance and area.
- QoS threshold, to define the transactions that are treated as high priority within the interconnect.
- Address widths.
- ID widths.
- Burst splitting option for slave interfaces.
- User-defined signal widths.
- The snoop filter RAM capacity to match connected processor cache sizes.
- The transport of data checksums, for example, parity or ECC.

The CCI-500 contains an address map that you can configure at reset-time. It includes options that you can use to interleave memory channels. Optionally, you can implement your own address decoder that defines any arbitrary addressing scheme. The CCI-500 includes a set of assertions that you can use with formal tools or in simulation to verify that your address decoder adheres to CCI-500 requirements.

# 1.7 Test features

The CCI-500 supports both scan cell insertion and MBIST methodologies for your SoC *Design For Test* (DFT) strategy. There are several additional signals to provide high coverage for your test strategy of the CCI-500 design and associated internal RAM cells.

The DFT control signals provide the following capabilities:

- Disabling of internal resets.
- Controlling architectural clock gating.
- Controlling of internal RAM chip-select control to preserve state.
- Controlling of internal RAM MBIST signals.
- Limiting of multi-cycle paths to enable delay testing.

See A.5 DFT signals on page Appx-A-83 for more information.

# 1.8 Product design flow and documentation

You must complete several processes to use the CCI-500. To obtain the best performance from the CCI-500, ARM recommends that you perform some of the implementation stages, including RAM integration, before integrating it into your wider SoC.

The processes you must perform are as follows:

#### Implementation

The implementer configures and synthesizes the RTL to produce a macrocell.

#### Integration

The integrator connects the implemented design into a SoC. This includes connecting it to a memory system and peripherals.

#### **Final SoC implementation**

The process of implementing the final, fully integrated SoC in silicon. ARM can provide only guidance relevant to its own products for this process. If ARM provides guidance on this process for your product, then a separate document is included in the implementation bundle for that product.

#### Programming

This is the last process. The system programmer develops the software required to configure and initialize the CCI-500, and tests the required application software.

For information on the CCI-500 documents that provide information on these processes, see *1.8.1 Documentation* on page 1-19.

Each process:

- Is separate, and a different person can complete it.
- Can include implementation and integration choices that affect the behavior and features of the CCI-500, and therefore the other tasks in the flow.

The operation of the final device depends on:

#### **Build configuration**

The implementer chooses the options that affect the preprocessing of the RTL source files. These options usually include or exclude the logic that affects one or more of the features, the area, or the maximum frequency and performance of the resulting macrocell. For example, the number of outstanding transactions that each master and slave interface supports.

#### **Configuration inputs**

The integrator configures some features of the CCI-500 by tying inputs to specific values. These configurations affect the start-up behavior before you specify the software configuration. They can also limit the options available to the software. For example, the ACCHANNELENSx signal inputs prevent AC coherency requests from being emitted from an unconnected slave interface.

#### Software configuration

The programmer configures the CCI-500 by programming particular values into registers. This affects the behavior of the CCI-500, for example, by enabling QoS features.

#### 1.8.1 Documentation

Each CCI-500 document has an intended audience and is associated with specific tasks in the design flow. These documents do not reproduce ARM architecture and protocol information.

For relevant protocol and architectural information that relates to this product, see *Additional reading* on page 8.

The CoreLink CCI-500 Cache Coherent Interconnect documentation is as follows:

#### **Technical Reference Manual**

The *Technical Reference Manual* (TRM) describes the functionality and the effects of functional options on the behavior of the CCI-500. It is required at all stages of the design flow. The choices made in the design flow can mean that some behaviors described in the TRM are not relevant. If you are programming the CCI-500, then contact:

- The implementer to determine:
  - The build configuration of the implementation.
  - What integration, if any, was performed before implementing the CCI-500.
- The integrator to determine the pin configuration of the device that you use.

#### **Configuration and Sign-off Guide**

The Configuration and Sign-off Guide (CSG) describes:

- A list of the design-time configuration options.
- The available build configuration options and related issues in selecting them.
- How to configure the Register Transfer Level (RTL) with the build configuration options.
- How to integrate RAM arrays.
- How to run test vectors.
- The processes to sign off the configured design.

The ARM product deliverables include reference scripts and information about using them to implement your design. Reference methodology flows supplied by ARM are example reference implementations. Contact your EDA vendor for EDA tool support.

The CSG is a confidential book that is only available to licensees.

#### **Integration Manual**

The Integration Manual (IM) describes how to integrate the CCI-500 into a SoC. It includes:

- A description of the CCI-500 features.
- A list of the reset-time configuration options.
- Considerations when integrating the CCI-500 into your system.

The IM is a confidential book that is only available to licensees.

# 1.9 Product revisions

There can be differences in functionality between different product revisions. ARM records these differences in this section.

r0p0 First release.

# Chapter 2 Functional Description

This chapter describes the functionality of the CoreLink CCI-500 Cache Coherent Interconnect.

It contains the following sections:

- 2.1 About the functions on page 2-23.
- 2.2 Interfaces on page 2-24.
- 2.3 Clocking and reset on page 2-27.
- 2.4 Operation on page 2-28.

# 2.1 About the functions

The CCI-500 is a coherent interconnect that enables hardware coherency. In hardware coherent systems, an operating system can run over multiple processor clusters without complicated cache maintenance software. This is a fundamental requirement for advanced ARM big.LITTLE processing models such as *Global Task Scheduling* (GTS).

In addition to the AXI and ACE interfaces, the CCI-500 provides interfaces that you can use for various system operations, such as:

- Programming the CCI-500 internal registers, debugging, and performance monitoring using an APB4 interface.
- Controlling clock and power states with P-Channel and Q-Channel to minimize power at low bandwidth.
- Logic and RAM testing, for manufacture test.

The CCI-500 includes snoop functionality that permits snooping of the ACE interfaces. A snoop filter provides efficient snoop transaction management by keeping a record of the addresses stored in the caches of the attached ACE masters. This means that the snoop filter can often resolve coherency messaging instead of broadcasting to all ACE interfaces. This can offer system power savings and reduce the latency in the case where data is not held in any of the upstream caches.

A *Performance Monitoring Unit* (PMU) provides events and counters that indicate CCI-500 run-time performance. Additional registers provide information on the current status of the interconnect and you can use these to help debug system deadlock. In addition, the CCI-500 provides a set of QoS regulation and control mechanisms.

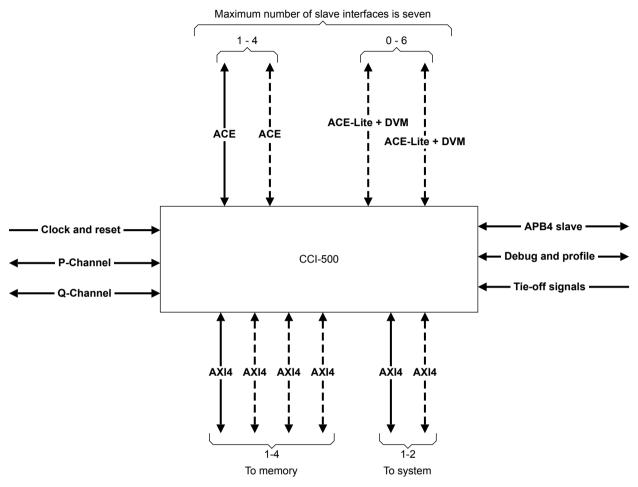
The CCI-500 supports Secure and Non-secure operations that can be used within a system that uses ARM TrustZone to provide Secure, Non-secure, and protected states.

The CCI-500 also supports cache maintenance operations and exclusive accesses.

## 2.2 Interfaces

The CCI-500 has several interfaces to connect it to a wider system.

The following figure shows the CCI-500 interfaces.



#### Figure 2-1 CCI-500 interfaces

All master and slave interfaces are numbered from 0, and the following table shows how many interfaces the CCI-500 can have.

```
_____ Note _____
```

You can have a minimum of two and a maximum of seven slave interfaces.

#### Table 2-1 Permitted CCI-500 interfaces

Interface type	Number of interfaces permitted by the CCI-500		
	Minimum	Maximum	
ACE slave.	1	4	
ACE-Lite + DVM slave.	0	6	
AXI4 master, to memory.	1	4	
AXI4 master, to system.	1	2	

This section contains the following subsections:

- 2.2.1 ACE interfaces on page 2-25.
- 2.2.2 ACE-Lite slave interfaces on page 2-25.
- 2.2.3 AXI4 master interfaces on page 2-25.
- 2.2.4 APB slave interface on page 2-25.
- 2.2.5 Clock and power control interfaces on page 2-25.
- 2.2.6 Debug and performance monitoring interface on page 2-26.
- 2.2.7 DFT interface on page 2-26.

#### 2.2.1 ACE interfaces

You can configure the CCI-500 to have up to four standard AXI Coherency Extensions (ACE) slave interfaces.

The CCI-500 supports the full ACE protocol, with a coherency granule of 64-bytes. For more information about how the CCI-500 handles snoop transactions, see *2.4.3 Snoop connectivity and control* on page 2-28.

See the ARM® AMBA® AXI and ACE Protocol Specification for more information.

#### 2.2.2 ACE-Lite slave interfaces

The ACE-Lite interfaces are a defined subset of the full ACE interfaces.

You can configure the CCI-500 to have up to six ACE-Lite slave interfaces. In addition to standard ACE-Lite functionality, these interfaces also support DVM messages that are passed to upstream system MMU components.

See the ARM<sup>®</sup> AMBA<sup>®</sup> AXI and ACE Protocol Specification.

#### 2.2.3 AXI4 master interfaces

The AMBA 4 protocol defines the AXI4 protocol that supports high-performance, high-frequency system designs.

Depending on configuration, the CCI-500 includes:

- Between one and four AXI4 master interfaces for connecting to memory.
- One or two AXI4 master interfaces for connecting to the rest of the system.

The primary difference between memory and system interfaces is the definition in the memory map. You can configure memory interfaces to be interleaved across the memory region, for example striping, to enable higher utilization of memory. See the *ARM*<sup>®</sup> *AMBA*<sup>®</sup> *AXI and ACE Protocol Specification* for more information about AXI4.

#### 2.2.4 APB slave interface

APB is a low-cost AMBA bus protocol that can reduce power consumption when connecting to the main system bus.

The CCI-500 has an APB slave interface for programming the internal registers and reading from the status, PMU and debug registers. This interface runs synchronously with the other CCI-500 interfaces.

#### 2.2.5 Clock and power control interfaces

The CCI-500 includes the ARM Q-Channel and P-Channel interfaces to control clock and power states.

To save power, the CCI-500 has:

- A Q-Channel interface that you can use to determine how to control the clock state of devices.
- A P-Channel interface that you can use to control the power state of the snoop filter RAMs within the CCI-500, for example, enabling a retention state.

See the *Low Power Interface Specification, ARM*<sup>®</sup> *Q-Channel and P-Channel Interfaces* for more information.

#### 2.2.6 Debug and performance monitoring interface

The CCI-500 supports debug and performance monitoring using a combination of standard signals and a dedicated error pin.

The CCI-500 includes standard ARM interface signals for:

- Debug configuration.
- Event outputs.
- PMU counter overflow interrupts.

In addition, the **nERRIRQ** output pin indicates a transaction error that cannot be signaled precisely. See 2.4.7 *Error responses* on page 2-38 for more information.

#### 2.2.7 DFT interface

The CCI-500 incorporates testing interfaces.

For in-silicon testing of logic and RAMs, the CCI-500 includes DFT and MBIST interfaces. See *A*. *5 DFT signals* on page Appx-A-83 for more information.

# 2.3 Clocking and reset

The CCI-500 uses ACLK and ARESETn signals for clock and reset, respectively.

The CCI-500 has a single main clock signal, ACLK, that it distributes to all sub-blocks. You must use external clock-domain-crossing bridges when masters and slaves connecting to the CCI-500 are in different clock domains.

The CCI-500 has a single reset domain with an active-LOW reset input signal, **ARESETn**. This is synchronized with the **ACLK** with a double-register on the input to the CCI-500 signal. You can replace this synchronizer with appropriate cells from your target library.

\_\_\_\_\_ Note \_\_\_\_\_

You must ensure that there is no activity on the slave interfaces, and the configuration inputs must be static for at least three **ACLK** cycles after the **ARESETn** signal goes HIGH.

The CCI-500 supports the following power saving features:

#### Internal regional clock gating

The CCI-500 automatically gates the clock to internal blocks that do not require the clock. **External architectural clock gating** 

The CCI-500 provides signaling to support implementation of your architectural clock gating strategy. The CCI-500 provides this support using the ARM Q-Channel.

When the Q-Channel is in the Q\_STOPPED state, you can safely disable the clock to the CCI-500. Any incoming transactions are stalled, and any changes to other inputs are not registered until the clock is reapplied. This clock gating can reduce dynamic power to zero when the system is idle.

#### External architectural power state control

The CCI-500 provides signaling to support implementing your architectural power state control. The CCI-500 provides this support using the ARM P-Channel.

See the Low Power Interface Specification, ARM<sup>®</sup> Q-Channel and P-Channel Interfaces for more information.

\_\_\_\_\_ Note \_\_\_\_

The Q-Channel and P-Channel supersede the AXI low-power interface that products such as the CCI-400 use. For legacy IP, the *Low Power Interface Specification, ARM® Q-Channel and P-Channel Interfaces* provides information on linking to IP that contains an AXI low-power interface.

# 2.4 Operation

This section groups information based on the operation of various features of the CCI-500.

This section contains the following subsections:

- 2.4.1 Connectivity and address map on page 2-28.
- 2.4.2 Snoop filter on page 2-28.
- 2.4.3 Snoop connectivity and control on page 2-28.
- 2.4.4 Performance Monitoring Unit on page 2-30.
- 2.4.5 In-silicon debug features on page 2-36.
- 2.4.6 Security on page 2-37.
- 2.4.7 Error responses on page 2-38.
- 2.4.8 Cache maintenance operations on page 2-39.
- 2.4.9 Barriers on page 2-39.
- 2.4.10 Exclusive accesses on page 2-39.
- 2.4.11 DVM messages on page 2-39.
- 2.4.12 Quality of Service on page 2-39.

#### 2.4.1 Connectivity and address map

The interconnect topology and the address map are factors that can affect whether a particular master can communicate with a particular slave.

The CCI-500 is a fully-connected interconnect, meaning that any master can communicate with any slave, subject to the user-defined address map. See your SoC documentation for the address map that your implementation supports.

#### 2.4.2 Snoop filter

The CCI-500 contains an inclusive snoop filter that records the addresses of data stored in the ACE master caches. This means that the filter can respond to the snoop in the case of a miss, and snoop appropriate masters only in the case of a hit. Snoop filter entries are maintained by observing transactions from ACE masters to determine when entries have to be allocated and deallocated.

The snoop filter can respond to many of the coherency requests without it being necessary to broadcast to all ACE interfaces. For example, if the address is not in any cache, the snoop filter responds with a miss and directs the request to memory. If the address is in a processor cache, then it is considered to be a hit and the snoop is directed to the appropriate ACE port containing that address in its cache.

The snoop filter is 8-way set associative and is usually configured to contain between 1.5 and 2 times the number of tags in the attached processor caches. In the case of a way conflict, the existing entry is evicted, and the snoop filter issues a CleanInvalid snoop to the processors that might be holding the evicted lines. This is known as a back-invalidation, and is expected to be a rare occurrence if you configure the snoop filter size as ARM recommends.

The snoop filter is updated by monitoring transactions from the attached masters, which allocate and deallocate data into their caches. In the ACE protocol, the deallocation of clean data is indicated using the Evict transaction. You must ensure that masters connected to the CCI-500 issue Evict transactions when they deallocate clean data. For ARM processors, you can control the issuing of Evict transactions using bit[3] of the L2 Auxiliary Control Register.

#### 2.4.3 Snoop connectivity and control

The CCI-500 has a fully-connected snoop interconnect and a snoop filter for efficient management of snoop request transactions.

You can control whether each interface is enabled for snoop requests and DVM message requests using the 3.3.8 Snoop Control Registers on page 3-62.

A shareable read request from an ACE master, that allocates data to the cache of the master, also allocates an entry in the snoop filter to record that the master has a copy of that data.

For requests for which it might be necessary to retrieve or invalidate data in the cache of another master, the CCI-500 looks up the address in the snoop filter. If the snoop filter indicates that a master has a copy of that data, then either:

- A snoop request is issued, if snoops to the master are enabled.
- The snoop filter entry is updated, if snoops to the master are disabled.

If the snoop filter indicates that no ACE master contains that address, then the request is directed to the appropriate master interface. DVM requests are broadcast through all slave interfaces that are enabled for DVM messages and do not interact with the snoop filter.

The programmable bits of the *3.3.8 Snoop Control Registers* on page 3-62 are LOW at reset. You must program them HIGH for each master in the shareable domain before the CCI-500 receives shareable transactions or DVM messages. Before disabling a master, you must disable snoop and DVM messages for the master by programming the relevant bits of the *3.3.8 Snoop Control Registers* on page 3-62 LOW.

If snoops are sent to interfaces where the master is disabled or not present, the system is likely to deadlock. A hardware mechanism of disabling snoops is provided to prevent software errors causing deadlocks in cases where masters are not present or do not support DVMs. Each slave interface has an **ACCHANNELENSx** signal input that controls whether snoops and DVM messages can be issued from that interface. This input overrides any programmable settings.

\_\_\_\_\_ Note \_\_\_\_

These bits are sampled only at reset, and changing them after ARESETn is HIGH has no effect.

#### Removing a master from the coherent domain

To ensure correct system operation, you must follow a specific process to remove the master from the CCI-500 coherent domain before powering down your master.

Several steps in this process require you to take action on the processor that you want to power down. For these steps, see the appropriate processor documentation.

#### Procedure

- 1. Configure the master so that it does not issue shareable data into its cache, for example by disabling the data cache.
- 2. Clean and invalidate all shareable data from the caches in the master.
- 3. Program the Snoop Control Register to prevent the CCI-500 from sending snoops or DVM messages to the master.
  - See 3.3.8 Snoop Control Registers on page 3-62 for more information.
- 4. Issue a barrier transaction to ensure that the previous step is complete.
- 5. Poll the Status Register to confirm that the Snoop Control Register changes are effected.

#### **Postrequisites**

After you complete these actions, the master is no longer in the coherent domain and you can power it down or disable it.

#### Adding a master to the coherent domain

To ensure correct system operation, you must follow a specific process to add a master to the CCI-500 coherent domain.

Before a master allocates any shareable data into its caches, you must add it to the CCI-500 coherent domain.

#### Procedure

- 1. Enable the master to respond to snoops.
- 2. Program the Snoop Control Register to enable snoops to the master being added.

See 3.3.8 Snoop Control Registers on page 3-62 for more information.

- 3. Issue a barrier transaction to ensure that the previous step is complete.
- 4. Poll the *3.3.3 Status Register* on page 3-55 to confirm that the Snoop Control Register changes are effected.
- 5. Configure the master so that it can issue cacheable, shareable transactions.

#### 2.4.4 Performance Monitoring Unit

The PMU events and counters indicate the run-time performance of the CCI-500.

The CCI-500 includes logic to gather various statistics on the operation of the interconnect during runtime, using events and counters. These events provide useful information about the behavior of the interconnect to use when debugging or profiling traffic.

The PMU provides eight counters. Each counter can count any of the events available in the CCI-500. To keep the PMU logic overhead to a minimum, the absolute count and timing of events might vary slightly. This has a negligible effect except when the counters are enabled for a very short time.

The PMU consists of:

- Performance event counters that are readable through the internal registers.
- A global start or stop bit that enables the counters to increment when HIGH. The default is LOW.
- A global reset bit that resets all counters to zero.
- A parallel event bus, **EVNTBUS**, that you can export from the CCI-500 to capture all events concurrently.
- Eight 32-bit event counters that you can program to count an event from the event bus.
- Input signals **DGBEN** and **NIDEN**. If either is HIGH, the counting and exporting of events is enabled.
- Input signals **DGBEN** and **SPIDEN**, that enable the counting of both Non-secure and Secure events.
- A set of counter overflow outputs, **nEVNTCNTOVERFLOW**, that can raise an interrupt when a number of events have occurred.

The PMU obeys the following rules:

- Each master and slave interface emits events separately from any other interface.
- The snoop filter emits events separately from any interface.
- Events that are marked ACE only, can only fire for ACE interfaces.
- Each event can only fire once per cycle.

This section describes:

- *PMU event list* on page 2-30.
- *PMU registers* on page 2-35.
- Using the PMU on page 2-35.

#### **PMU event list**

The CCI-500 can generate a wide range of events, attributed to a specific interface or globally where they apply to central functions. A list of these events and interface identifiers enables you to identify and then program the events and source locations you want to monitor.

To program the CCI-500 use the code column in each respective table to identify the value to program in to each register field. If you monitor events using the **EVNTBUS**, then use the **EVNTBUS** offset column to identify each position of the bit.

Each event has a 9-bit configuration identifier comprising a source identifier and an event number concatenated {source,number}. Each source identifier is allocated a 4-bit code that indicates the

interface that generated the 5-bit event code. The following tables show the event codes that you can choose for each interface type or global event respectively:

- Table 2-3 5-bit event codes, sources: slave interfaces on page 2-31.
- Table 2-4 5-bit event codes, sources: master interfaces on page 2-33.
- *Table 2-5 5-bit event codes, sources: Global events* on page 2-34.

The following table shows the possible 4-bit source identification codes that you can monitor events on.

#### Table 2-2 Identification code for events

Code[8:5]	Source
0x0	Slave interface 0, SI0
0x1	Slave interface 1, SI1
0x2	Slave interface 2, SI2
0x3	Slave interface 3, SI3
0x4	Slave interface 4, SI4
0x5	Slave interface 5, SI5
0x6	Slave interface 6, SI6
0x7	Reserved
0x8	Master interface 0, MI0
0x9	Master interface 1, MI1
ØxA	Master interface 2, MI2
ØxB	Master interface 3, MI3
ØxC	Master interface 4, MI4
0xD	Master interface 5, MI5
ØxE	Reserved
ØxF	Global

\_\_\_\_\_ Note \_\_\_\_

As CCI-500 is a configurable product not all interfaces might be present, but the source encodings remain the same. If you select an interface that is not present in the specific implementation, then no events are generated.

The following tables show the 5-bit event list and codes for slave interfaces.

#### Table 2-3 5-bit event codes, sources: slave interfaces

Slave event	Code[4:0]	EVNTBUS offset	Secure exempt	ACE only
Read request handshake: any <b>ARVALID</b> , <b>ARREADY</b> HIGH	0x00	0	-	-
Read request:device	0x01	1	-	-
Read request handshake: normal, non-shareable	0x02	2	-	-
Read request handshake: normal, shareable, non- allocating, for example ReadOnce	0x03	3	-	-

#### Table 2-3 5-bit event codes, sources: slave interfaces (continued)

Slave event	Code[4:0]	EVNTBUS offset	Secure exempt	ACE only
Read request handshake: normal, shareable allocating, for example ReadClean, ReadShared, ReadNotSharedDirty, ReadUnique	0x04	4	-	Y
Read request handshake: invalidation, for example MakeUnique, CleanUnique	0x05	5	-	Y
Read request handshake: cache maintenance operation, for example CleanInvalid, MakeInvalid, CleanShared	0x06	6	-	-
Read request handshake: DVM message, any	0x07	7	-	-
Read data handshake: any	0x08	8	Y	-
RVALID, RREADY HIGH				
Read data handshake with <b>RLAST</b> set, for a snoop hit	0x09	9	Y	-
Write request: any	0x0A	10	-	-
AWVALID, AWREADY HIGH				
Write request: device	0x0B	11	-	-
Write request: non-shareable	0x0C	12	-	-
Write request handshake: shareable, for example WriteBack, WriteClean	0x0D	13	-	Y
Write request handshake: shareable, for example WriteLineUnique	0x0E	14	-	-
Write request handshake: shareable, for example WriteUnique	0x0F	15	-	-
Write request handshake, for example Evict	0x10	16	-	Y
Write request handshake, for example WriteEvict Note WriteEvict in not supported in the CCL 500, so does	0x11	17	-	Y
WriteEvict is not supported in the CCI-500, so does not fire.				
Write data beat: any	0x12	18	Y	-
WVALID, WREADY HIGH				
Snoop request: any	0x13	19	-	-
ACVALID, ACREADY HIGH				
Snoop request: read, for example ReadOnce, ReadClean, ReadNotSharedDirty, ReadShared, ReadUnique	0x14	20	-	Y
Snoop request: clean or invalidate, for example MakeInvalid, CleanInvalid, CleanShared	0x15	21	-	Y
Snoop request: Data Transfer bit CRRESP[0] LOW	0x16	22	Y	-
Read request stall:	0x17	23	-	-
ARVALID HIGH				
ARREADY LOW				

Table 2-3 5-bit event codes, se	ources: slave interfaces (continued)
---------------------------------	--------------------------------------

Slave event	Code[4:0]	EVNTBUS offset	Secure exempt	ACE only
Read data stall:	0x18	24	Y	-
RVALID HIGH				
RREADYLOW				
Write request stall:	0x19	25	-	-
AWVALID HIGH				
AWREADY LOW				
Write data stall:	0x1A	26	Y	-
WVALID HIGH				
WREADY LOW				
Write response stall:	0x1B	27	Y	-
BVALID HIGH				
BREADY LOW				
Snoop request stall:	0x1C	28	-	-
ACVALID HIGH				
ACREADY LOW				
Snoop data stall:	0x1D	29	Y	Y
CDVALID HIGH				
CDREADY LOW				
Request stall cycle because of OT transaction limit	0x1E	30	-	-
Read stall because of arbitration	0x1F	31	-	-

The following table shows the 5-bit event list and codes for master interfaces.

#### Table 2-4 5-bit event codes, sources: master interfaces

Master event	Code[4:0]	EVNTBUS offset	Secure exempt
Read data beat: any	0x00	0	Y
Write data beat: any	0x01	1	Y
Read request stall:	0x02	2	-
ARVALID HIGH			
ARREADY LOW			
Read data stall:	0x03	3	Y
RVALID HIGH			
RREADY LOW			
Write request stall:	0x04	4	-
AWVALID HIGH			
AWREADY LOW			

#### Table 2-4 5-bit event codes, sources: master interfaces (continued)

Master event	Code[4:0]	EVNTBUS offset	Secure exempt	
Write data stall:	0x05	5	Y	
WVALID HIGH				
WREADY LOW				
Write response stall:	0x06	6	Y	
<b>BVALID</b> HIGH				
BREADY LOW				

The following table shows the 5-bit event list and codes for Global events.

#### Table 2-5 5-bit event codes, sources: Global events

Global event	Code[4:0]	EVNTBUS offset	Secure exempt
Access to snoop filter bank 0 or 1, any response.	0x00	0	-
Access to snoop filter bank 2 or 3, any response.	0x01	1	-
Access to snoop filter bank 4 or 5, any response.	0x02	2	-
Access to snoop filter bank 6 or 7, any response.	0x03	3	-
Access to snoop filter bank 0 or 1, miss response.	0x04	4	-
Access to snoop filter bank 2 or 3, miss response.	0x05	5	-
Access to snoop filter bank 4 or 5, miss response.	0x06	6	-
Access to snoop filter bank 6 or 7, miss response.	0x07	7	-
Back invalidation from snoop filter.	0x08	8	-
Requests that allocate into a snoop filter bank might be stalled because all ways are used. The snoop filter RAM might be too small.	0x09	9	Y
Stall because TT full, increase TT_DEPTH parameter to avoid performance degradation.	0x0A	10	-
CCI-generated write request.	0x0B	11	-
CD handshake in snoop network, use this to measure snoop data bandwidth. Each event corresponds to 16 bytes of snoop data.	0x0C	12	Y
Request stall because of address hazard.	0x0D	13	-
Snoop request stall because of snoop TT being full.	0x0E	14	Y
Snoop request type override for TZMP1 protection.	0x0F	15	Y

#### **Event bus**

The CCI-500 exports a vector of event signals providing information from the *Performance Monitor Unit* (PMU) using the **EVNTBUS** signal. The width of this bus varies depending on the number of master and slave interfaces in your CCI-500 implementation.

The **EVNTBUS** output is a concatenation of all events that is, global events and events on each *Master Interface* (MI) and *Slave Interface* (SI). The global events are always the least significant bits from [14:0]

irrespective of the number of interfaces. The bit offsets in the **EVNTBUS** output can be found in *Table 2-5 5-bit event codes, sources: Global events* on page 2-34.

— Note —

By default, only events for Non-secure transactions are recorded. However, if the **SPNIDEN** input signal is HIGH, or if both **DBGEN** and **SPIDEN** inputs are HIGH, then the CCI-500 counts and exports both Secure and Non-secure events. Events marked in the tables as Secure exempt do not have a security classification, so they are counted and exported in either case.

#### **PMU registers**

The CCI-500 contains the following performance-related registers:

- 3.3.13 Event Select Registers on page 3-67.
- 3.3.14 Event Count Registers on page 3-68.
- 3.3.15 Count Control Registers on page 3-68.

#### Using the PMU

You can run performance and monitor tests to check the CCI-500 performance.

For each performance and monitor test that you run, you can:

- Select a maximum of eight events to monitor during the test.
- Read the value of each event counter at the end of the test.
- Detect counter overflows.

----- Note -------

The CCI-500 PMU does not include a clock counter because the clock can be disabled to save power. To make time-related measurements, you must use another system timer, for example, the clock counter in the processor PMU.

Use the following registers to set up your test, and to monitor each event:

- Event Select Register to select the event.
- Event Count Register to indicate how many events occur.
- Event Counter Control Register to enable or disable the event counter.
- Event Overflow Flag Status Register to detect the event counter overflow.

#### Example of how to use the PMU

Use the following example to run a test scenario and show how to use the PMU to measure the snoop hit rate for shareable read requests for one ACE master and one ACE-Lite master.

In this example, it is assumed that the ACE master is connected to slave interface 3 and the ACE-Lite master is connected to slave interface 2.

#### Procedure

- 1. Set up the performance counters as follows:
  - a. Program the Event Select Registers as follows:

- Program the Event\_select\_0 register to count shareable, non-allocating read requests through slave interface 3:
  - Program bits[8:5] to 0x3 to select slave interface 3.
  - Program bits[4:0] to 0x03 to select the event for Read request handshake: normal, shareable, non-allocating.
- Program the Event\_select\_1 register to count shareable, allocating read requests through slave interface 3:
  - Program bits[8:5] to 0x3 to select slave interface 3.
  - Program bits[4:0] to 0x04 to select the event for Read request handshake: normal, shareable, non-allocating.
- Program the Event\_select\_2 register to count slave interface 3 snoop hits:
  - Program bits[8:5] to 0x3.
  - Program bits[4:0] to  $0 \times 09$ .
- Program the Event\_select\_3 register to count shareable non-allocating read requests through slave interface 2:
  - Program bits[8:5] to 0x2.
  - Program bits[4:0] to  $0 \times 03$ .
- Program the Event\_select\_4 register to count slave interface 2 snoop hits:
  - Program bits[8:5] to 0x2.
  - Program bits [4:0] to  $0 \times 09$ .
- b. Enable all of the event counters:
  - Program the Count\_control\_0 register bit[0] to 0b1.
  - Program the Count\_control\_1 register bit[0] to 0b1.
  - Program the Count\_control\_2 register bit[0] to 0b1.
  - Program the Count\_control\_3 register bit[0] to 0b1.
  - Program the Count\_control\_4 register bit[0] to 0b1.
- 2. Ensure that the NIDEN and SPNIDEN input are HIGH.
- 3. Program the following bits in the Performance Monitor Control Register (PMCR):
  - Bit[1] to 0b1 to reset event counters.
  - Bit[0] to 0b1 to enable all counters.
- 4. Permit the test to run for an appropriate amount of time.
- 5. Program the PMCR bit[0] to 0b0 to disable all counters to stop the test:
- 6. Read the results of the test from the event counters:
  - Counter 0 and 1 hold the number of shareable reads for slave interface 3.
  - Counter 2 holds the number of snoop hits for slave interface 3.
  - Counter 3 holds the number of shareable reads for slave interface 4.
  - Counter 4 holds the number of snoop hits for slave interface 4.
- 7. Check the overflow bits of all counters and adjust your results accordingly.

#### 2.4.5 In-silicon debug features

The CCI-500 has monitors on all slave and master interfaces that you can use to observe interface status. Each monitor records the number of outstanding read, write, and snoop transactions. It also records the status of the handshake signal from each channel.

This feature can be helpful in the case of a deadlock, for example, the monitors can help to determine outstanding transactions or where back-pressure is being applied.

The monitors are situated inside the outermost registers of the CCI-500, meaning that the values indicated are affected by the numbers of pipeline stages configured in a specific implementation.

—— Note –

If the debug registers are accessed through the CCI-500, you might not be able to read the registers in the case of a deadlock.

## 2.4.6 Security

To build a system based on the Secure and Non-secure capabilities that ARM TrustZone technology provides, you must consider the following security issues.

This section describes:

- Security status of the internal programmers view on page 2-37.
- Making a non-TrustZone aware master Secure on page 2-37.
- Security of master interfaces on page 2-37.
- *Security considerations for the PMU* on page 2-37.

#### Security status of the internal programmers view

You can configure the programmers view of the CCI-500 for access by Secure or Non-secure requests.

With the exception of the PMU registers, the programmers view defaults to Secure access only, as follows:

- Non-secure reads of Secure registers receive an error response and zeroed data.
- Non-secure writes to Secure registers receive an error response and are Write-Ignored (WI).

You can change the security model by writing to the 3.3.2 Secure Access Register on page 3-54. This enables Non-secure access to all registers except the 3.3.1 Control Override Register on page 3-53 and the 3.3.2 Secure Access Register on page 3-54. You can also make the PMU registers accessible to Secure requests only.

## Making a non-TrustZone aware master Secure

For a master that is not TrustZone-aware, you can tie the **ARPROT[1]** and **AWPROT[1]** input signals LOW to place it permanently in the Secure domain. This means that the master can access Secure data in the caches of the ACE masters and secure registers in the CCI-500, so the resulting system might not be secure under all circumstances.

## Security of master interfaces

Transactions from the CCI-500 master interfaces always retain the security setting of the originating transactions.

The security settings of the originating transactions apply to:

- Non-shareable transactions.
- Shareable transactions that miss in the snoop filter or receive a snoop miss response.
- Writes generated by the CCI-500.

## Security considerations for the PMU

You can configure the PMU to count only Non-secure events or both Secure and Non-secure events, depending on the **SPNIDEN**, **SPIDEN** and **DBGEN** input signals.

For more information on configuring, and the input signals **SPNIDEN**, **SPIDEN** and **DBGEN**, see the 2.4.4 Performance Monitoring Unit on page 2-30 section.

If you configure the PMU to count both Secure and Non-secure events, then there is a potential security risk because Non-secure software can observe Secure activity through the performance counters. ARM recommends that you consider the security to be breached for devices placed in this state and take appropriate action.

If the PMU changes from counting all events to counting only Non-secure events, the counters can contain information relating to Secure transactions. Therefore, ARM recommends that the software sets the event counters to zero after changing the configuration to avoid a potential security risk.

\_\_\_\_\_ Note \_\_\_\_

Unlike ARM processors, the CCI-500 makes no distinction between events from user or privileged transactions.

## Support for TrustZone Media Protection

In systems that require hardware protection of media data, you can configure the CCI-500 to support ARM TZMP1.

To differentiate between Protected and Non-Trusted entities, ARM defines 16 states that mark all processes within hardware and software. These states are defined using the Non-secure Access ID (NSAID), and each initiating device in the SoC has one or more NSAID values assigned in hardware. The NSAID enables other components to identify the initiating device for a particular transaction, and to identify whether the device is treated as Non-protected and therefore permitted to read data from other Non-protected masters.

## 2.4.7 Error responses

The CCI-500 uses a combination of precise and imprecise error responses. Precise errors are signaled on the response to the request that caused the error. With the exception of DVM or Evict requests, for accesses to regions that are not mapped in the address decoder, the CCI-500 generates a DECERR response and any snoops, or snoop filter updates, are suppressed. The address map is implementation specific. See your platform documentation for more information.

A snoop error response to a CleanInvalid, CleanShared, or MakeInvalid transaction generates a SLVERR response to the originating device.

There are certain circumstances when it is not possible to signal an error precisely. In these cases, the CCI-500 signals an error imprecisely, using the **nERRIRQ** output pin. You can identify the interface that received the error response by reading the *3.3.4 Imprecise Error Register* on page 3-57.

The following table shows the errors that are signaled as imprecise. All other sources of error are signaled precisely.

----- Note ---

An error is signaled either precisely or imprecisely, but never both.

## Table 2-6 Imprecise errors

Error condition	Channel receiving error	Imprecise error indicator from
A snoop hit response with the error bit set, where data from another snooped master is returned instead of this one.	CR	Slave interface receiving the CR response.
A snoop miss response with the error bit set.	CR	Slave interface receiving the CR response.
Write access generated by the CCI-500.	В	Master interface receiving the B response.
A snoop response with the error bit set where the snoop was generated from a WriteLineUnique or WriteUnique transaction.	CR	Master interface receiving the CR response
A snoop response with the error bit set where the snoop was generated from a back-invalidation.	CR	Slave interface receiving the CR response.

The CCI-500 generates a precise error response for a security violation on a CCI-500 register access. See *2.4.6 Security* on page 2-37.

## 2.4.8 Cache maintenance operations

The CCI-500 supports snooping of cache-maintenance operations based on the Snoop Control Register.

You can use snooping and cache maintenance to manage Level 1 and Level 2 caches within the same domain as the CCI-500. The CCI-500 does not support the propagation of cache maintenance operations downstream of its master interfaces.

## 2.4.9 Barriers

The CCI-500 does not support barrier transactions. You must ensure that barriers are terminated upstream of the CCI. For example, set **SYSBARDISABLE** HIGH in Cortex-A processors.

## 2.4.10 Exclusive accesses

The CCI-500 supports the propagation of exclusive accesses to shareable and non-shareable locations. It does not contain master or slave exclusive access monitors, but does have *Point of Serialization* (PoS) exclusive monitors to avoid livelock.

------ Note -----

- See the *ARM*<sup>®</sup> *AMBA*<sup>®</sup> *AXI and ACE Protocol Specification* for more information on shareable and Non-shareable locations and PoS exclusive monitors.
- The *ARM*<sup>®</sup> *AMBA*<sup>®</sup> *AXI* and *ACE Protocol Specification* permits shareable exclusive accesses on ACE interfaces only.

## 2.4.11 DVM messages

All slave interfaces on the CCI-500 support DVM messages. For ACE-Lite interfaces, this is through the addition of AC and CR channels. Each slave interface has a hardware enable input and programmable enable bit to determine whether it supports the issuing of AC requests for DVM messages.

The 3.3.8 Snoop Control Registers on page 3-62 and 3.3.1 Control Override Register on page 3-53 control DVM message requests.

- Note -

A master that issues DVM messages must also be able to receive DVM messages. The slave interface through which the master connects must have DVM messages enabled.

## 2.4.12 Quality of Service

The CCI-500 provides a set of QoS regulation and control mechanisms.

The following mechanisms are supported:

- *QoS value as a priority indicator* on page 2-39. This is the reservation of resource based on a QoS threshold.
- *Regulation based on outstanding transactions* on page 2-40.

## QoS value as a priority indicator

The CCI-500 uses the QoS value as a priority indicator for arbitration of requests. The QoS value can be from an input to a slave interface, or it can be overwritten by a programmed value.

The CCI-500 uses the QoS value when selecting the request to admit into the main transaction queue. Requests with the highest QoS have the highest priority unless an anti-starvation mechanism is activated. The CCI-500 uses a *Least Recently Granted* (LRG) scheme when two or more transactions share the

highest priority. The arbiter has starvation avoidance mechanisms to prevent high bandwidth requests from stalling lower priority requests indefinitely.

The CCI-500 propagates QoS values. This determines the service rate when downstream interconnect and slave devices are sensitive to the QoS value. The NIC-400 Network Interconnect is sensitive to the QoS value.

----- Note -----

Ensure that you balance the relative priorities of all slave interfaces. For example, setting each one to the highest QoS value reduces the arbitration to LRG, and there is no advantage in using the QoS value.

You can override the **ARQOS** and **AWQOS** input signals from each slave interface by using a programmable register. The value from this register is only applied if the relevant static input signal, **QOSOVERRIDE[6:0]**, is HIGH. CCI-500-generated transactions use the QoS value of the trigger transaction or the override value if the **QOSOVERRIDE** signal is set.

\_\_\_\_\_ Note –

The **QOSOVERRIDE** signal only applies to transactions for which the **ARQOS** or **AWQOS** signals are set to a value of zero. Therefore, each interface can have a mixture of overridden traffic and other traffic, with an unaffected non-zero QoS value.

## Regulation based on outstanding transactions

Each slave interface has a programmable mechanism for limiting the number of outstanding read and write transactions.

An *Outstanding Transaction* (OT) is a read request that has not yet received its last beat of read data, or a write request that has not yet received a response. You can use this mechanism in conjunction with QoS value mechanisms or when the system is not sensitive to the QoS value.

There is a combined OT count for read and write transactions, and this includes all possible request types. Two-part DVM messages count as two outstanding transactions, and transactions that the CCI-500 splits into 64-byte granules count as multiple transactions.

When programming the OT register, the hardware implementation sets the maximum value. This is the value of the register from reset. The minimum value for the OT register is 4, and this is the number of tracker slots reserved for requests from each slave interface to prevent deadlock. If you write a value outside these limits, then the limited value is set and read back.

The OT limit sets a maximum bandwidth for the attached master, based on the average response latency expected from downstream. For ACE masters, from the response to the **RACK** or **WACK** acknowledgement must be included in the response latency. This is approximately:

• OT limit = maximum bandwidth \* average latency / bytes per request

For example, if the average latency between arrival at the main CCI-500 tracking structures and downstream response is 128ns, the maximum required bandwidth is 8GB/s, and requests are 64 bytes in length, then the necessary OT limit for an ACE-Lite master assuming a negligible hit rate is:

• max OT = 8 \* 128 / 64 = 16

In this way, you can allocate memory bandwidth resource amongst various masters in the system.

# Chapter 3 Programmers Model

This chapter describes the programmers model of the CoreLink CCI-500 Cache Coherent Interconnect

It contains the following sections:

- *3.1 About this programmers model* on page 3-42.
- 3.2 Register summary on page 3-43.
- 3.3 Register descriptions on page 3-53.
- 3.4 Address map on page 3-73.

# 3.1 About this programmers model

This section provides general information about the CCI-500 register properties.

The following information applies to the CCI-500 registers:

- The base address is not fixed, and can be different for any particular system implementation. The offset of each register from the base address is fixed.
- Do not attempt to access reserved or unused address locations. Attempting to access these locations can result in UNPREDICTABLE behavior.
- Unless otherwise stated in the accompanying text:
  - Do not modify undefined register bits.
  - Ignore undefined register bits on reads.
  - All register bits are reset to 0 by a system or powerup reset.
- Access type is described as follows:

RW	Read and write.
RO	Read only.
WO	Write only.
RAZ	Read as zero.
WI	Write ignored.

- Bit positions described as reserved are:
  - In an RW register, RAZ/WI
  - In an RO register, RAZ
  - In a WO register, WI.

The CCI-500 registers are accessed using the APB slave interface and cannot be accessed directly through the ACE or ACE-Lite slave interfaces.

The programmers model contains regions for control, slave interface, and performance counter registers. Accesses to unmapped or reserved registers are WI/RAZ. Non-secure accesses to secure registers are WI/RAZ.

The programmers model is not affected by the number of interfaces on a specific CCI-500 configuration. Writing to registers pertaining to interfaces that are not present on your specific implementation has no effect.

# 3.2 Register summary

The register summary lists all CCI-500 registers and some key characteristics.

The following table shows the registers in offset order. The base address of the CCI-500 is not fixed, and can be different for any particular system implementation. Consult your SoC implementation documentation for more information. The offset of each register from the base address is fixed.

## Table 3-1 Register summary

Offset	Name	Туре	Reset	Width	Description
0×00000	Control_Override_Register	RW	0x0000000	32	3.3.1 Control Override Register on page 3-53
0x00008	Secure_Access_Register	RW	0x0000000	32	3.3.2 Secure Access Register on page 3-54
0x0000C	Status_Register	RO	0x00000000 Note	32	<i>3.3.3 Status Register</i> on page 3-55
0x00010	Imprecise_Error_Register	RW	0x0000000	32	3.3.4 Imprecise Error Register on page 3-57
0x00100	Performance_Monitor_Control _Register	RW	0x00004000	32	3.3.5 Performance Monitor Control Register (PMCR) on page 3-60
0x00104	Interface_Monitor_Control _Register	RW	0x0000000	32	3.3.6 Interface Monitor Control Register on page 3-61
0x00FE0	Peripheral_ID0	RO	0x00000022	32	3.3.7 Component and Peripheral ID Registers on page 3-61
0x00FE4	Peripheral_ID1	RO	0x00000B4	32	3.3.7 Component and Peripheral ID Registers on page 3-61
0x00FE8	Peripheral_ID2	RO	0x000000B	32	3.3.7 Component and Peripheral ID Registers on page 3-61
0x00FEC	Peripheral_ID3	RO	0x0000000	32	3.3.7 Component and Peripheral ID Registers on page 3-61

Offset	Name	Туре	Reset	Width	Description
0x00FD0	Peripheral_ID4	RO	0x0000084	32	3.3.7 Component and Peripheral ID Registers on page 3-61
0x00FD4	Peripheral_ID5	RO	0x0000000	32	3.3.7 Component and Peripheral ID Registers on page 3-61
0x00FD8	Peripheral_ID6	RO	0x0000000	32	3.3.7 Component and Peripheral ID Registers on page 3-61
0x00FDC	Peripheral_ID7	RO	0x0000000	32	3.3.7 Component and Peripheral ID Registers on page 3-61
0x00FF0	Component_ID0	RO	0x000000D	32	3.3.7 Component and Peripheral ID Registers on page 3-61
0x00FF4	Component_ID1	RO	0x000000F0	32	3.3.7 Component and Peripheral ID Registers on page 3-61
0x00FF8	Component_ID2	RO	0x0000005	32	3.3.7 Component and Peripheral ID Registers on page 3-61
0x00FFC	Component_ID3	RO	0x00000B1	32	3.3.7 Component and Peripheral ID Registers on page 3-61
0x01000	Snoop_Control_Register_SI0	RW	[31:30] ACCHANNELENS0 [29:0] 0x00000000	32	3.3.8 Snoop Control Registers on page 3-62. Applies to slave interface 0.
0x01004	Shareable_Override_Register _SI0	RW	0x0000000	32	3.3.9 Shareable Override Register on page 3-64. Applies to slave interface 0.

Offset	Name	Туре	Reset	Width	Description
0x01100	Read_QoS_Override_Register _SI0	RW	0x0000000	32	3.3.10 Read Channel QoS Value Override Register on page 3-65. Applies to slave interface 0.
0x01104	Write_QoS_Override_Register _SI0	RW	0x0000000	32	3.3.11 Write Channel QoS Value Override Register on page 3-66. Applies to slave interface 0.
0x01110	Max_OT_Register_SI0	RW	Maximum number of OT supported by this implementation, based on SIO_RW_MAX parameter	32	3.3.12 Maximum Outstanding Transactions Registers on page 3-66. Applies to slave interface 0.
0x02000	Snoop_Control_Register_SI1	RW	[31:30] ACCHANNELENS1 [29:0] Øx0000000	32	3.3.8 Snoop Control Registers on page 3-62. Applies to slave interface 1.
0x02004	Shareable_Override_Register _SI1	RW	0x0000000	32	3.3.9 Shareable Override Register on page 3-64. Applies to slave interface 1.
0x02100	Read_Qos_Override_Register _SI1	RW	0x0000000	32	3.3.10 Read Channel QoS Value Override Register on page 3-65. Applies to slave interface 1.
0x02104	Write_QoS_Override_Register _SI1	RW	0x0000000	32	3.3.11 Write Channel QoS Value Override Register on page 3-66. Applies to slave interface 1.
0x02110	Max_OT_Register_SI1	RW	Maximum number of OT supported by this implementation, based on SI1_RW_MAX parameter	32	3.3.12 Maximum Outstanding Transactions Registers on page 3-66. Applies to slave interface 1.

Offset	Name	Туре	Reset	Width	Description
0x03000	Snoop_Control_Register_SI2	RW	[31:30] ACCHANNELENS2 [29:0] Øx00000000	32	3.3.8 Snoop Control Registers on page 3-62. Applies to slave interface 2.
0x03004	Shareable_Override_Register _SI2	RW	0x0000000	32	3.3.9 Shareable Override Register on page 3-64. Applies to slave interface 2.
0x03100	Read_QoS_Override_Register _SI2	RW	0x0000000	32	3.3.10 Read Channel QoS Value Override Register on page 3-65. Applies to slave interface 2.
0x03104	Write_QoS_Override_Register _SI2	RW	0x0000000	32	3.3.11 Write Channel QoS Value Override Register on page 3-66. Applies to slave interface 2.
0x03110	Max_OT_Register_SI2	RW	Maximum number of OT 32 supported by this implementation, based on SI2_RW-MAX parameter		3.3.12 Maximum Outstanding Transactions Registers on page 3-66. Applies to slave interface 2.
0x04000	Snoop_Control_Register_SI3	RW	[31:30] ACCHANNELENS3 [29:0] 0x00000000	32	3.3.8 Snoop Control Registers on page 3-62. Applies to slave interface 3.
0x04004	Shareable_Override_Register _SI3	RW	0x0000000	32	3.3.9 Shareable Override Register on page 3-64. Applies to slave interface 3.
0x04100	Read_QoS_Override_Register _SI3	RW	0x0000000	32	3.3.10 Read Channel QoS Value Override Register on page 3-65 Applies to slave interface 3.

Offset	Name	Туре	Reset	Width	Description
0x04104	Write_QoS_Override_Register _SI3	RW	0x0000000	32	3.3.11 Write Channel QoS Value Override Register on page 3-66. Applies to slave interface 3.
0x04110	Max_OT_Register_SI3	RW	Maximum number of OT supported by this implementation, based on SI3_RW_MAX parameter	32	3.3.12 Maximum Outstanding Transactions Registers on page 3-66. Applies to slave interface 3.
0x05000	Snoop_Control_Register_SI4	RW	[31:30] ACCHANNELENS4[29:0] 0x0000000	32	3.3.8 Snoop Control Registers on page 3-62. Applies to slave interface 4.
0x05004	Shareable_Override_Register _SI4	RW	0x0000000	32	3.3.9 Shareable Override Register on page 3-64. Applies to slave interface 4.
0x05100	Read_QoS_Override_Register _SI4	RW	0x0000000	32	3.3.10 Read Channel QoS Value Override Register on page 3-65. Applies to slave interface 4.
0x05104	Write_QoS_Override_Register _SI4	RW	0x0000000	32	3.3.11 Write Channel QoS Value Override Register on page 3-66. Applies to slave interface 4.
0x05110	Max_OT_Register_SI4	RW	Maximum number of OT supported by this implementation, based on SI4_RW_MAX parameter	32	3.3.12 Maximum Outstanding Transactions Registers on page 3-66. Applies to slave interface 4.
0x06000	Snoop_Control_Register_SI5	RW	[31:30] ACCHANNELENS5[29:0] 0x0000000	32	3.3.8 Snoop Control Registers on page 3-62. Applies to slave interface 5.

Offset	Name	Туре	Reset	Width	Description
0x06004	Shareable_Override_Register _SI5	RW	0x0000000	32	3.3.9 Shareable Override Register on page 3-64. Applies to slave interface 5.
0x06100	Read_QoS_Override_Register _SI5	RW	0x0000000	32	3.3.10 Read Channel QoS Value Override Register on page 3-65. Applies to slave interface 5.
0x06104	Write_QoS_Override_Register _SI5	RW	0x000000	32	3.3.11 Write Channel QoS Value Override Register on page 3-66. Applies to slave interface 5.
0x06110	Max_OT_Register_SI5	RW	Maximum number of OT supported by this implementation, based on SI5_RW_MAX parameter	32	3.3.12 Maximum Outstanding Transactions Registers on page 3-66. Applies to slave interface 5.
0x07000	Snoop_Control_Register_SI6	RW	[31:30] ACCHANNELENS6 [29:0] 0x0000000	32	3.3.8 Snoop Control Registers on page 3-62. Applies to slave interface 6.
0x07004	Shareable_Override_Register _SI6	RW	0x0000000	32	3.3.9 Shareable Override Register on page 3-64. Applies to slave interface 6.
0x07100	Read_QoS_Override_Register _SI6	RW	0x0000000	32	3.3.10 Read Channel QoS Value Override Register on page 3-65. Applies to slave interface 6.
0x07104	Write_QoS_Override_Register _SI6	RW	0x0000000	32	3.3.11 Write Channel QoS Value Override Register on page 3-66. Applies to slave interface 6.

Offset	Name	Туре	Reset	Width	Description
0x07110	Max_OT_Register_S6	RW	Maximum number of OT supported by this implementation, based on SI6_RW_MAX parameter	32	3.3.12 Maximum Outstanding Transactions Registers on page 3-66. Applies to slave interface 6.
0x10000	Event_Select_0	RW	0x0000000	32	3.3.13 Event Select Registers on page 3-67
0x10004	Event_Count_0	RW	0x0000000	32	3.3.14 Event Count Registers on page 3-68
0x10008	Count_Control_0	RW	0x0000000	32	3.3.15 Count Control Registers on page 3-68
0x1000C	Count_Overflow_0	RW	0x0000000	32	3.3.16 Overflow Flag Status Registers on page 3-69
0x20000	Event_Select_1	RW	0x0000000	32	3.3.13 Event Select Registers on page 3-67
0x20004	Event_Count_1	RW	0x0000000	32	3.3.14 Event Count Registers on page 3-68
0x20008	Count_Control_1	RW	0x0000000	32	3.3.15 Count Control Registers on page 3-68
0x2000C	Count_Overflow_1	RW	0x0000000	32	3.3.16 Overflow Flag Status Registers on page 3-69
0x30000	Event_Select_2	RW	0×0000000	32	3.3.13 Event Select Registers on page 3-67
0x30004	Event_Count_2	RW	0x0000000	32	3.3.14 Event Count Registers on page 3-68
0x30008	Count_Control_2	RW	0x0000000	32	3.3.15 Count Control Registers on page 3-68
0x3000C	Count_Overflow_2	RW	0x0000000	32	3.3.16 Overflow Flag Status Registers on page 3-69

Offset	Name	Туре	Reset	Width	Description
0x40000	Event_Select_3	RW	0×0000000	32	3.3.13 Event Select Registers on page 3-67
0x40004	Event_Count_3	RW	0x0000000	32	3.3.14 Event Count Registers on page 3-68
0x40008	Count_Control_3	RW	0x0000000	32	3.3.15 Count Control Registers on page 3-68
0x4000C	Count_Overflow_3	RW	0×00000000	32	3.3.16 Overflow Flag Status Registers on page 3-69
0x50000	Event_Select_4	RW	0x0000000	32	3.3.13 Event Select Registers on page 3-67
0x50004	Event_Count_4	RW	0x0000000	32	3.3.14 Event Count Registers on page 3-68
0x50008	Count_Control_4	RW	0x0000000	32	3.3.15 Count Control Registers on page 3-68
0x5000C	Count_Overflow_4	RW	0x0000000	32	3.3.16 Overflow Flag Status Registers on page 3-69
0×60000	Event_Select_5	RW	0×0000000	32	3.3.13 Event Select Registers on page 3-67
0x60004	Event_Count_5	RW	0×0000000	32	3.3.14 Event Count Registers on page 3-68
0x60008	Count_Control_5	RW	0×0000000	32	3.3.15 Count Control Registers on page 3-68
0x6000C	Count_Overflow_5	RW	0×0000000	32	3.3.16 Overflow Flag Status Registers on page 3-69
0x70000	Event_Select_6	RW	0x0000000	32	3.3.13 Event Select Registers on page 3-67
0x70004	Event_Count_6	RW	0×00000000	32	3.3.14 Event Count Registers on page 3-68
0x70008	Count_Control_6	RW	0x0000000	32	3.3.15 Count Control Registers on page 3-68

Offset	Name	Туре	Reset	Width	Description
0x7000C	Count_Overflow_6	RW	0x0000000	32	3.3.16 Overflow Flag Status Registers on page 3-69
0×80000	Event_Select_7	RW	0×00000000	32	3.3.13 Event Select Registers on page 3-67
0x80004	Event_Count_7	RW	0x0000000	32	3.3.14 Event Count Registers on page 3-68
0×80008	Count_Control_7	RW	0x0000000	32	3.3.15 Count Control Registers on page 3-68
0×8000C	Count_Overflow_7	RW	0x0000000	32	3.3.15 Count Control Registers on page 3-68
0x90000	Interface Monitor Register SI0	R0	0x0000000	32	3.3.17 Slave Interface Monitor Registers on page 3-69
0x90004	Interface Monitor Register SI1	RO	0x0000000	32	3.3.17 Slave Interface Monitor Registers on page 3-69
0x90008	Interface Monitor Register SI2	RO	0x0000000	32	3.3.17 Slave Interface Monitor Registers on page 3-69
0x9000C	Interface Monitor Register SI3	RO	0×0000000	32	3.3.17 Slave Interface Monitor Registers on page 3-69
0x90010	Interface Monitor Register SI4	RO	0×0000000	32	3.3.17 Slave Interface Monitor Registers on page 3-69
0x90014	Interface Monitor Register SI5	RO	0x0000000	32	3.3.17 Slave Interface Monitor Registers on page 3-69
0x90018	Interface Monitor Register SI6	RO	0×00000000	32	3.3.17 Slave Interface Monitor Registers on page 3-69
0x90100	Interface Monitor Register MI0	RO	0x0000000	32	3.3.17 Slave Interface Monitor Registers on page 3-69
0x90104	Interface Monitor Register MI1	RO	0x0000000	32	3.3.18 Master Interface Monitor Registers on page 3-71

Offset	Name	Туре	Reset	Width	Description
0x90108	Interface Monitor Register MI2	RO	0x0000000	32	3.3.18 Master Interface Monitor Registers on page 3-71
0x9010C	Interface Monitor Register MI3	RO	0x0000000	32	3.3.18 Master Interface Monitor Registers on page 3-71
0x90110	Interface Monitor Register MI4	RO	0x0000000	32	3.3.18 Master Interface Monitor Registers on page 3-71
0x90114	Interface Monitor Register MI5	RO	0x0000000	32	3.3.18 Master Interface Monitor Registers on page 3-71

# 3.3 Register descriptions

Each register description provides information about the register, such as usage constraints, configurations, attributes, and bit assignments.

This section contains the following subsections:

- 3.3.1 Control Override Register on page 3-53.
- 3.3.2 Secure Access Register on page 3-54.
- 3.3.3 Status Register on page 3-55.
- 3.3.4 Imprecise Error Register on page 3-57.
- 3.3.5 Performance Monitor Control Register (PMCR) on page 3-60.
- 3.3.6 Interface Monitor Control Register on page 3-61.
- 3.3.7 Component and Peripheral ID Registers on page 3-61.
- 3.3.8 Snoop Control Registers on page 3-62.
- 3.3.9 Shareable Override Register on page 3-64.
- 3.3.10 Read Channel QoS Value Override Register on page 3-65.
- 3.3.11 Write Channel QoS Value Override Register on page 3-66.
- 3.3.12 Maximum Outstanding Transactions Registers on page 3-66.
- 3.3.13 Event Select Registers on page 3-67.
- 3.3.14 Event Count Registers on page 3-68.
- 3.3.15 Count Control Registers on page 3-68.
- 3.3.16 Overflow Flag Status Registers on page 3-69.
- 3.3.17 Slave Interface Monitor Registers on page 3-69.
- 3.3.18 Master Interface Monitor Registers on page 3-71.

## 3.3.1 Control Override Register

This register provides a fail-safe override for some CCI-500 functions. Use this register to resolve problems that you cannot work around in another way.

## Usage constraints

If you have to write to this register, you must do so before issuing any shareable transactions or DVM messages to the CCI-500. For example, you can do it very early in the boot sequence before installing any Secure OS.

You can access the Control Override Register using Secure transactions only, irrespective of the programming of the *3.3.8 Snoop Control Registers* on page 3-62.

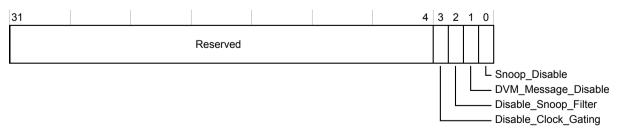
## Configurations

Available in all CCI-500 configurations.

## Attributes

See Table 3-1 Register summary on page 3-43.

The following figure shows the bit assignments.



## Figure 3-1 Control Override Register bit assignments

The following table shows the bit assignments.

#### Table 3-2 Control Override Register bit assignments

		Function			
[31:4]	Reserved	-			
[3]	Disable_Clock_Gating	Disable regional clock gating:			
		<b>0b0</b> Regional clock gating operates in the CCI-500. See 1.7 Test features on page 1-18 and 2.3 Clocking and reset on page 2-27.			
		<b>0b1</b> Disables regional clock gating in the CCI-500.			
[2]	Disable_Snoop_Filter	Disable the snoop filter:			
		<b>0b0</b> Snoop filter operation is defined by the power state input, <b>PSTATE</b> .			
		<b>0b1</b> Disable snoop filter operation.			
[1]	DVM_Message_Disable	DVM message disable:			
		<b>0b0</b> Send DVM messages according to the Snoop Control Registers. See <i>3.3.8 Snoop Control Registers</i> on page 3-62.			
		<b>0b1</b> Disable propagation of all DVM messages.			
[0]	Snoop_Disable	Snoop Disable:			
		<b>0b0</b> Send snoop requests according to the Snoop Control Registers. See <i>3.3.8 Snoop Control Registers</i> on page 3-62.			
		<b>0b1</b> Disables all snoops but not DVM messages.			

#### 3.3.2 Secure Access Register

This register controls whether only Non-secure transactions can read and program the CCI-500 registers.

#### Usage constraints

You can only write to this register only using Secure transactions.

#### Configurations

Available in all CCI-500 configurations.

## Attributes

See Table 3-1 Register summary on page 3-43.

## ——— Warning —

This register enables Non-secure access for all masters to the CCI-500 registers. This compromises the security of your system.

The following figure shows the bit assignments.



## Figure 3-2 Secure Access Register bit assignments

The following table shows the bit assignments.

#### Table 3-3 Secure Access Register bit assignments

Bits	Name	Function		
[31:2]	Reserved	-		
[1]	Debug_monitor_security_override	Debug monitor security override:		
		<b>0b0</b> Enable Non-secure access to the PMU and Interface Monitor Registers.		
		<b>0b1</b> Disable Non-secure access to the PMU and Interface Monitor Registers, unless overridden by bit[0].		
[0]	Secure_access_control	Non-secure register access override:		
		<b>0b0</b> Disable Non-secure access to the CCI-500 registers.		
		<b>0b1</b> Enable Non-secure access to the CCI-500 registers.		

### 3.3.3 Status Register

This register permits snooping to be enabled and disabled safely by indicating when changes made to the Enable\_snoops or Enable\_DVMs bits in the Snoop Control Registers have not taken effect for all transactions outstanding in the system.

When changing these bits, the CCI-500 goes through a transition period where a mixture of transactions with the old value and transactions with the new value are in flight. During this time, the Change\_pending bit stays HIGH. You must wait for the Change\_pending to go LOW before removing or adding masters into the coherency domain. See also *2.4.3 Snoop connectivity and control* on page 2-28.

#### — Note —

You must wait for the completion of write to the Snoop Control Register before testing the Change\_pending bit.

This register indicates whether:

- There are any changes to the enables that have not yet been applied.
- A slave interface has been disabled for future snoop and DVM messages, but has outstanding AC requests.

Other bits in the Status Register indicate:

- Current power state.
- Requested power state.
- Power state change pending.
- Snoop filter initialization phase.

## Usage constraints

There are no usage constraints.

#### Configurations

Available in all CCI-500 configurations.

## Attributes

See Table 3-1 Register summary on page 3-43.

The following figure shows the bit assignments.

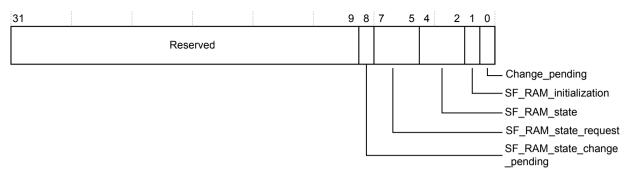


Figure 3-3 Status Register bit assignments

The following table shows the bit assignments.

## Table 3-4 Status Register bit assignments

Name	Function		
Reserved	-		
SF_RAM_state_change_pending	•	Snoop filter RAM power state change pending. This bit reads back the <b>PREQ</b> input.	
	<b>0b0</b> No change or denied.	pending, any previous requests have been accepted	
	<b>0b1</b> State chang	ge is pending and might be accepted or denied.	
SF_RAM_state_request	This is the last requested power state of the snoop filter RAMs Encoding as SF_RAM_state.		
SF_RAM_state	The snoop filter RAM power states are:		
	0b000	Off.	
	0b001	Static snoop filter RAM retention.	
	0b010	Reserved.	
	0b011	Dynamic snoop filter RAM retention.	
	0b100	On.	
	0b101-0b111	Reserved.	
	Not	e	
	This register is re Dynamic retention	eadable only when the interconnect is in either the on or the On state.	
	Reserved SF_RAM_state_change_pending SF_RAM_state_request	Reserved       -         SF_RAM_state_change_pending       Snoop filter RAM the PREQ input.         Øb0       No change or denied.         Øb1       State change         SF_RAM_state_request       This is the last re Encoding as SF_         SF_RAM_state       The snoop filter I         Øb000       Øb011         Øb010       Øb011         Øb101       Øb101         Øb101       Øb101         Øb101       Øb101         Øb101       Øb101	

#### Table 3-4 Status Register bit assignments (continued)

Bits	Name	Function
[1]	SF_RAM_initialization	Indicates when the snoop filter RAM is initialized. Shareable requests are not serviced during this period.
		<b>0b0</b> Snoop filter RAM initialization is complete.
		<b>0b1</b> Snoop filter RAM initialization is in progress.
		Note
		If you use the interconnect to access the CCI-500 registers when the trackers are full of shareable requests waiting for initialization completion, it might not be possible to read this register until initialization is complete.
[0]	Change_pending	Indicates whether any changes to the <i>3.3.8 Snoop Control Registers</i> on page 3-62 or the <i>3.3.1 Control Override Register</i> on page 3-53 are pending in the CCI-500:
		<b>0b0</b> No changes are pending.
		<b>0b1</b> Changes are pending.

## 3.3.4 Imprecise Error Register

This register records the CCI-500 interfaces that have encountered an error that cannot be signaled precisely.

A register bit corresponding to a CCI-500 interface is set when one or more error responses are detected on that interface. Each bit is reset on a write of 1 to that bit.

#### Usage constraints

Accessible using only Secure accesses, unless you set the Secure Access Register. See *3.3.2 Secure Access Register* on page 3-54.

## Configurations

Available in all CCI-500 configurations.

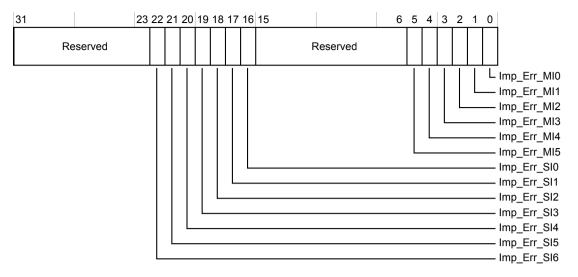
#### Attributes

See Table 3-1 Register summary on page 3-43.

\_\_\_\_\_ Note \_\_\_\_\_

If any Imprecise Error Register bits are set, the **nERRIRQ** signal is asserted, active LOW.

The following figure shows the bit assignments.



#### Figure 3-4 Imprecise Error Register bit assignments

The following table shows the bit assignments.

#### Table 3-5 Imprecise Error Register bit assignments

Bits	Name	Function
[31:23]	Reserved_2	-
[22]	Imp_Err_SI6	Imprecise error indicator for slave interface 6:
		<b>0b0</b> No error from the time this bit was last reset.
		<b>0b1</b> An error response has been received, but not signaled precisely.
[21]	Imp_Err_SI5	Imprecise error indicator for slave interface 5:
		<b>0b0</b> No error from the time this bit was last reset.
		<b>0b1</b> An error response has been received, but not signaled precisely.
[20]	Imp_Err_SI4	Imprecise error indicator for slave interface 4:
		<b>0b0</b> No error from the time this bit was last reset.
		<b>0b1</b> An error response has been received, but not signaled precisely.
[19]	Imp_Err_SI3	Imprecise error indicator for slave interface 3:
		<b>0b0</b> No error from the time this bit was last reset.
		<b>0b1</b> An error response has been received, but not signaled precisely.
[18]	Imp_Err_SI2	Imprecise error indicator for slave interface 2:
		<b>0b0</b> No error from the time this bit was last reset.
		<b>0b1</b> An error response has been received, but not signaled precisely.

## Table 3-5 Imprecise Error Register bit assignments (continued)

Bits	Name	Function
[17]	Imp_Err_SI1	Imprecise error indicator for slave interface 1:
		<b>0b0</b> No error from the time this bit was last reset.
		<b>0b1</b> An error response has been received, but not signaled precisely.
[16]	Imp_Err_SI0	Imprecise error indicator for slave interface 0:
		<b>0b0</b> No error from the time this bit was last reset.
		<b>0b1</b> An error response has been received, but not signaled precisely.
[15:6]	Reserved_1	-
[5]	Imp_Err_MI5	Imprecise error indicator for master interface 5:
		<b>0b0</b> No error from the time this bit was last reset.
		<b>0b1</b> An error response has been received, but not signaled precisely.
[4]	Imp_Err_MI4	Imprecise error indicator for master interface 4:
		<b>0b0</b> No error from the time this bit was last reset.
		<b>0b1</b> An error response has been received, but not signaled precisely.
[3]	Imp_Err_MI3	Imprecise error indicator for master interface 3:
		<b>0b0</b> No error from the time this bit was last reset.
		<b>0b1</b> An error response has been received, but not signaled precisely.
[2]	Imp_Err_MI2	Imprecise error indicator for master interface 2:
		<b>0b0</b> No error from the time this bit was last reset.
		<b>0b1</b> An error response has been received, but not signaled precisely.
[1]	Imp_Err_MI1	Imprecise error indicator for master interface 1:
		<b>0b0</b> No error from the time this bit was last reset.
		<b>0b1</b> An error response has been received, but not signaled precisely.
[0]	Imp_Err_MI0	Imprecise error indicator for master interface 0:
		<b>0b0</b> No error from the time this bit was last reset.
		<b>0b1</b> An error response has been received, but not signaled precisely.

## 3.3.5 Performance Monitor Control Register (PMCR)

This register controls the PMU.

Usage constraints

There are no usage constraints.

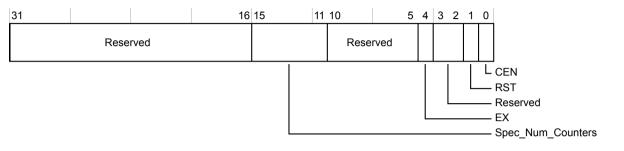
Configurations

Available in all CCI-500 configurations.

## Attributes

See Table 3-1 Register summary on page 3-43.

The following figure shows the bit assignments.



## Figure 3-5 Performance Monitor Register bit assignments

The following table shows the bit assignments.

## Table 3-6 Performance Monitor Control Register bit assignments

Bits	Name	Access	Function
[31:16]	Reserved	-	-
[15:11]	Spec_Num_Counters	RO	Specifies the number of counters implemented
[10:5]	Reserved	-	-
[4]	EX	RW	Enables export of the events to the event bus, <b>EVNTBUS</b> , to permit an external monitoring block to trace events:
			<b>0b0</b> Do not export <b>EVNTBUS</b> .
			Øb1Export EVNTBUS.
[3:2]	Reserved	-	
[1]	RST	RAZ/W	Performance counter reset:
			0b0 No action.
			<b>0b1</b> Reset all performance counters to zero.
[0]	CEN	RW	Enable bit:
			<b>0b0</b> Disable all event counters.
			<b>0b1</b> Enable all event counters.

The following table shows the relationship between the debug enable inputs, **NIDEN** and **DBGEN**, and the PMCR register settings.

—— Note ——

In this table, X can be any value.

NIDEN OR DBGEN	PMCR.CEN	PMCR.EX	Event counters enabled	Events exported
0	Х	Х	No	No
1	0	Х	No	No
1	1	0	Yes	No
1	1	1	Yes	Yes

## Table 3-7 Relationship between NIDEN and DBGEN, and PMCR register settings

## 3.3.6 Interface Monitor Control Register

This register enables all interface monitor control.

Usage constraints There are no usage constraints. Configurations Available in all CCI-500 configurations. Attributes See Table 3-1 Register summary on page 3-43. The following figure shows the bit assignments.



#### Figure 3-6 Interface Monitor Control Register bit assignments

The following table shows the bit assignments.

#### Table 3-8 Interface Monitor Control Register bit assignments

Bits	Name	Access	Function	
[31:1]	Reserved	-	-	
[0]	Enable_Interface_Monitors	RW	Enable bit:	
			<ul><li>0b0 Interface Monitor counters and flags are set to 0.</li><li>0b1 Enable all Interface Monitors.</li></ul>	

## 3.3.7 Component and Peripheral ID Registers

The component and peripheral identity registers are standard JEP106 registers. They provide key information about the CCI-500 hardware, including the product and associated revision number. They also identify ARM as the manufacturer.

These registers are all read-only. Each field is a single byte. This means you must read the most significant 24 bits as zero and only the least significant byte is valid. The least significant 8 bits of the four Component ID registers form a single 32-bit conceptual ID register. In a similar way, the defined fields of the eight Peripheral ID registers form a conceptual 64-bit ID register.

Register	Offset	Bits	Value	Function
Peripheral_ID4	0xFD0	[7:4]	0x8	4KB region count.
		[3:0]	0x4	JEP106 continuation code for ARM.
Peripheral_ID5	0xFD4	[7:0]	0x00	Reserved.
Peripheral_ID6	0xFD8	[7:0]	0x00	Reserved.
Peripheral_ID7	0xFDC	[7:0]	0x00	Reserved.
Peripheral_ID0	0xFE0	[7:0]	0x22	Part number[7:0]
Peripheral_ID1	0xFE4	[7:4]	ØxB	JEP106 ID code[3:0] for ARM.
		[3:0]	0x4	Part number[11:8].
Peripheral_ID2	0xFE8	[7:4]	0x0	CCI-500 revision.
		[3]	0x1	IC uses a manufacturer's identity code allocated by JEDEC according to the JEP106 specification.
		[2:0]	0x3	JEP106 ID code[6:4] for ARM.
Peripheral_ID3	0xFEC	[7:4]	0x0	ARM-approved ECO number. Use the <b>ECOREVNUM</b> inputs to modify this value.
		[3:0]	0x0	Customer modification number. You must not modify this number unless you have permission from ARM.
Component_ID0	0xFF0	[7:0]	0x0D	These values identify the CCI-500 as an ARM component.
Component_ID1	0xFF4	[7:0]	0xF0	
Component_ID2	0xFF8	[7:0]	0x05	
Component_ID3	ØxFFC	[7:0]	0xB1	

#### Table 3-9 Component and Peripheral ID registers bit assignments

#### **ECO revision number**

To track any *Engineering Change Order* (ECO) fixes in the CCI-500, you can change part of the peripheral ID register using the **ECOREVNUM** input pins. You must tie these signals LOW unless you have an ECO from ARM.

The **ECOREVNUM[3:0]** input corresponds to bits[7:4] of the Peripheral ID3 register, MSB to MSB. Driving an input bit HIGH inverts the associated Peripheral ID3 bit.

- Note

ARM recommends that each of the signal drivers is distinct and readily identifiable to facilitate possible metal layer modification.

## 3.3.8 Snoop Control Registers

These registers control the issuing of snoop and DVM requests on slave interfaces.

You can read the register to determine if the interface supports snoops or DVM messages. Enabling snoop or DVM requests on an interface that does not support them has no effect. One Snoop Control Register exists for each slave interface.

#### **Usage constraints**

Accessible using only Secure accesses, unless you set the Secure Access Register to permit Non-secure accesses. See *3.3.2 Secure Access Register* on page 3-54.

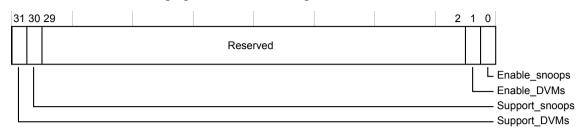
## Configurations

Available in all CCI-500 configurations.

Attributes

## See Table 3-1 Register summary on page 3-43.

The following figure shows the bit assignments.



#### Figure 3-7 Snoop Control Registers bit assignments

The following table shows the bit assignments.

#### Table 3-10 Snoop Control Registers bit assignments

Bits	Name	Reset	Access	Function
[31]	Support_DVMs	ACCHANNELENSx[0] input	R/WI	Slave interface supports DVM messages.
				This is overridden to $0 \times 0$ if you set the Control Override Register bit[1]. See <i>3.3.1 Control Override Register</i> on page 3-53.
[30]	Support_Snoops	ACCHANNELENSx[1] input for ACE interfaces. 0x0 for ACE-Lite interfaces.	R/WI	Slave interface supports snoop requests. This is overridden to $0 \times 0$ if you set the Control Override Register bit[0]. See 3.3.1 Control Override Register on page 3-53. Note This bit only affects the operation of ACE interfaces.
[29:2]	Reserved		-	-

#### Table 3-10 Snoop Control Registers bit assignments (continued)

Bits	Name	Reset	Access	Function
[1]	Enable_DVMs	060	RW	Enable issuing of DVM message requests from this slave interface. RAZ/WI for interfaces not supporting DVM messages:
				<b>0b0</b> Disable DVM message requests.
				<b>0b1</b> Enable DVM message requests.
				Note
				This bit is writable only when bit[31] is set to 1.
[0]	Enable_Snoops	0b0	RW	Enable issuing of snoop requests from this slave interface. RAZ/WI for interfaces not supporting snoops:
				<b>0b0</b> Disable snoop requests.
				<b>0b1</b> Enable snoop requests.
				Note
				<ul> <li>This bit only affects the operation of ACE interfaces.</li> <li>This bit is writable only when bit[30] is set to 1.</li> </ul>

## 3.3.9 Shareable Override Register

This register overrides the shareability characteristics of Normal transactions received on the relevant interface. Overriding of the shareability settings does not occur for FIXED-type bursts, Device transactions, or DVM message transactions.

#### **Usage constraints**

This register is for ACE-Lite slave interfaces only.

Accessible using only Secure accesses, unless you set the Secure Access Register to permit Non-secure accesses. See *3.3.2 Secure Access Register* on page 3-54.

## Configurations

Available in all CCI-500 configurations.

## Attributes

See Table 3-1 Register summary on page 3-43.

---- Note --

Exclusive accesses must not be issued on an interface that is being overridden as shareable. If the CCI-500 is programmed to override transactions as shareable, exclusive accesses are overridden to normal accesses. An exclusive write then receives an OKAY response to indicate that the slave does not support exclusive accesses.

The following figure shows the bit assignments.

31					2	1	0
		R	eserved				
						_	

Domain\_Override —

#### Figure 3-8 Shareable Override Register bit assignments

The following table shows the bit assignments.

#### Table 3-11 Shareable Override Register bit assignments

Bits	Name	Function	
[31:2]	Reserved	-	
[1:0]	Domain_Override	Shareable ov	verride for slave interface:
		0b00-0b01	Do not override AxDOMAIN inputs.
		0b10	<ul> <li>Override AxDOMAIN inputs to 0b00, meaning that all transactions are treated as non-shareable:</li> <li>ReadOnce becomes ReadNoSnoop.</li> <li>WriteUnique and WriteLineUnique become WriteNoSnoop.</li> <li>CleanShared, CleanInvalid, and MakeInvalid transactions do not generate snoops.</li> </ul>
		0b11	<ul> <li>Override AxDOMAIN inputs to 0b01, meaning that all Normal transactions are treated as shareable:</li> <li>ReadNoSnoop becomes ReadOnce.</li> <li>WriteNoSnoop becomes WriteUnique.</li> <li>CleanShared, CleanInvalid, and MakeInvalid transactions generate snoops.</li> </ul>

## 3.3.10 Read Channel QoS Value Override Register

This register stores the override value for the **ARQOS** signal when there is a separate register for each slave interface. This value is applied to transactions when the **QOSOVERRIDE** input signal bit is HIGH for the relevant slave interface and the **ARQOS** input is zero for that request.

#### Usage constraints

Accessible using only Secure accesses, unless you set the *3.3.2 Secure Access Register* on page 3-54 to permit Non-secure accesses.

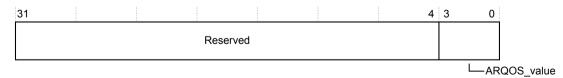
#### Configurations

Available in all CCI-500 configurations.

### Attributes

See Table 3-1 Register summary on page 3-43.

The following figure shows the bit assignments.



#### Figure 3-9 Read Channel QoS Value Override Register bit assignments

The following table shows the bit assignments.

#### Table 3-12 Read Channel QoS Value Override Register bit assignments

Bits	Name	Function
[31:4]	Reserved	-
[3:0]	ARQOS_value	ARQOS value override for the slave interface.
		Note
		This value is applied to transactions with an <b>ARQOS</b> value of zero, if the <b>QOSOVERRIDE</b> input is HIGH for this interface.

## 3.3.11 Write Channel QoS Value Override Register

This register stores the override value for the **AWQOS** signal. If the **QOSOVERRIDE** input bit is HIGH for the relevant slave interface, this override value is applied to requests that have an **AWQOS** value of zero.

The Write Channel QoS Value Override Register characteristics are:

#### Usage constraints

Accessible using only Secure accesses, unless you set the 3.3.2 Secure Access Register on page 3-54 to permit Non-secure accesses.

#### Configurations

Available in all CCI-500 configurations. A separate register exists for each slave interface. Attributes

See Table 3-1 Register summary on page 3-43.

The following figure shows the bit assignments.

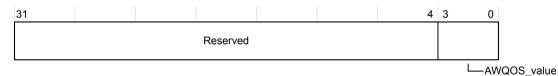


Figure 3-10 Write Channel QoS Value Override Register bit assignments

The following table shows the bit assignments.

## Table 3-13 Write Channel QoS Value Override Register bit assignments

Bits	Name	Function
[31:4]	Reserved	-
[3:0]	AWQOS_value	AWQOS value override for the slave interface.
		Note
		This value is applied to transactions with an <b>AWQOS</b> value of zero, if the <b>QOSOVERRIDE</b> input is HIGH for this interface.

## 3.3.12 Maximum Outstanding Transactions Registers

These registers determine how many *Outstanding Transactions* (OTs) are permitted when the OT regulator is enabled for the relevant slave interface.

#### Usage constraints

If you set the maximum OT size greater than that configured in the RTL, then the value of SIx\_RW\_MAX is written into the register. The minimum value of SIx\_RW\_MAX is 4. Writing values lower than this writes a value of 4 into the register.

Accessible using only Secure accesses, unless you set the *3.3.2 Secure Access Register* on page 3-54 to permit Non-secure accesses.

## Configurations

Available in all CCI-500 configurations. One register exists for each slave interface.

#### Attributes

See Table 3-1 Register summary on page 3-43.

The following figure shows the bit assignments.

31				8 7		0
		Reserved			Max_OT	

#### Figure 3-11 Max OT Register bit assignments

The following table shows the bit assignments.

#### Table 3-14 Max OT Register bit assignments

Bits	Reset	Name	Function
[31:8]	-	Reserved	-
[7:0]	Max_OT	-	The maximum number of OTs for the interface. This is a combined issuing limit. It represents the maximum number of transactions that the upstream master can issue when the AR and AW channels are considered as one issuing source.
			Note
			Additional transactions can be issued into the CCI-500 at the boundary of the device. This is because of the presence of configurable registering between the boundary and the main trackers

## 3.3.13 Event Select Registers

These registers determine the event that a particular counter tracks.

#### Usage constraints

There are no usage constraints.

#### Configurations

Available in all CCI-500 configurations.

One register exists per counter.

#### Attributes

See Table 3-1 Register summary on page 3-43.

— Note —

You can use event counters in different ways, for example:

- To measure traffic across all interfaces by using a counter for each interface.
- To analyze a particular interface by using all the counters to measure a different aspect of the interface.

The following figure shows the bit assignments.

31				9	8	5	4	0
		Reserved			MON_I	NTx	MON_EV	Тx

## Figure 3-12 Event Select Register bit assignments

The following table shows the bit assignments.

#### Table 3-15 Event Select Register bit assignments

Bits	Name	Function
[31:9]	Reserved	-
[8:5]	MON_INTx	Event code that defines the interface to monitor. See <i>PMU event list</i> on page 2-30.
[4:0]	MON_EVTx	Event code that defines the event to monitor. See <i>PMU event list</i> on page 2-30.

## 3.3.14 Event Count Registers

These registers are 32-bit read and write. There is one for each of the eight corresponding event counters. You can reset all event counter values to zero by writing a 1 to the **PMCR** bit[1].

## 3.3.15 Count Control Registers

These registers enable or disable the event counters.

	Usage constrain	ts			
	There an	e no usage constra	aints.		
	Configurations				
	Availab	le in all CCI-500 c	configuration	ns.	
	One reg	ister exists per cou	unter.		
	Attributes	-			
	See Tab	le 3-1 Register su	<i>mmary</i> on p	age 3-43.	
	The following fig	gure shows the bit	assignment	s.	
31					1 0
·		Reserved			

### Figure 3-13 Count Control Register bit assignments

The following table shows the bit assignments.

#### Table 3-16 Count Control Register bit assignments

Name	Function
Reserved	-
CNT_EN	Counter enable:
	<b>0b0</b> Counters disabled.
	<b>0b1</b> Counters enabled.
	Reserved

## 3.3.16 Overflow Flag Status Registers

These registers contain the state of the overflow flags for the event counters.

	Usage constraints There are no usage constraints. Configurations Available in all CCI-500 configurations. One register exists for each event counter. Attributes See Table 3-1 Register summary on page 3-43.
1	The following figure shows the bit assignments.
31	Reserved 1 0

## Figure 3-14 Count Overflow register bit assignments

L CNT\_OVFL\_X

Where x = 0-7.

The following table shows the bit assignments.

#### Table 3-17 Overflow Flag Status Register bit assignments

Bits	Name	Function
[31:1]	Reserved	-
[0]	CNT_OVFL_x	Event counter and overflow flag

When reading this register, any overflow flag that reads as 0 indicates that the counter has not overflowed. An overflow flag that reads as 1 indicates that the counter has overflowed.

When writing to this register, any overflow flag written with a value of 0 is ignored, that is, no change. An overflow flag written with a value of 1 clears the counter overflow flag. The negated counter overflow bits are exported from the CCI-500 on the **nEVNTCNTOVERFLOW**[7:0] signal. You can use this to trigger interrupts. The MSB corresponds to the cycle count overflow.

## 3.3.17 Slave Interface Monitor Registers

These 32-bit RO registers monitor each slave interface.

#### **Usage constraints**

There are no usage constraints.

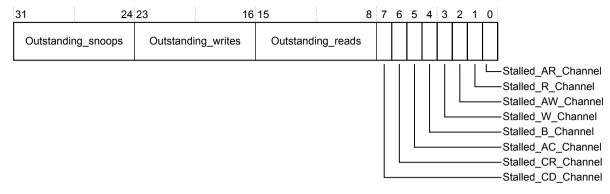
## Configurations

Available in all CCI-500 configurations.

#### Attributes

See Table 3-1 Register summary on page 3-43.

The following figure shows the bit assignments.



## Figure 3-15 Slave Interface Monitor Register bit assignments

The following table shows the bit assignments.

## Table 3-18 Slave Interface Monitor Register bit assignments

Bits	Name	Function
[31:24]	Outstanding_snoops	Number of outstanding snoop requests or DVM messages. From request handshake to response or snoop data for a hit.
[23:16]	Outstanding_writes	Number of outstanding write transactions. From request handshake to response for ACE Lite interfaces or <b>WACK</b> for ACE interfaces.
[15:8]	Outstanding_reads	Number of outstanding read transactions. From request handshake to response or <b>RACK</b> for ACE interfaces.
[7]	Stalled_CD_channel	A transfer is stalled on the CD channel.
		CDVALID is HIGH.
		<b>CDREADY</b> is LOW.
		ACE slave only.
[6]	Stalled_CR_channel	A transfer is stalled on the CR channel.
		CRVALID is HIGH.
		<b>CRREADY</b> is LOW.
[5]	Stalled_AC_channel	A transfer is stalled on the AC channel.
		ACVALID is HIGH.
		ACREADY is LOW.
[4]	Stalled_B_channel	A transfer is stalled on the B channel. <b>BVALID</b> is HIGH
		BREADY is LOW.
[3]	Stalled_W_channel	A transfer is stalled on the W channel.
		WVALID is HIGH.
		WREADY is LOW.
[2]	Stalled_AW_channel	A transfer is stalled on the AW channel.
		AWVALID is HIGH.
		AWREADY is LOW.

#### Table 3-18 Slave Interface Monitor Register bit assignments (continued)

Bits	Name	Function
[1]	Stalled_R_channel	A transfer is stalled on the R channel.
		<b>RVALID</b> is HIGH.
		<b>RREADY</b> is LOW.
[0]	Stalled_AR_channel	A transfer is stalled on the AR channel. ARVALID is HIGH.
		ARREADY is LOW.

#### 3.3.18 Master Interface Monitor Registers

These 32-bit RO registers monitor each master interface.

Usage	constraints
-------	-------------

There are no usage constraints.

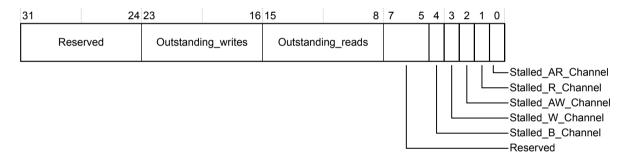
## Configurations

Available in all CCI-500 configurations.

## Attributes

See Table 3-1 Register summary on page 3-43.

The following figure shows the bit assignments.



## Figure 3-16 Master Interface Monitor Register bit assignments

The following table shows the bit assignments.

## Table 3-19 Master Interface Monitor Register bit assignments

Bits	Name	Function
[31:24]	Reserved	-
[23:16]	Outstanding_writes	Number of outstanding write transactions. From request handshake to response.
[15:8]	Outstanding_reads	Number of outstanding read transactions. From request handshake to response.
[7:5]	Reserved	-
[4]	Stalled_B_channel	A transfer is stalled on the B channel. <b>BVALID</b> is HIGH
		BREADY is LOW.
[3]	Stalled_W_channel	A transfer is stalled on the W channel.
		WVALID is HIGH.
		WREADY is LOW.

## Table 3-19 Master Interface Monitor Register bit assignments (continued)

Bits	Name	Function
[2]	Stalled_AW_channel	A transfer is stalled on the AW channel.
		AWVALID is HIGH.
		AWREADY is LOW.
[1]	Stalled_R_channel	A transfer is stalled on the R channel.
		<b>RVALID</b> is HIGH.
		<b>RREADY</b> is LOW.
[0]	Stalled_AR_channel	A transfer is stalled on the AR channel. <b>ARVALID</b> is HIGH.
		ARREADY is LOW.

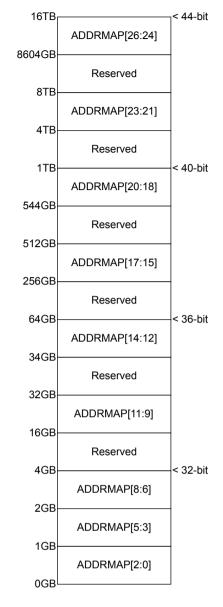
# 3.4 Address map

The CCI-500 has a global address map. This means that, for the example decoder provided with the CCI-500, each master attached to the interconnect has the same view of memory. You can modify the supplied address decoder to define different address maps for each slave interface, or for reads and writes.

\_\_\_\_\_ Note \_\_\_\_\_

The following text and diagram describes the operation of the address decoder supplied with CCI-500. However, the implementer is permitted to modify the address decoder. See your platform documentation to determine the address map for a particular implementation.

The example decoder supplied with the CCI-500 defines address regions, as specified in the document *Principles of ARM® Memory Maps*. The CCI-500 address width is configurable, so the address map extends up to 44 bits. If your implementation uses a smaller address width, then some regions are not addressable. *Figure 3-17 Example decoder address regions* on page 3-74 shows the region sizes and offsets, with associated **ADDRMAP** inputs and address width limits.



## Figure 3-17 Example decoder address regions

In the example decoder, accesses to reserved regions generate a DECERR response.

For each non-reserved region, accesses are mapped to one of the master interfaces, or can be striped across a number of master interfaces. The mapping is determined using configuration input signals, **ADDRMAP**, that are sampled at reset. In the example address map, there are three **ADDRMAP** bits per region, with the encoding defined as:

#### Table 3-20 Decoder mapping

ADDRMAP[2:0	ADDRMAP[2:0] Decode			
0b000	Master interface 0			
0b001	Master interface 1			
0b010	Master interface 2			
0b011	Master interface 3			
0b100	Master interface 4			

# Table 3-20 Decoder mapping (continued)

ADDRMAP[2:0]	Decode
0b101	Master interface 5
0b110	Reserved
Øb111	<ul> <li>The behavior depends on the number of memory interfaces that are configured:</li> <li>With one memory port, all accesses in the region are to that port.</li> <li>With two memory ports, striping occurs across both ports.</li> <li>With three memory ports, striping occurs across the two highest numbered ports.</li> <li>With four memory ports, striping occurs across all four ports.</li> </ul>

If the **ADDRMAP** input maps a region to a master interface that is not present in the specific configuration, then a DECERR response is generated for any requests that target that region.

The example decoder uses a stripe size of 256 bytes.

# Appendix A Signal Descriptions

This appendix describes the external signals of the CoreLink CCI-500 Cache Coherent Interconnect.

It contains the following sections:

- A.1 Clock and reset signals on page Appx-A-77.
- *A.2 Power and clock control signals* on page Appx-A-78.
- A.3 Configuration signals on page Appx-A-80.
- *A.4 Debug signals* on page Appx-A-82.
- *A.5 DFT signals* on page Appx-A-83.
- *A.6 APB4 signals* on page Appx-A-84.
- *A.7 ACE and ACE-Lite slave interface signals* on page Appx-A-85.
- A.8 AXI Master interface signals on page Appx-A-90.
- A.9 Miscellaneous signals on page Appx-A-93.

# A.1 Clock and reset signals

The CCI-500 uses a single set of standard clock and reset signals.

The following table shows the clock and reset signals.

# Table A-1 Clock and reset signals

Signal	Direction	Description
ACLK	Input	Global clock.
ARESETn	Input	Global reset.

# A.2 Power and clock control signals

The CCI-500 uses a range of signals to communicate with the Q-Channel and P-Channel interfaces. The following table shows the power and clock control signals.

## Table A-2 Power and clock control signals

Signal	Direction	Description	
AWAKEUPSx	Input	of the associated CCI-500 takes A clock is enabled. Where $x = 0-6$ , c	Asfers are pending on the AR, AW, or W channels slave interface. If any of these inputs is HIGH, the CLKQACTIVE HIGH to request that the CCI . There is one input for each slave interface. depending on the configuration. te
AWAKEUPMy	Output	of the associated request that the There is one out	
PWAKEUP	Input	Indicates that the transaction is inc	e APB interface requires a clock because a coming.
ACLKQREQn	Input		le the ACLK input. If the clock control channel is e ACLKQREQn to 0b1.
ACLKQACCEPTn	Output	Clock disable ac	ceptance response.
ACLKQDENY	Output	Clock disable de	nial response.
ACLKQACTIVE	Output	Indicates that the	e CCI-500 requires the ACLK input to run.
PREQ	Input	Request to change power state.	
PSTATE[2:0]	Input	Required power	state.
		The encodings for	or this are:
		0b000	Off.
		0b001	Static snoop filter RAM retention.
		0b010	Reserved.
		0b011	Dynamic snoop filter RAM retention.
		0b100	On.
		0b101-0b111	Reserved.
		If the P channel state.	is not used, you must tie <b>PSTATE</b> to <b>0b100</b> ,On
РАССЕРТ	Output	Power state trans	sition acceptance.

Signal	Direction	Description	
PDENY	Output	Power state transition denial.	
PACTIVE[4:0]	Output	Hint from the CCI-500 to indicate the power states that it can accept.	
		Each bit corresponds to a power state, if HIGH then that state is a legal power transition:	
		[0] Off.	
		[1] Static snoop filter RAM retention.	
		[2] Reserved.	
		[3] Dynamic snoop filter RAM retention.	
		[4] On.	

# Table A-2 Power and clock control signals (continued)

# A.3 Configuration signals

The CCI-500 samples configuration signals only when the **ARESETn** signal transitions from LOW to HIGH.

The following table shows the configuration signals.

## Table A-3 Configuration signals

Signal	Direction	Description
ADDRMAPx[ADDRMAP_WIDTH-1:0]	Input	Configuration inputs that you can use to define the mapping scheme of the address decoder. In the example decoder, there are 3 bits for each of the possible nine address regions.
		Note
		It is the reset sampled version of the ADDRMAP that is passed to the address decode irrespective of whether it is the ARM supplied address map or a modified version.
QOSOVERRIDE[n:0]	Input	If HIGH, the internally generated values override the <b>ARQOS</b> and <b>AWQOS</b> input signals. See <i>2.4.12 Quality of Service</i> on page 2-39 for more information.
		One bit exists for each slave interface.
ACCHANNELENSx[1:0]	Input	AC channel enables, one input per slave interface. These
for ACE interfaces		inputs override any software enables.
ACCHANNELENSx[0]		Bit[0], DVM message enable
for ACE-Lite interfaces This bit is encoded as follows:		This bit is encoded as follows:
		0b0 DVM messages disabled.
		<b>0b1</b> DVM messages enabled.
		Bit[1], Snoop enable
		This bit applies to ACE interfaces only, and is encoded as follows:
		<b>0b0</b> Snoop requests disabled.
		<b>0b1</b> Snoop requests enabled.
		Note
		Snoops and DVM messages must still be enabled in the Snoop Control Registers.

# Table A-3 Configuration signals (continued)

Signal	Direction	Description
ORDERED_WRITE_OBSERVATION[n:0]	Input	Controls whether the slave interface supports the Ordered Write Observation property.
		This bit is encoded as follows:
		<b>0b0</b> Interface does not support Ordered Write Observation.
		<b>0b1</b> Interface supports Ordered Write Observation.
		One bit exists for each slave interface.
BURST_SPLIT_ALL[n:0]	Input	If HIGH, all incoming requests are split into 64-byte transfers, rather than shareable requests only. This signal has no effect on an interface where the SIX_BURST_SPLITTER parameter is set to 0.
		One bit exists for each ACE-Lite slave interface.

# A.4 Debug signals

The inputs can change at runtime and you must synchronize them to the CCI-500 clock to prevent timing hazards.

The following table shows the debug signals.

## Table A-4 Debug signals

Signal	Direction	Description
NIDEN	Input	Non-invasive debug enable. If HIGH, the signal enables counting and export of PMU events.
SPNIDEN	Input	Secure privileged non-invasive debug enable. If both <b>SPNIDEN</b> and <b>NIDEN</b> are HIGH, the signal enables counting of both Non-secure and Secure events.
DBGEN	Input	Invasive debug enable. If HIGH, enables the counting and export of PMU events.
SPIDEN	Input	Secure invasive debug enable. If both <b>SPIDEN</b> and <b>DBGEN</b> are HIGH, enables the counting of both Non-secure and Secure events.
EVNTBUS[n:0]	Output	The CCI-500 events exported if enabled in the PMCR. See <i>PMU</i> event list on page 2-30 for information on pin allocations of this vector.
		The vector width is dependent on the number of master interfaces, M, and the number of slave interfaces, S, and is defined as:
		S*32 + M*7 + 15.
nEVNTCNTOVERFLOW[7:0]	Output	Overflow flags for the PMU clock and counters. This is an active- LOW signal. Each bit represents the overflow for the event counter with that number.
nERRIRQ	Output	Indicates that an error response, DECERR or SLVERR, is received on the <b>RRESP</b> , <b>BRESP</b> , or <b>CRRESP</b> input signals, and it cannot be signaled precisely. If LOW, the signal indicates that an error has occurred.

See the *ARM*<sup>®</sup> *CoreSight*<sup>™</sup> *Architecture Specification* for more information.

# A.5 DFT signals

The CCI-500 uses the *Design For Test* (DFT) signals to communicate with the DFT and MBIST interfaces.

The following table shows the DFT signals.

#### Table A-5 DFT signals

Signal	Direction	Description
DFTRSTDISABLE	Input	Disables reset during scan shift.
DFTCGEN	Input	Assert HIGH during scan shift to enable architectural clock gates for ACLK clocks.
DFTRAMHOLD	Input	Blocks chip select to RAMs to preserve state.
DFTMCPHOLD	Input	Limits number of multi-cycle path toggles during ATPG delay test.
nMBISTRESET0	Input	Resets MBIST mode.
MBISTREQ0	Input	Selects MBIST mode.
MBISTACK0	Output	Acknowledges MBIST mode.

# A.6 APB4 signals

The following table shows the APB4 Slave Interface signals.

Signal	Direction	Description
PADDR[31:0]	Input	Address.
PPROT[2:0]	Input	Protection type.
PSEL	Input	Peripheral select.
PENABLE	Input	Enable for transfer.
PWRITE	Input	Write transaction indicator
PWDATA[31:0]	Input	Write data.
PSTRB[3:0]	Input	Write data strobe.
PREADY	Output	Transfer ready.
PRDATA[31:0]	Output	Read data.
PSLVERR	Output	Error response.

## Table A-6 APB4 signals

# A.7 ACE and ACE-Lite slave interface signals

The CCI-500 has a configurable number of ACE and ACE-Lite slave interfaces. The suffix is Sx, where x is 0-6, depending on the configuration.

This section contains the following subsections:

- A.7.1 Write address channel signals on page Appx-A-85.
- A.7.2 Write data channel signals on page Appx-A-85.
- A.7.3 Write data response channel signals on page Appx-A-86.
- *A.7.4 Read address channel signals* on page Appx-A-86.
- *A.7.5 Read data channel signals* on page Appx-A-87.
- A.7.6 Coherency address channel signals on page Appx-A-87.
- *A.7.7 Coherency response channel signals* on page Appx-A-88.
- *A.7.8 Coherency data channel signals for ACE interfaces* on page Appx-A-88.
- *A.7.9 Acknowledge signals for ACE interfaces* on page Appx-A-88.

# A.7.1 Write address channel signals

These signals carry control information that describes the nature of the data to be transferred. The data is transferred between master and slave using either a read data channel or a write data channel.

The following table shows the write address channel signals.

## Table A-7 Write address channel signals

Signal	Direction	Description
AWIDSx[n:0]	Input	Write address ID. You can configure the width of this signal.
AWADDRSx[n:0]	Input	Write address. You can configure the CCI-500 to support between a 32-bit and a 44-bit signal width.
AWREGIONSx[3:0]	Input	Write address region. You can tie this signal LOW if the master does not drive it.
AWLENSx[7:0]	Input	Write burst length.
AWSIZESx[2:0]	Input	Write burst size.
AWBURSTSx[1:0]	Input	Write burst type.
AWLOCKSx	Input	Write lock type.
AWCACHESx[3:0]	Input	Write cache type.
AWPROTSx[2:0]	Input	Write protection type.
AWSNOOPSx[2:0]	Input	Write snoop request type.
AWDOMAINSx[1:0]	Input	Write domain.
AWQOSSx[3:0]	Input	Write QoS value.
AWUSERSx[n:0]	Input	Specified extension to AW payload. You can configure the width of this signal.
NSAIDWSx[3:0]	Input	Optional extension to AW payload, that transmits the Non-secure access identifier for a request.
AWVALIDSx	Input	Write address valid.
AWREADYSx	Output	Write address ready.

## A.7.2 Write data channel signals

Write data channel signals carry the write data from the master to the slave, and include the data bus and a byte lane strobe signal.

The following table shows the write data channel signals.

### Table A-8 Write data channel signals

Signal	Direction	Description
WDATASx[127:0]	Input	Write data.
WSTRBSx[15:0]	Input	Write byte-lane strobes.
WLASTSx	Input	Write last. This signal indicates the last transfer in a write burst.
WUSERSx[n:0]	Input	The specified extension to the W payload.
WCHECKSUMSx[n:0]	Input	An optional extension to the W payload that can transmit checksum or parity information for the data. You can configure the width of this signal.
WVALIDSx	Input	Write data is valid.
WREADYSx	Output	Write data is ready.

## A.7.3 Write data response channel signals

A slave uses the write response channel to respond to write transactions. All write transactions require completion signaling on the write response channel.

The following table shows the write data response channel signals.

#### Table A-9 Write data response channel signals

Signal	Direction	Description
BIDSx[n:0]	Output	Write response ID. You can configure the width.
BRESPSx[1:0]	Output	Write response.
BUSERSx[n:0]	Output	The specified extension to the B payload.
BVALIDSx	Output	Write response is valid.
BREADYSx	Input	Write response is ready.

# A.7.4 Read address channel signals

The following table shows the read address channel signals.

#### Table A-10 Read address channel signals

Signal	Direction	Description
ARIDSx[n:0]	Input	Read address ID. You can configure the width of this signal.
ARADDRSx[n:0]	Input	Read address. You can configure the CCI-500 to support 40-bit [39:0], or 44-bit [43:0] DVM transactions.
ARREGIONSx[3:0]	Input	Read address region. You can tie this signal LOW if the master does not drive it.
ARLENSx[7:0]	Input	Read burst length.
ARSIZESx[2:0]	Input	Read burst size.
ARBURSTSx[1:0]	Input	Read burst type.
ARLOCKSx	Input	Read lock type.

#### Table A-10 Read address channel signals (continued)

Direction	Description
Input	Read cache type.
Input	Read protection type.
Input	Read domain.
Input	Read snoop request type.
Input	Read QoS.
Input	The specified extension to the AR payload.
Input	Optional extension to AR payload, that transmits the Non-secure access identifier for a request.
Input	Read address is valid.
Output	Read address is ready.
	Input Input Input Input Input Input Input Input Input Input

#### A.7.5 Read data channel signals

Read data channel signals carry the read data and the read response information from the slave to the master, and include the data bus and a read response signal.

The following table shows the read data channel signals.

#### Table A-11 Read data channel signals

Signal	Direction	Description	
RIDSx[n:0]	Output	Read data ID. You can configure the width of this signal.	
RDATASx[127:0]	Output	Read data.	
RRESPSx[3:0]	Output	Read data response for ACE interfaces S3 and S4.	
RRESPSx[1:0]	Output	Read data response for ACE-Lite interfaces S0, S1, and S2.	
RLASTSx	Output	Read last. This signal indicates the last transfer in a read burst.	
RUSERSx[n:0]	Output	The specified extension to the R payload.	
RCHECKSUMSx[n:0]	Output	An optional extension to the R payload that can transmit checksum or parity information for the data.	
		You can configure the width of this signal.	
RVALIDSx	Output	Read data is valid.	
RREADYSx	Input	Read data is ready.	

## A.7.6 Coherency address channel signals

Coherency address channel signals provide address and associated control information for snoop transactions.

The following table shows the coherency address channel signals.

Table A-12 Coherency address channel signals

Signal	Direction	Description	
ACADDRSx[n:0]	Output	Snoop address. You can configure the CCI-500 to support 40-bit [39:0], or 44-bit [43:0] snoop transactions.	
ACPROTSx[2:0]	Output	Snoop protection type.	
ACSNOOPSx[3:0]	Output	Snoop request type.	
ACVALIDSx	Output	Snoop address is valid.	
ACREADYSx	Input	The master interface is ready to accept the snoop address.	

## A.7.7 Coherency response channel signals

Coherency response channel signals provide the response to snoop transactions.

The following table shows the coherency response channel signals.

#### Table A-13 Coherency response channel signals

Signal	Direction	Description	
CRRESPSx[4:0]	Input	Snoop response.	
NSAIDCRSx[3:0]	Input	Optional extension to CR payload, that transmits the Non-secure access identifier for a snoop response.	
CRVALIDSx	Input	Snoop response is valid.	
CRREADYSx	Output	The slave interface is ready to accept the snoop response.	

# A.7.8 Coherency data channel signals for ACE interfaces

Coherency data channel signals, if used, pass snoop data out from an ACE master.

The following table shows the coherency data channel signals for ACE interfaces.

# Table A-14 Coherency data channel signals, ACE interfaces

Signal	Direction	Description
CDDATASx[127:0]	Input	Snoop data.
CDLASTSx	Input	Snoop data last transfer.
CDCHECKSUMSx[n:0]	Input	An optional extension to the CD payload that can transmit checksum or parity information for the data.
		You can configure the width of this signal.
CDVALIDSx	Input	Snoop data is valid.
CDREADYSx	Output	The slave interface is ready to accept snoop data.

# A.7.9 Acknowledge signals for ACE interfaces

Acknowledge signals indicate that a master or slave has completed a read or write transaction.

The following table shows the acknowledge signals for ACE interfaces.

Table A-15 Acknowledge signals, ACE interfaces

Signal	Direction	Description
RACKSx	Input	Read acknowledge.
WACKSx	Input	Write acknowledge.

# A.8 AXI Master interface signals

The CCI-500 has a configurable number of AXI4 master interfaces. The suffix is My, where y is 0-5, depending on the configuration.

This section contains the following subsections:

- A.8.1 Write address channel signals on page Appx-A-90.
- A.8.2 Write data channel signals on page Appx-A-90.
- A.8.3 Write data response channel signals on page Appx-A-91.
- A.8.4 Read address channel signals on page Appx-A-91.
- *A.8.5 Read data channel signals* on page Appx-A-92.

#### A.8.1 Write address channel signals

These signals carry control information that describes the nature of the data to be transferred. The data is transferred between master and slave using either a read data channel or a write data channel.

The following table shows the write address channel signals.

#### Table A-16 Write address channel signals

Signal	Direction	Description	
AWIDMy[n:0]	Output	Write address ID. The width is the maximum <b>AWID</b> signal width across the slave interfaces + 3 bits, and it is a minimum of 9 bits.	
AWADDRMy[n:0]	Output	Write address. You can configure the CCI-500 to support between a 32-bit and a 44-bit signal width.	
AWREGIONMy[3:0]	Output	Write address region.	
AWLENMy[7:0]	Output	Write burst length.	
AWSIZEMy[2:0]	Output	Write burst size.	
AWBURSTMy[1:0]	Output	Write burst type.	
AWLOCKMy	Output	Write lock type.	
AWCACHEMy[3:0]	Output	Write cache type.	
AWPROTMy[2:0]	Output	Write protection type.	
AWQOSMy[3:0]	Output	Write QoS value.	
AWUSERMy[n:0]	Output	The specified extension to the AW payload.	
AWVALIDMy	Output	Write address is valid.	
NSAIDWMy[3:0]	Output	Optional extension to AW payload, that transmits the Non-secure access identifier for a request.	
AWREADYMy	Input	Write address is ready.	

## A.8.2 Write data channel signals

Write data channel signals carry the write data from the master to the slave, and include the data bus and a byte lane strobe signal.

The following table shows the write data channel signals where y represents the master interface number.

#### Table A-17 Write data channel signals

Signal	Direction	Description
WIDMy[n:0]	Output	Write ID tag. Included to help connection to AXI3 slave interfaces.
WDATAMy[127:0]	Output	Write data.
WSTRBMy[15:0]	Output	Write strobes.
WLASTMy	Output	Write last.
WUSERMy[n:0]	Output	User signal.
WCHECKSUMMy[n:0]	Output	An optional extension to the W payload that you can use to transmit checksum or parity information for the data. You can configure the width of this signal.
WVALIDMy	Output	Write valid.
WREADYMy	Input	Write ready.

#### A.8.3 Write data response channel signals

A slave uses the write response channel to respond to write transactions. All write transactions require completion signaling on the write response channel.

The following table shows the write data response channel signals.

#### Table A-18 Write data response channel signals

Signal	Direction	Description
BIDMy[n:0]	Input	Write response ID.
BRESPMy[1:0]	Input	Write response.
BUSERMy[n:0]	Input	The specified extension to the B payload.
BVALIDMy	Input	Write response is valid.
BREADYMy	Output	Write response is ready.

#### A.8.4 Read address channel signals

These signals carry control information that describes the nature of the data to be transferred. The data is transferred between master and slave using either a read data channel or a write data channel.

The following table shows the read address channel signals.

#### Table A-19 Read address channel signals

Signal	Direction	Description	
ARIDMy[n:0]	Output	Read address ID. The width is the maximum <b>ARID</b> signal width across slave interfaces + 3 bi and it is a minimum of 6 bits.	
ARADDRMy[n:0]	Output	Read address. You can configure the CCI-500 to support between a 32-bit and a 44-bit signal width.	
ARREGIONMy[3:0]	Output	Read address region.	
ARLENMy[7:0]	Output	Read burst length.	
ARSIZEMy[2:0]	Output	Read burst size.	
ARBURSTMy[1:0]	Output	Read burst type.	

## Table A-19 Read address channel signals (continued)

Signal	Direction	Description
ARLOCKMy	Output	Read lock type.
ARCACHEMy[3:0]	Output	Read cache type.
ARPROTMy[2:0]	Output	Read protection type.
ARQOSMy[3:0]	Output	Read QoS value.
ARUSERMy[n:0]	Output	The specified extension to the AR payload.
ARVALIDMy	Output	Read address is valid.
NSAIDRMy[3:0]	Output	Optional extension to AR payload, that transmits the Non-secure access identifier for a request
ARREADYMy	Input	Read address is ready.

# A.8.5 Read data channel signals

Read data channel signals carry the read data and the read response information from the slave to the master.

The following table shows the read data channel signals.

#### Table A-20 Read data channel signals

Signal	Direction	Description
RIDMy[n:0]	Input	Read data ID.
RDATAMy[127:0]	Input	Read data.
RRESPMy[3:0]	Input	Read data response.
RLASTMy	Input	Read last. This signal indicates the last transfer in a read burst.
RUSERMy[n:0]	Input	The specified extension to the R payload.
RCHECKSUMMy[n:0]	Input	An optional extension to the R payload that you can use to transmit checksum or parity information for the data. You can configure the width of this signal.
RVALIDMy	Input	Read data is valid.
RREADYMy	Output	Read data is ready.

# A.9 Miscellaneous signals

Some signals in the CCI-500 are not applicable for general operation, implementation, or debug of the CCI-500. However, if there is a problem to resolve, you might require these signals.

The following table shows a signal that you can use to indicate an ECO applies to this implementation of the CCI-500.

### Table A-21 Miscellaneous signals

Signal	Direction	Description	
ECOREVNUM[3:0]	1	Use this signal to update the revision field register if you have an ECO to apply to the CCI-500. See <i>ECO revision number</i> on page 3-62.	

# Appendix B **Revisions**

This appendix describes the technical changes between released issues of this book.

It contains the following sections:

• *B.1 Revisions* on page Appx-B-95.

# B.1 Revisions

This appendix describes the technical changes between released issues of this book.

#### Table B-1 Issue 00

Change	Location	Affects
First release	-	-

#### Table B-2 Differences between Issue 00 and Issue 01

Change	Location	Affects
Clarified information about adding and removing masters from the coherent domain.	2.4.3 Snoop connectivity and control on page 2-28.	All revisions.
Updated information about relationship between debug enable inputs and PMCR settings.	3.3.5 Performance Monitor Control Register (PMCR) on page 3-60.	All revisions.